Thomas Fernique

Moscow, Spring 2011

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3 Random assembly

4 Random sampling

Random assembly

Random sampling

Quenching (first quasicrystals)



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Free energy minimization

Stability: minimal free energy F = E - TS

Because of the high temperature in the melt and the rapid cooling:

Minimal free energy $F \Leftrightarrow Maximal \text{ entropy } S$.

Matching rules modelling interactions (the energy E): forgotten!

Entropy for tilings?

Configuration entropy

Entropy of the tilings of a domain $D \subset \mathbb{R}^2$:

 $S := \log(\text{number of tilings of } D).$

Entropy per tile of a tiling T of D:

$$s(T):=\frac{S(T)}{|T|}.$$



Example: compute the entropy of these tilings.

Typical tilings of maximal entropy

We are mainly interested in two questions.

- Which domains do maximize the entropy?
- Ob the tilings of such domains have a "typical look"?

Typical look: properties satisfied by most of large tilings, e.g.,

- presence (or not) of some patterns;
- proportions of different tiles (phason strain);
- perp-space fluctuations;
- . . .

 \rightsquigarrow These are properties that one can expect after quenching!

A Simple case

Let W_n be the set of words of length *n* over the alphabet $\{1, 2\}$.

Entropy of $w \in W_n$, with $|w|_1 = a$ and $|w|_2 = b$:

$$s(w) = \frac{1}{n} \log\left(\frac{n!}{a!b!}\right).$$

Words of maximal entropy: a = b (balanced words).

Random balanced words are known to have fluctuations in \sqrt{n} .

 \rightsquigarrow seen as broken lines on \mathbb{Z}^2 , they stay close to the line y = x.

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Further cases?



Which of these tilings by Penrose rhombi has the largest entropy?

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Further cases?



How close to a Penrose tiling can be expected to be a random tiling?











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Dimer tilings



Dimer tiling: pairing of adjacent cells of a part of the triangular grid.

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Dimer tilings



Dimer tiling: pairing of adjacent cells of a part of the triangular grid.

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Dimer tilings



3-dimensional viewpoint: remind the shadows of rhombus tilings!

Random assembly

Random sampling

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Counting with determinants



Bottom edges of vertical tiles \rightsquigarrow family of non-intersecting paths.

Random assembly

Random sampling

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Counting with determinants



Bottom edges of vertical tiles \rightsquigarrow family of non-intersecting paths.

Random assembly

Random sampling

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Counting with determinants



Conversely: family of non-intersecting paths \rightsquigarrow unique dimer tiling.

Random assembly

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Counting with determinants



Conversely: family of non-intersecting paths \rightsquigarrow unique dimer tiling.

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Counting with determinants



Let $N := det(\lambda_{i,j})$, where $\lambda_{i,j}$ is the number of paths from *i* to *j*.

Random assembly

Random sampling

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Counting with determinants



Theorem: there are N distinct families of non-intersecting paths.

Random assembly

Random sampling

Counting with determinants



 $\sum_{\sigma \in S_n} \varepsilon(\sigma) \lambda_{1,\sigma(1)} \cdots \lambda_{n,\sigma(n)}$: weighted sum of all the path families.

Random assembly

Random sampling

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Counting with determinants



Order grid vertices. Consider the first intersection of a "bad" family.

Random assembly

Random sampling

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Counting with determinants



Uncross paths at this vertex ~>> pairing with another bad family.

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Counting with determinants



Paired families have opposite signature $\rightsquigarrow N$ counts the good ones!

The Dimer case

Random assembly

Random sampling

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Entropy of boxed tilings (MacMahon 1916, Elser 1984)



The Dimer case

Random assembly

Random sampling

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Random assembly

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The Dimer case

Random assembly

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Entropy of boxed tilings (MacMahon 1916, Elser 1984)



The Dimer case

Random assembly

Random sampling

Entropy of boxed tilings (MacMahon 1916, Elser 1984)



	252	210	120	45	10
	210	252	210	120	45
N=	120	210	252	210	120
	45	120	210	252	210
	10	45	120	210	252

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The Dimer case

Random assembly

Random sampling

Entropy of boxed tilings (MacMahon 1916, Elser 1984)



252	210	120	45	10
210	252	210	120	45
120	210	252	210	120
45	120	210	252	210
10	45	120	210	252
	252 210 120 45 10	 252 210 210 252 120 210 45 120 10 45 	252210120210252210120210252451202101045120	25221012045210252210120120210252210451202102521045120210

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Entropy of this size 75 tiling? $s = \frac{\log(N)}{75} = \frac{\log(267\,227\,532)}{75} \simeq 0.259$.

Entropy of boxed tilings (MacMahon 1916, Elser 1984)



For $a \ge b \ge c$: $N_{a,b,c} = \left| \begin{pmatrix} a+b \\ a+i-j \end{pmatrix}_{1 \le i,j \le c} \right|$

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Entropy of boxed tilings (MacMahon 1916, Elser 1984)



For
$$a \ge b \ge c$$
:
 $N_{a,b,c} = \left| \begin{pmatrix} a+b\\ a+i-j \end{pmatrix}_{1 \le i,j \le c} \right|$
 $N_{a,b,c} = \prod_{i=1}^{a} \prod_{j=1}^{b} \prod_{k=1}^{c} \frac{i+j+k-1}{i+j+k-2}$

Entropy of this size $3n^2$ tiling? $s_{a,b,c} \leq s_{n,n,n} \xrightarrow[n\infty]{} \log \frac{3\sqrt{3}}{4} \simeq 0.262$.

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The Arctic circle (Cohn-Larsen-Propp, 1997)



The Arctic circle (Cohn-Larsen-Propp, 1997)



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The Arctic circle (Cohn-Larsen-Propp, 1997)



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The Arctic circle (Cohn-Larsen-Propp, 1997)


The Arctic circle (Cohn-Larsen-Propp, 1997)



Theorem

Normalized random boxed tilings converge in probability towards a limit surface (exponentially fast). This surface is characterized:

• flat outside an arctic circle;

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subtly wavy inside it.

Do the boxed tilings have a typical look?

A Variational principle (Cohn-Kenyon-Propp, 2001)

Height function of a tiling: distance to the plane x + y + z = 0.

Theorem

Let $R \subset \mathbb{R}^2$ be bounded by a piecewise smooth simple closed curve. If, for $n \ge 0$, R_n is a tileable domain which approximates nR, then

$$\lim_{n\to\infty} s(R_n) = \sup_{h} \frac{1}{|R|} \iint_R \operatorname{ent}\left(\frac{\partial h}{\partial x}, \frac{\partial h}{\partial y}\right) \, dx \, dy,$$

where ent : $\mathbb{R}^2 \to \mathbb{R}$ and h is any 2-Lipschitz real function on R. Moreover, the normalized R_n 's random height functions converge in probability (exponentially fast) towards the integral-maximizing h.

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where ent : $\mathbb{R}^2 \to \mathbb{R}$ and h is any 2-Lipschitz real function on R. Moreover, the normalized R_n 's random height functions converge in probability (exponentially fast) towards the integral-maximizing h.

Note: this does not tell anything about the integral-maximizing h. However, one checks that ent is concave with a maximum in (0,0).

Dimer tilings of maximal entropy



The entropy is thus maximal for flat boundary height functions. Typical dimer tilings then stay close to the plane x + y + z = 0. The Dimer case

Random assembly

Random sampling

Dimer tilings of maximal entropy



Entropy:
$$s_{\text{flat}} = -\frac{2}{\pi} \int_0^{\frac{\pi}{3}} \log(2\cos(t)) dt \simeq 0.338 > s_{\text{boxed}} = 0.262.$$

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1 Random tilings







Principle

Energy:

Can explain quasiperiodic tiling stability at low temperature. Mostly fails to explain growth of quasiperiodic tiling.

Entropy:

Can explain random tiling stability at high temperature. Some random tilings may be close to quasiperiodic tilings.

Question:

Can we explain the growth of random tilings?

A simple case

Growth of a two-letter word via a Bernoulli process of parameter *p*:





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A simple case

Growth of a two-letter word via a Bernoulli process of parameter *p*:

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A simple case

Growth of a two-letter word via a Bernoulli process of parameter *p*:

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A simple case

Growth of a two-letter word via a Bernoulli process of parameter *p*:



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A simple case

Growth of a two-letter word via a Bernoulli process of parameter *p*:



This is a biased random walk which, once normalized, converges in probability (exponentially fast) towards a line segment (slope $\frac{1-p}{p}$).

The growth is local and yields either periodic or non-periodic lines! However, the stability decreases when the slope goes away from 1.

A more complex case (Elser-Joseph 1997)

Chemical potential: parameter $\mu > 0$.

Energy of a tiling (\sim surface tension):

$$E := \sum_{x \text{ vertex}} \mu - \theta(x),$$

where $\theta(x)$ denotes the total angle subtended by tiles around x.

Growth: at each step, add or remove a randomly chosen edge with probability 1 if *E* decreases or $\exp(-\Delta E/T)$ if *E* increases by ΔE .

Claim: the growth rate tends to decrease with μ , increase with T.



Growth on a cylinder with seed base. A fast growth can create tears.

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Tears may be healed by melting, *i.e.*, by playing suitably on T and μ .

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Tears may be healed by melting, *i.e.*, by playing suitably on T and μ .

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Tears may be healed by melting, *i.e.*, by playing suitably on T and μ .

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Some comments

Very handwavy framework, no formal result or even conjecture.

Simulations yield interesting tilings which look like random. To what extent? Are they entropically stabilized?

Growth governed by energy minimization at positive temperature. But should not entropy maximization play a role in such a growth?

Entropy could help to explain growth because of its non-locality. But it is not easy to compute ΔS when adding/removing a tile...

Lot of open questions!

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2 The Dimer case

3 Random assembly



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Problem



How to draw uniformly at random such rhombus tilings?

Flip (or phason-flip)

Definition (Flip)

A flip is a 180° rotation of three rhombi which form a hexagon.



Theorem (Kenyon 1993)

Rhombus tilings of a simply connected domain are flip-connected.

Discrete-time Markov chain on rhombus tilings: Choose uniformly at random a vertex, a direction, and try to flip.





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Discrete-time Markov chain on rhombus tilings: Choose uniformly at random a vertex, a direction, and try to flip.





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Discrete-time Markov chain on rhombus tilings: Choose uniformly at random a vertex, a direction, and try to flip.



Flip-connectedness \rightsquigarrow ergodic chain \rightsquigarrow convergence in probability. Symmetric chain \rightsquigarrow uniform stationary (equilibrium) distribution.

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Flip-connectedness \rightsquigarrow ergodic chain \rightsquigarrow convergence in probability. Symmetric chain \rightsquigarrow uniform stationary (equilibrium) distribution.

Mixing time

Total variation between two distributions:

$$||\mu - \nu|| := \frac{1}{2} \sum_{x \in \Omega} |\mu(x) - \nu(x)|.$$

Measure of the convergence towards the stationary distribution:

$$d(t) = \max_{x \in \Omega} ||P^t(x, \cdot) - \pi||.$$

Mixing time:

$$\tau := \min\{t \mid d(t) \leq 1/(2e)\}.$$

Theorem (Half-life)

For an ergodic Markov chain, one has: $d(t) \leq \exp(-\lfloor t/\tau \rfloor)$.

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The two-letter case



On a length n word: choose a vertex, a direction, and try to flip.

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The two-letter case



On a length n word: choose a vertex, a direction, and try to flip.

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The two-letter case



On a length n word: choose a vertex, a direction, and try to flip.

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The two-letter case



Coupling of two words x and $y \rightsquigarrow$ coalescence time T_{xy} .
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The two-letter case



Coupling of two words x and $y \rightsquigarrow$ coalescence time T_{xy} .

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The two-letter case



Theorem: $\tau = O(\mathbb{E}T)$, where $T = \max_{x,y} T_{xy}$.

The two-letter case



If x is completely below y at time t, then this holds for t' > t.

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The two-letter case



Top and bottom words thus sandwich all the other words $\rightsquigarrow T$.

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The two-letter case



Let A_t be the sandwiched area at time t. $A_0 = n^2$. $A_T = 0$.

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The two-letter case



Let $\Delta A_t := A_{t+1} - A_t$. Claim: $\mathbb{E}(\Delta A_t | A_t) \leq 0$ (count flips by type).

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The two-letter case



Theorem: $(\mathbb{E}(\Delta A_t) \leq 0, \operatorname{Pr}(|\Delta A_t| \geq 1) > \alpha) \Rightarrow \mathbb{E}T \leq \frac{\max A_t^2}{\alpha}.$

The two-letter case



Here: $\max A_t^2 = A_0^2 = n^4$ and $\alpha \ge \frac{1}{2n} \rightsquigarrow \mathbb{E}T \le n^5 \rightsquigarrow \tau = O(n^5)$.

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The dimer case (Luby-Randall-Sinclair 1995)



Dimer \sim family of non-intersecting paths \rightsquigarrow *n* two-letter words.

The dimer case (Luby-Randall-Sinclair 1995)



Flip on the dimer \rightsquigarrow flip on one of these words.

The dimer case (Luby-Randall-Sinclair 1995)



Flip on the dimer \rightsquigarrow flip on one of these words.

The dimer case (Luby-Randall-Sinclair 1995)



Flip on the dimer \rightsquigarrow flip on one of these words.

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The dimer case (Luby-Randall-Sinclair 1995)



The converse does not always hold \rightsquigarrow Markov chain modification.

The dimer case (Luby-Randall-Sinclair 1995)



Choose a vertex, a direction, and try to flip a tower.

The dimer case (Luby-Randall-Sinclair 1995)



Choose a vertex, a direction, and try to flip a tower.

The dimer case (Luby-Randall-Sinclair 1995)



Choose a vertex, a direction, and try to flip a tower.

The dimer case (Luby-Randall-Sinclair 1995)



More precisely, flip a tower of height h with probability $\frac{1}{h}$.

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The dimer case (Luby-Randall-Sinclair 1995)



Still ergodicity and symmetric \rightsquigarrow uniform stationary distribution.

The dimer case (Luby-Randall-Sinclair 1995)



 $\mathsf{Coupling} \rightsquigarrow V_t := \sum_i A^i_t \rightsquigarrow \mathbb{E}(\Delta V_t | V_t) \leq 0 \text{ thanks to the factor } \tfrac{1}{h}.$

Random assembly

Random sampling

The dimer case (Luby-Randall-Sinclair 1995)



Here: max $V_t \leq V_0 = n^3$ and $\Pr(|\Delta V_t| \geq 1) \geq \frac{1}{2n} \rightsquigarrow \tau = O(n^7)$.

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Further cases? Other methods?

Tight bounds (Wilson 2001):

- two-letter case: $\tau = \Theta(n^3 \log(n));$
- dimer case: $\tau = \Theta(n^4 \log(n))$.

Open question: Mixing times for general rhombus tilings?

Random sampling without flip?

- two-letter words: k updates $1 \rightarrow 2$ uniformly at random on 1^n ;
- random boxed dimer tiling in $O(n^5)$ (Borodin-Gorin 2008);
- for more general rhombus tilings?

Some references for this lecture:

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These slides and the above references can be found there:

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http://www.lif.univ-mrs.fr/~fernique/qc/
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