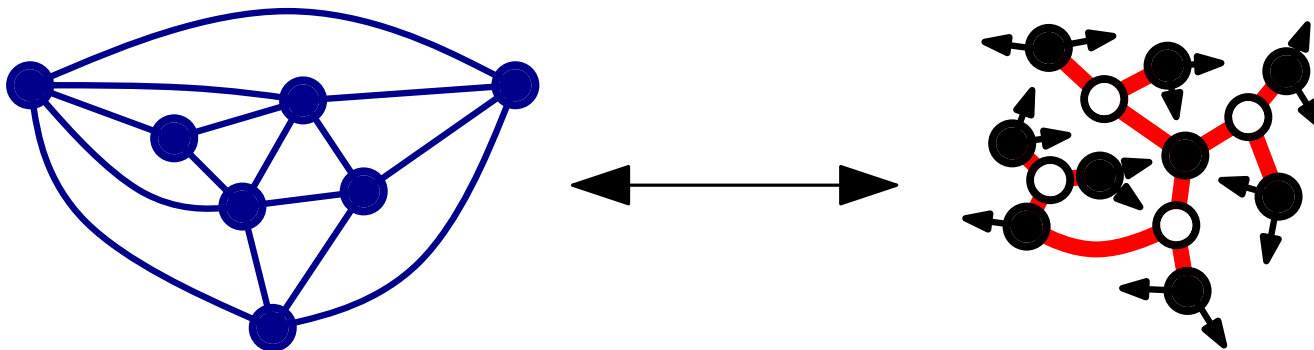


Bijections for planar maps with boundaries

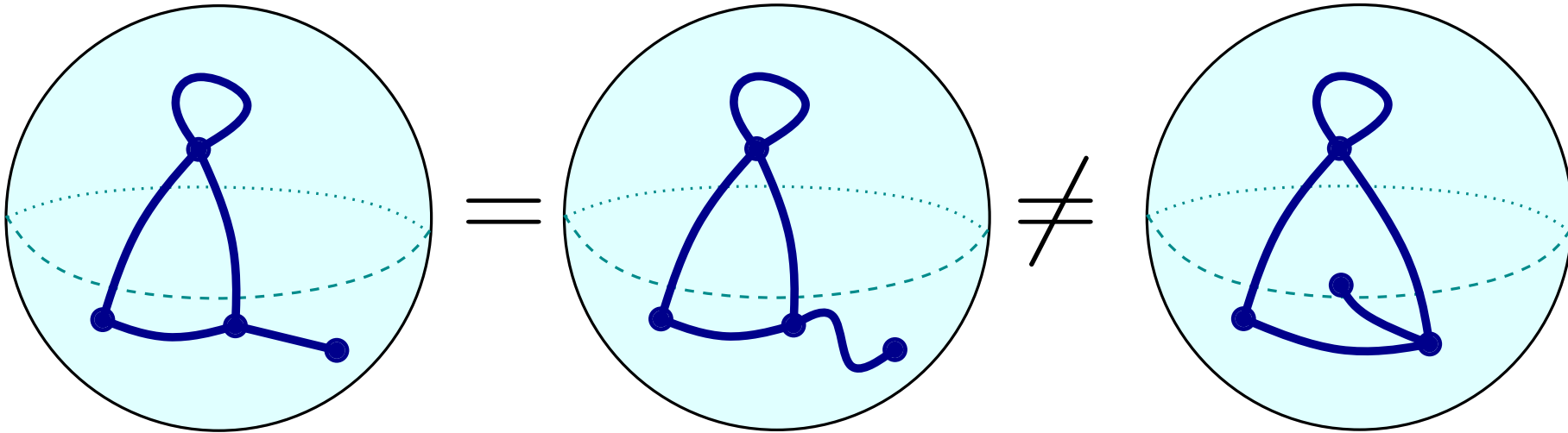
Éric Fusy (CNRS/LIX)

Joint work with Olivier Bernardi (Brandeis Univ.)

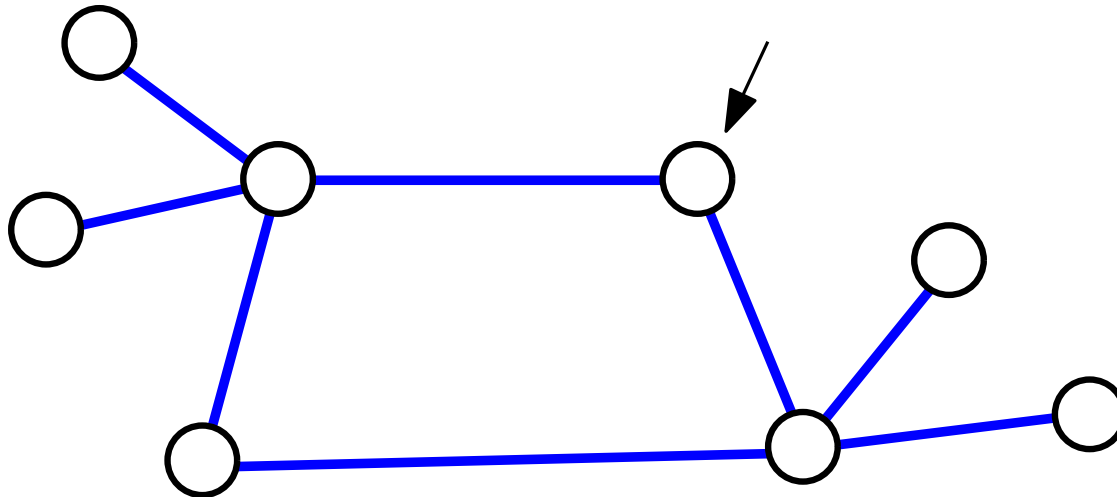


Planar maps

- **Planar map** = connected graph embedded on the **sphere**, considered up to continuous deformation



- **Rooted map** = map with a marked corner



A rooted map

Counting formulas for rooted maps

- **Beautiful counting formulas** discovered by Tutte

Maps with n edges

$$\frac{2 \cdot 3^n \cdot (2n)!}{n!(n+2)!}$$

Bipartite maps
with n edges

$$\frac{3 \cdot 2^{n-1} \cdot (2n)!}{n!(n+2)!}$$

2-connected maps
with n edges

$$\frac{4 \cdot (3n-3)!}{(n-1)!(2n)!}$$

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- **Tutte's slicings formula (1962):**

Let $B[n_1, n_2, \dots, n_k]$ be the number of rooted bipartite maps with n_i faces of degree $2i$ for $i \in [1..k]$. Then

$$B[n_1, \dots, n_k] = 2 \frac{e!}{v!} \prod_{i=1}^k \frac{1}{n_i!} \binom{2i-1}{i-1}^{n_i}$$

where $e = \#edges = \sum_i i n_i$ and $v = \#vertices = e - k + 2$

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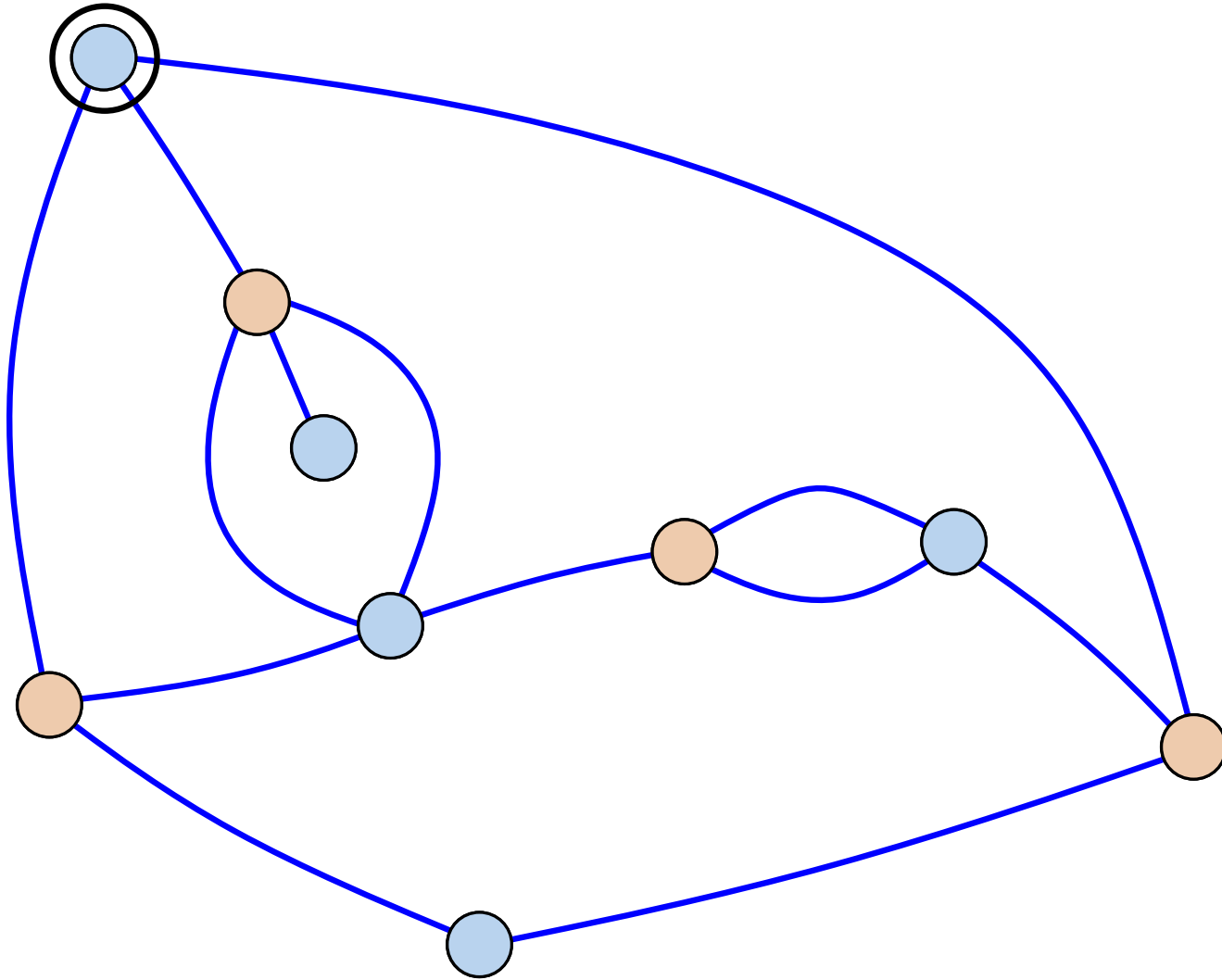
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Counting methods: recursive method, matrix integrals, **bijections**

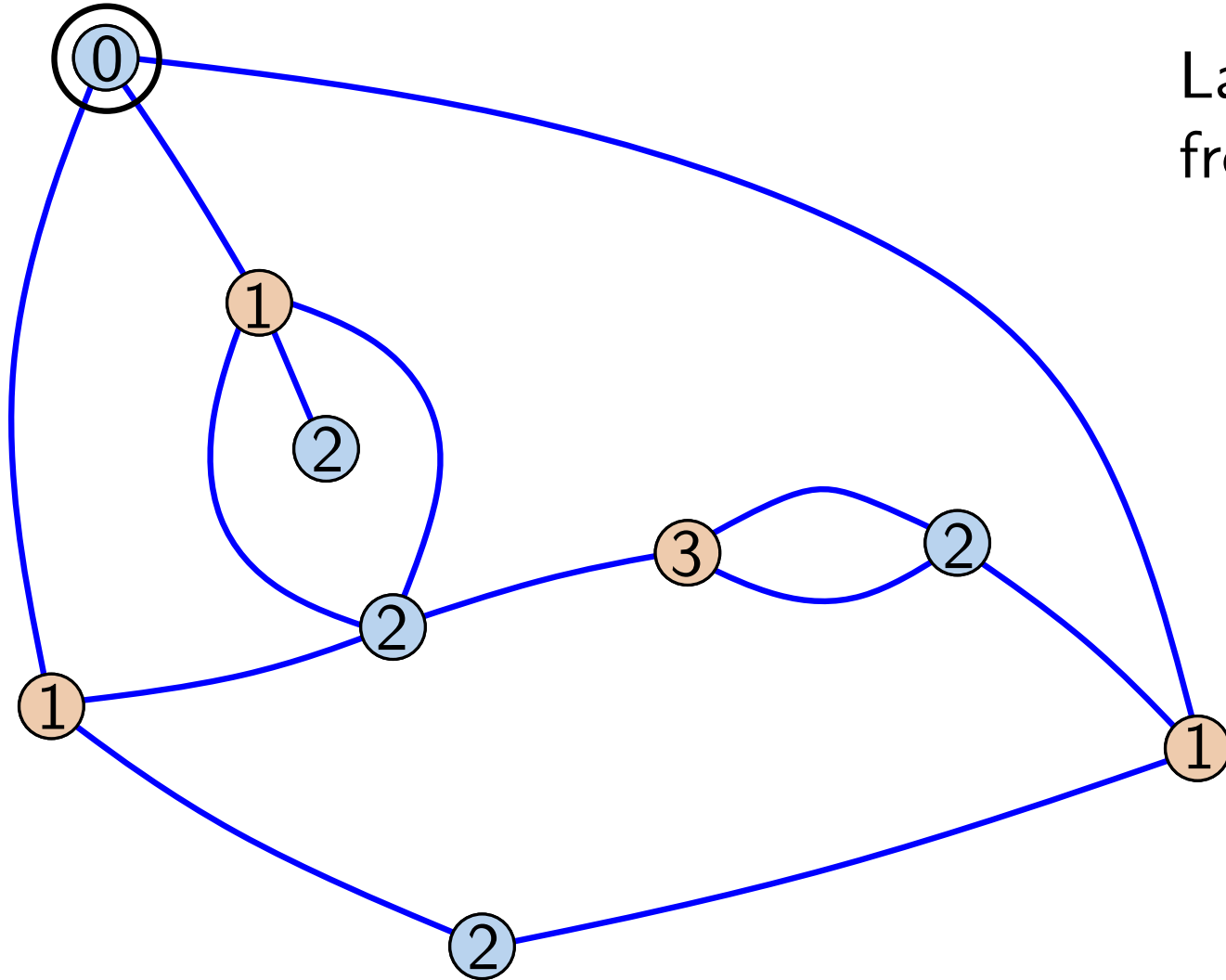
The BDG bijection for pointed bipartite maps

[Bouttier, Di Francesco, Guitter'04]



The BDG bijection for pointed bipartite maps

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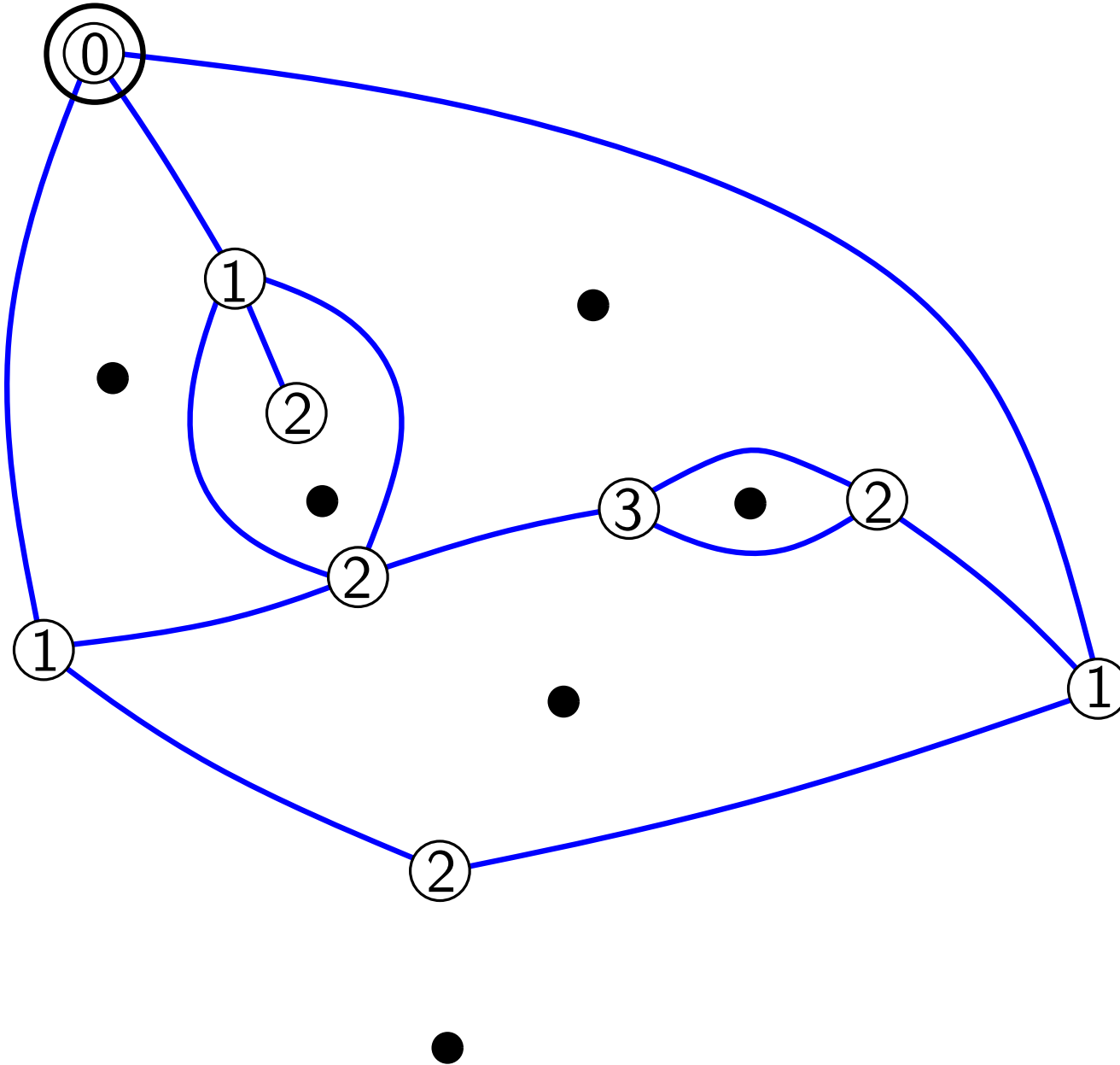
Label vertices by distance from the marked vertex

The BDG bijection for pointed bipartite maps

[Bouttier, Di Francesco, Guitter'04]

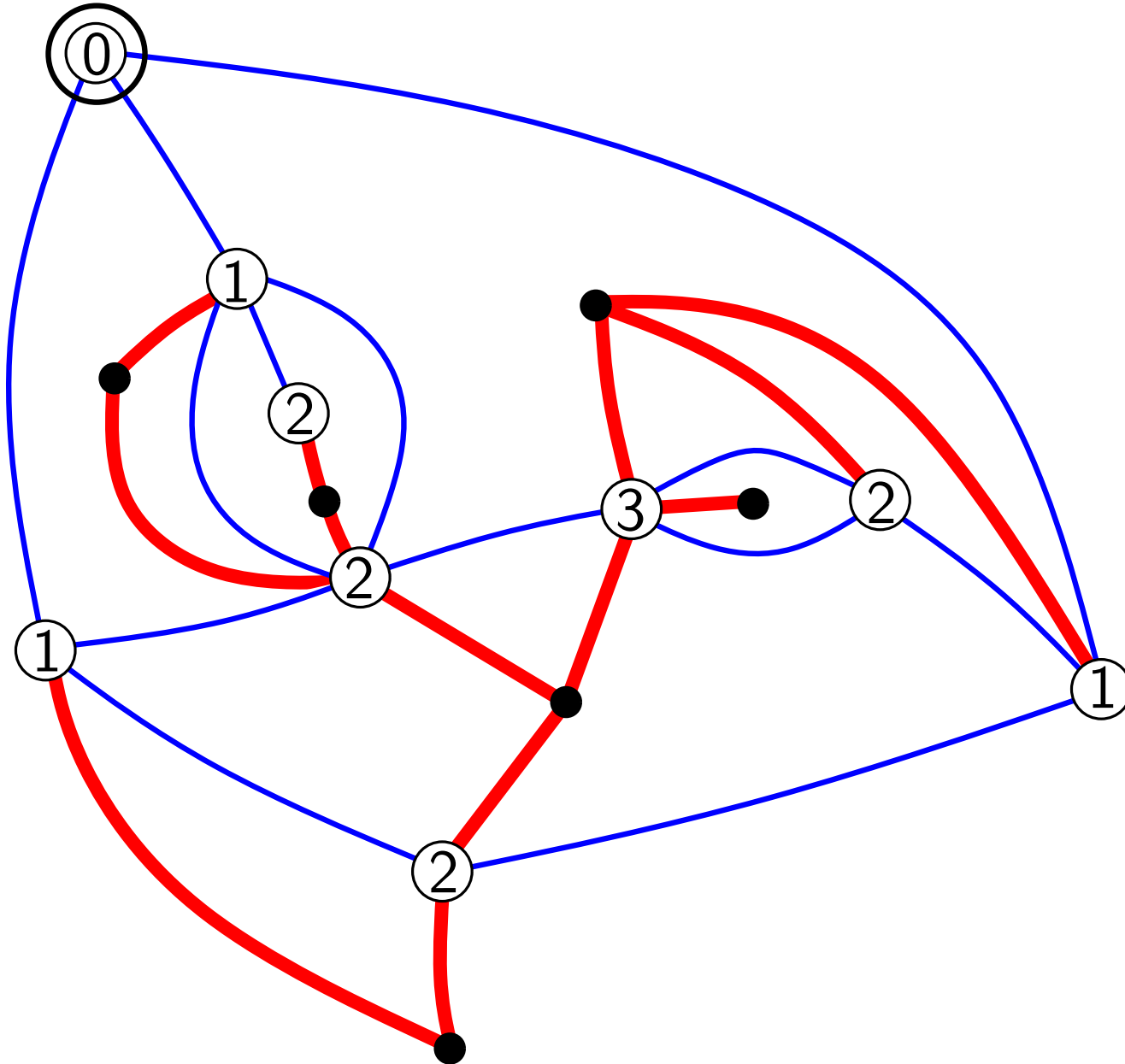
**Construction of a
labeled mobile**

(i) Add a black vertex
in each face



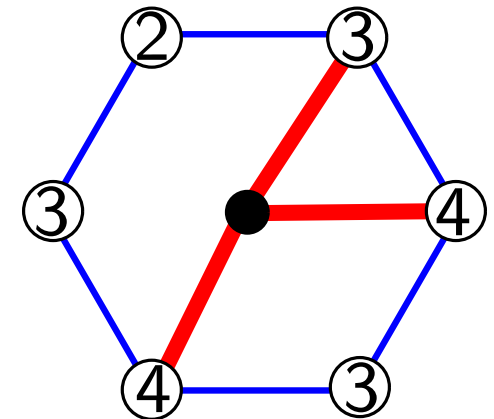
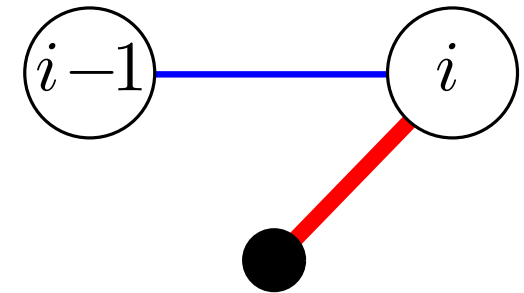
The BDG bijection for pointed bipartite maps

[Bouttier, Di Francesco, Guitter'04]



Construction of a labeled mobile

- (i) Add a black vertex in each face
- (ii) Each map-edge gives a mobile-edge using the local rule

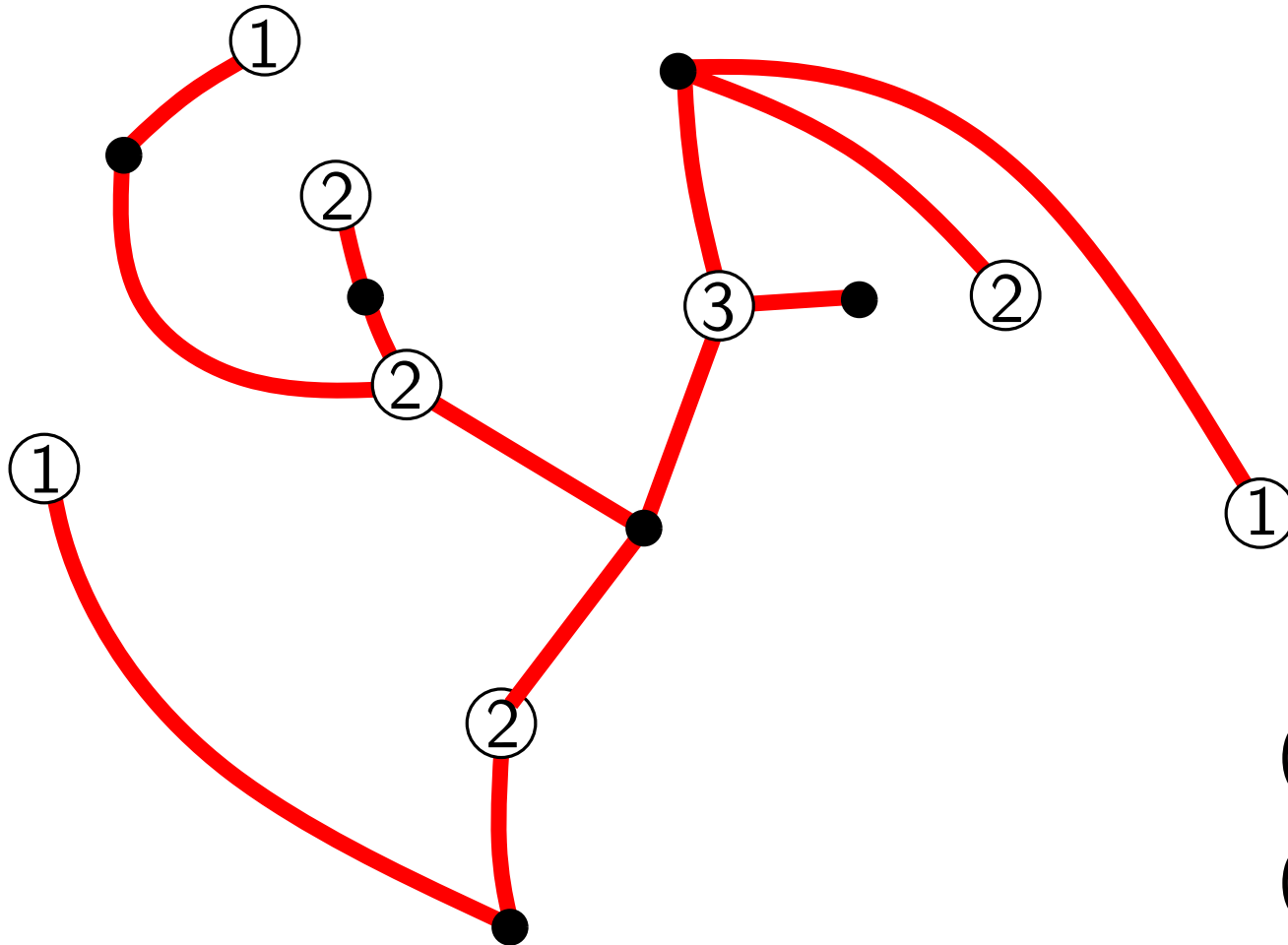


The BDG bijection for pointed bipartite maps

[Bouttier, Di Francesco, Guitter'04]

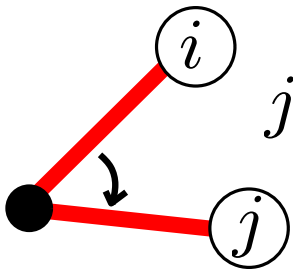
①

remove the map-edges and the marked vertex ①



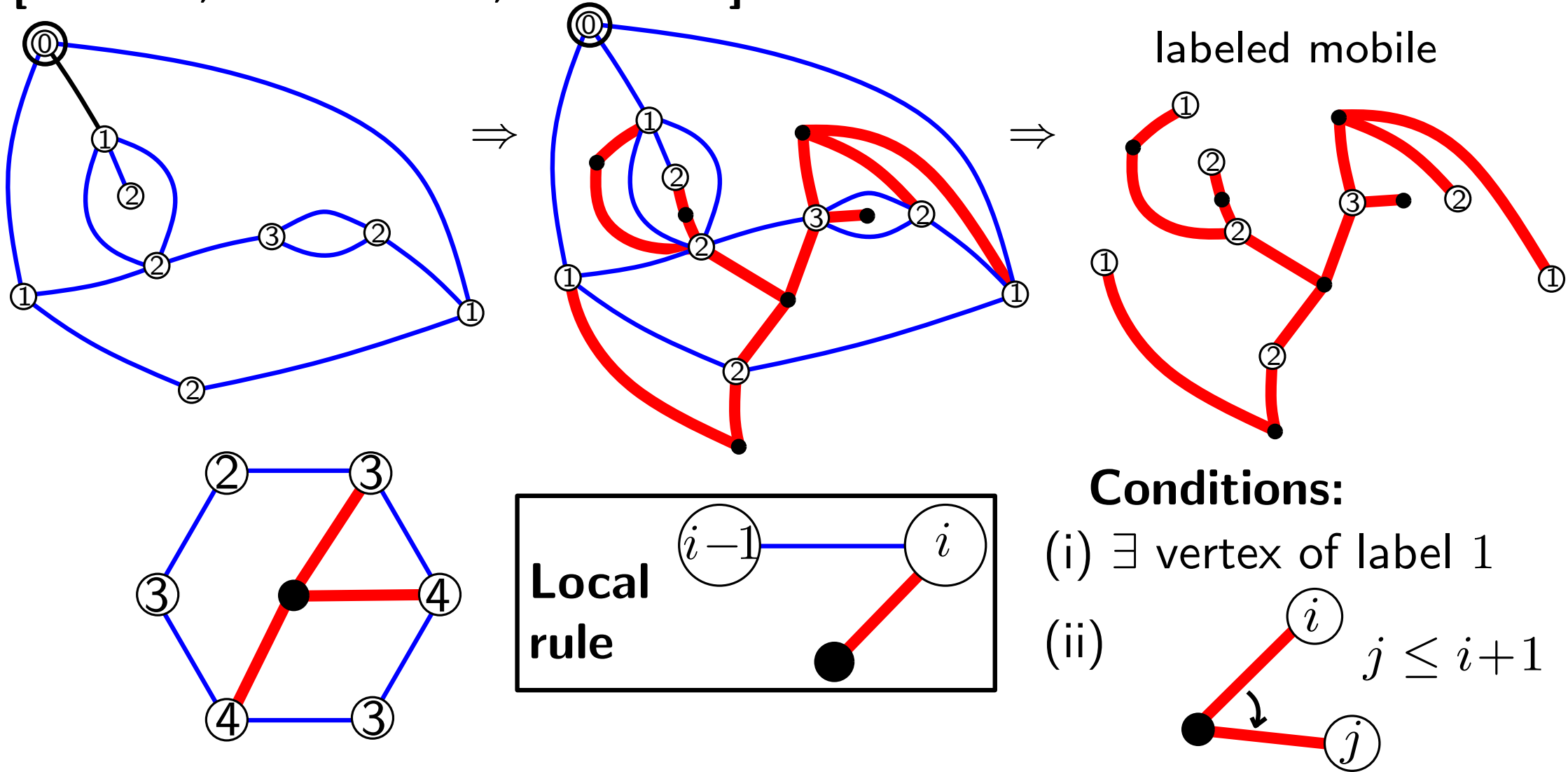
Conditions:

(i) \exists vertex of label 1

(ii)  $j \leq i+1$

The BDG bijection for pointed bipartite maps

[Bouttier, Di Francesco, Guitter'04]

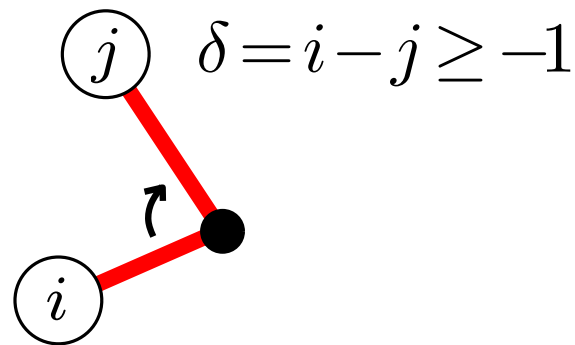
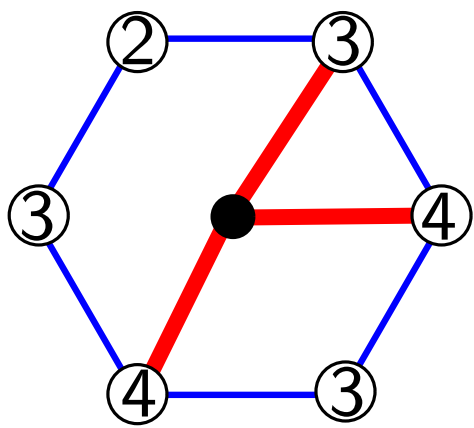
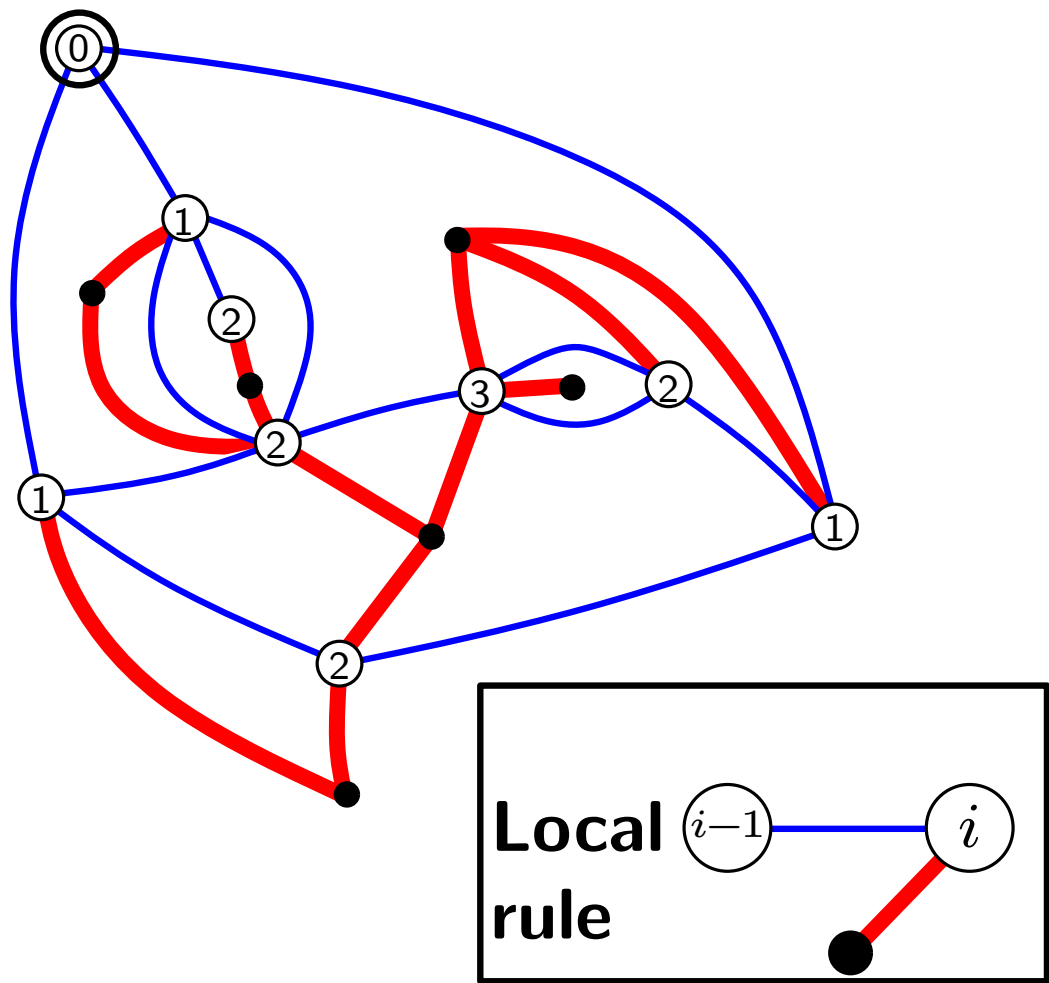


Theorem: The mapping is a **bijection**.

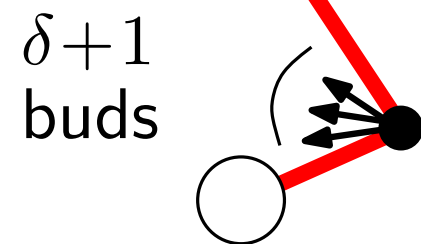
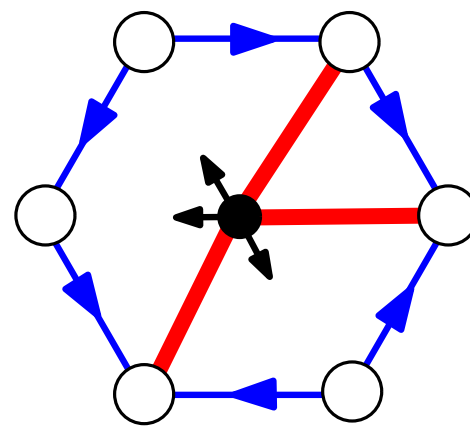
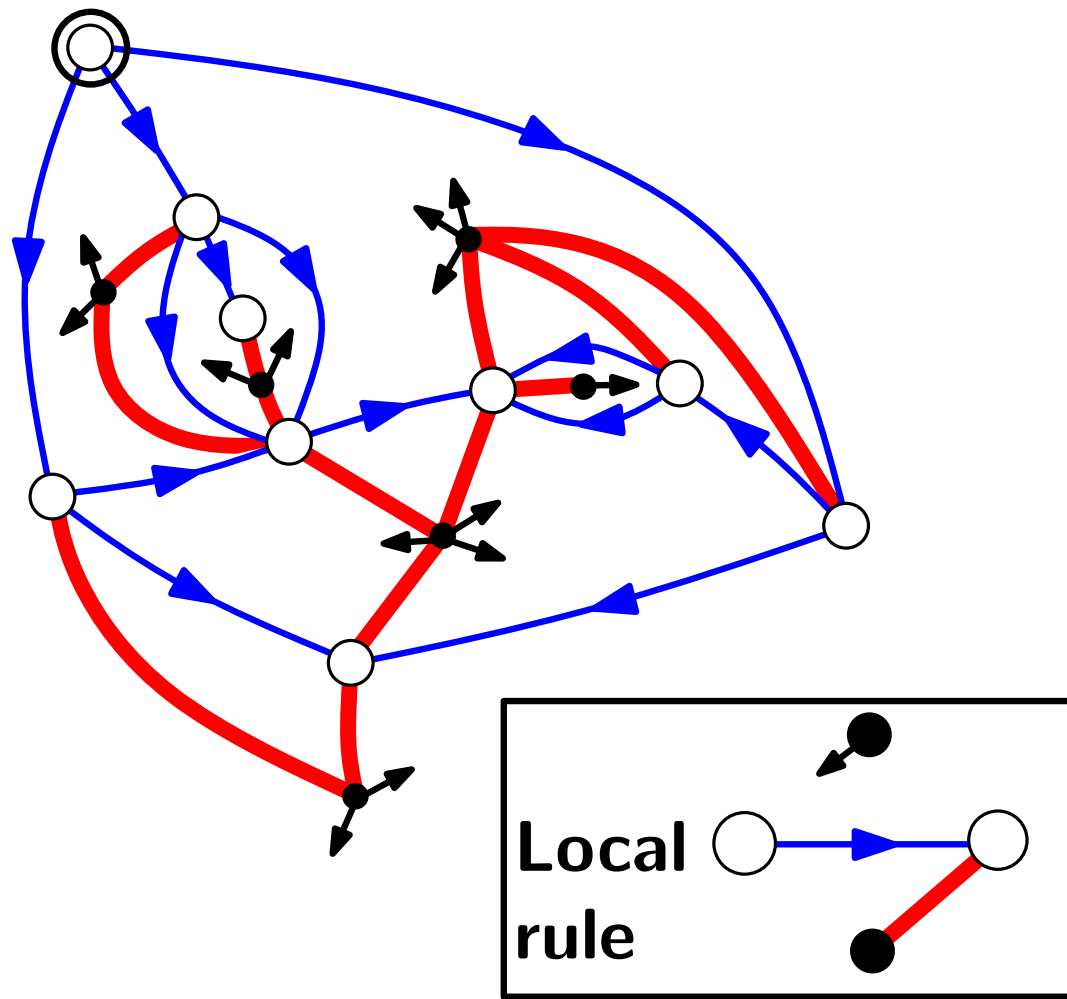
face of degree $2i$ \longleftrightarrow black vertex of degree i

Reformulation with orientations

Distance-labeling



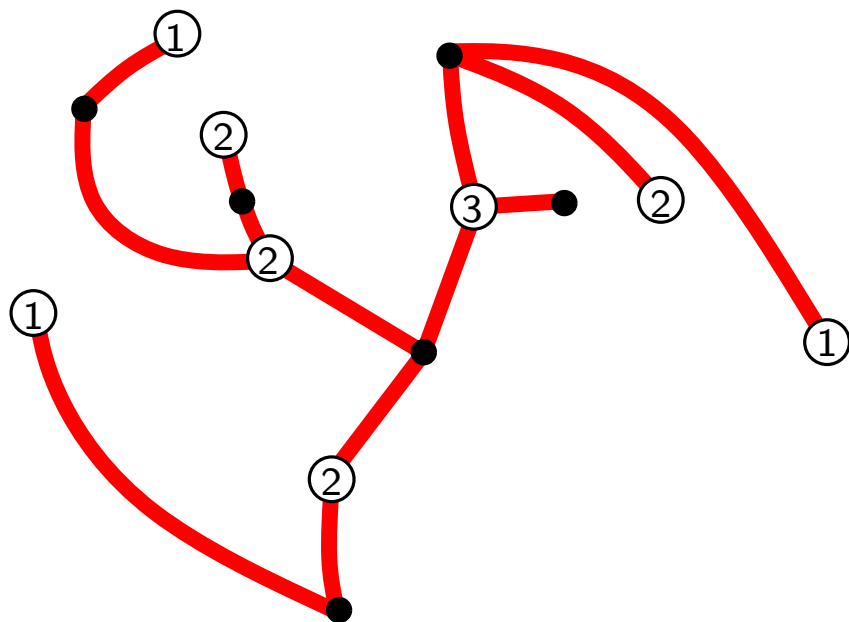
Geodesic orientation



Mobile conditions in the two formulations

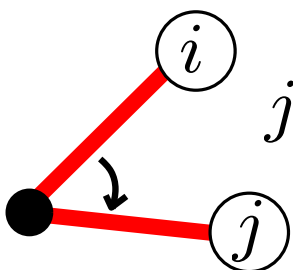
Formulation with labels

gives a labeled mobile



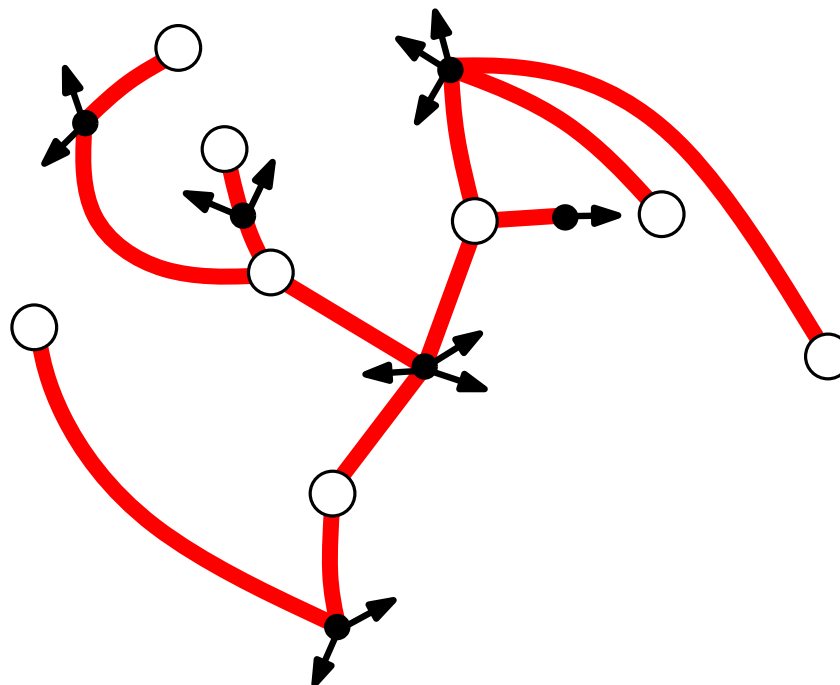
with the conditions:

(i) \exists node of label 1

(ii)  $j \leq i + 1$

Formulation with orientations

gives a "blossoming" mobile



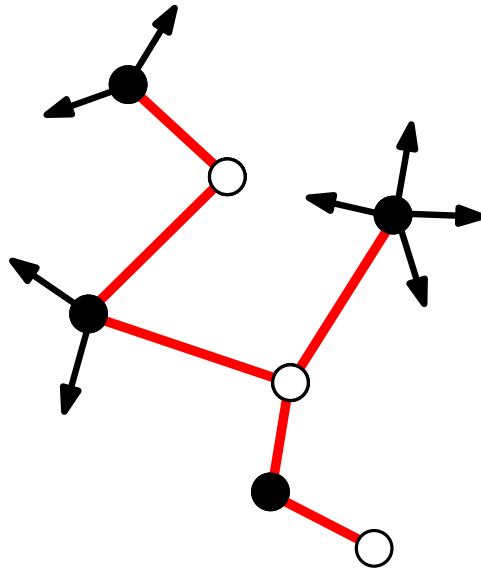
with the condition:

each black vertex has as many buds as neighbors

Definition of blossoming mobiles

- **Blossoming mobile** = bipartite tree (black/white vertices) where each corner at a black vertex carries $i \geq 0$ buds

excess = number of edges - number of buds

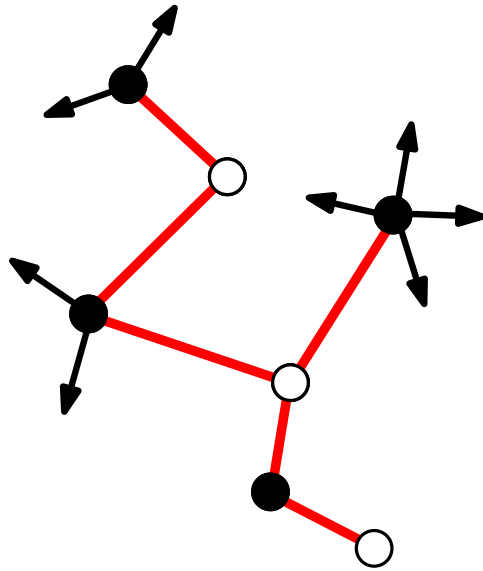


a blossoming mobile of excess -2

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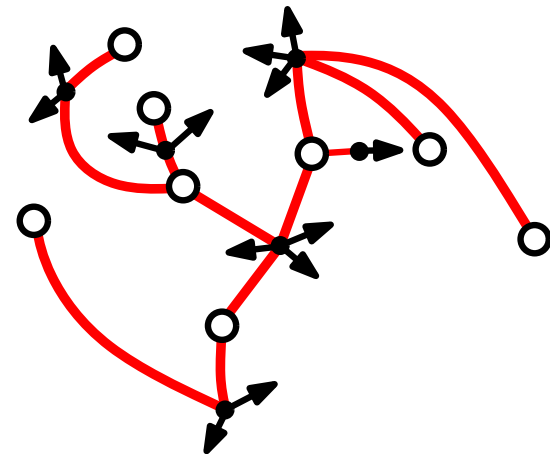
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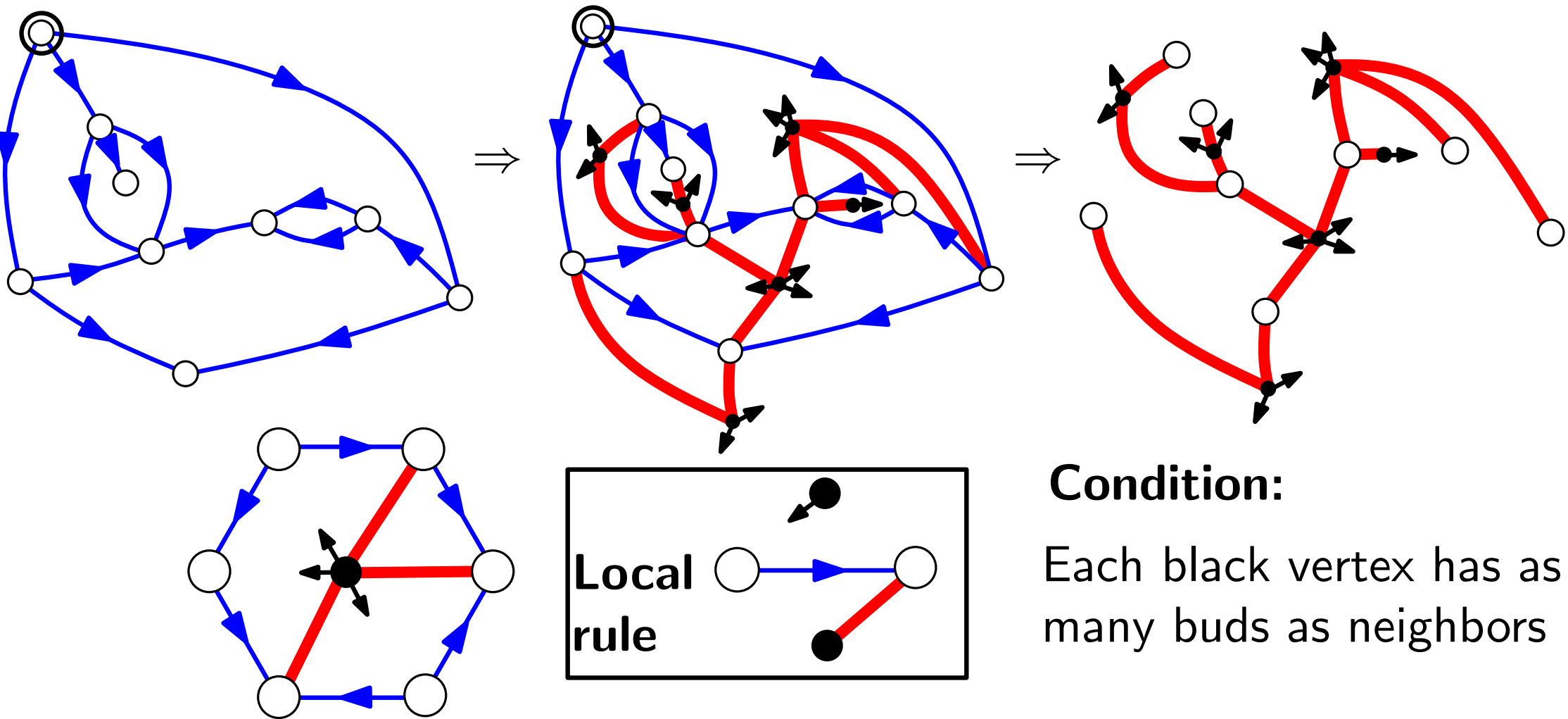
a blossoming mobile of excess -2

- A blossoming mobile is called **balanced** iff each black vertex has as many buds as neighbors

Rk: implies that the excess is 0



Summary of the reformulation



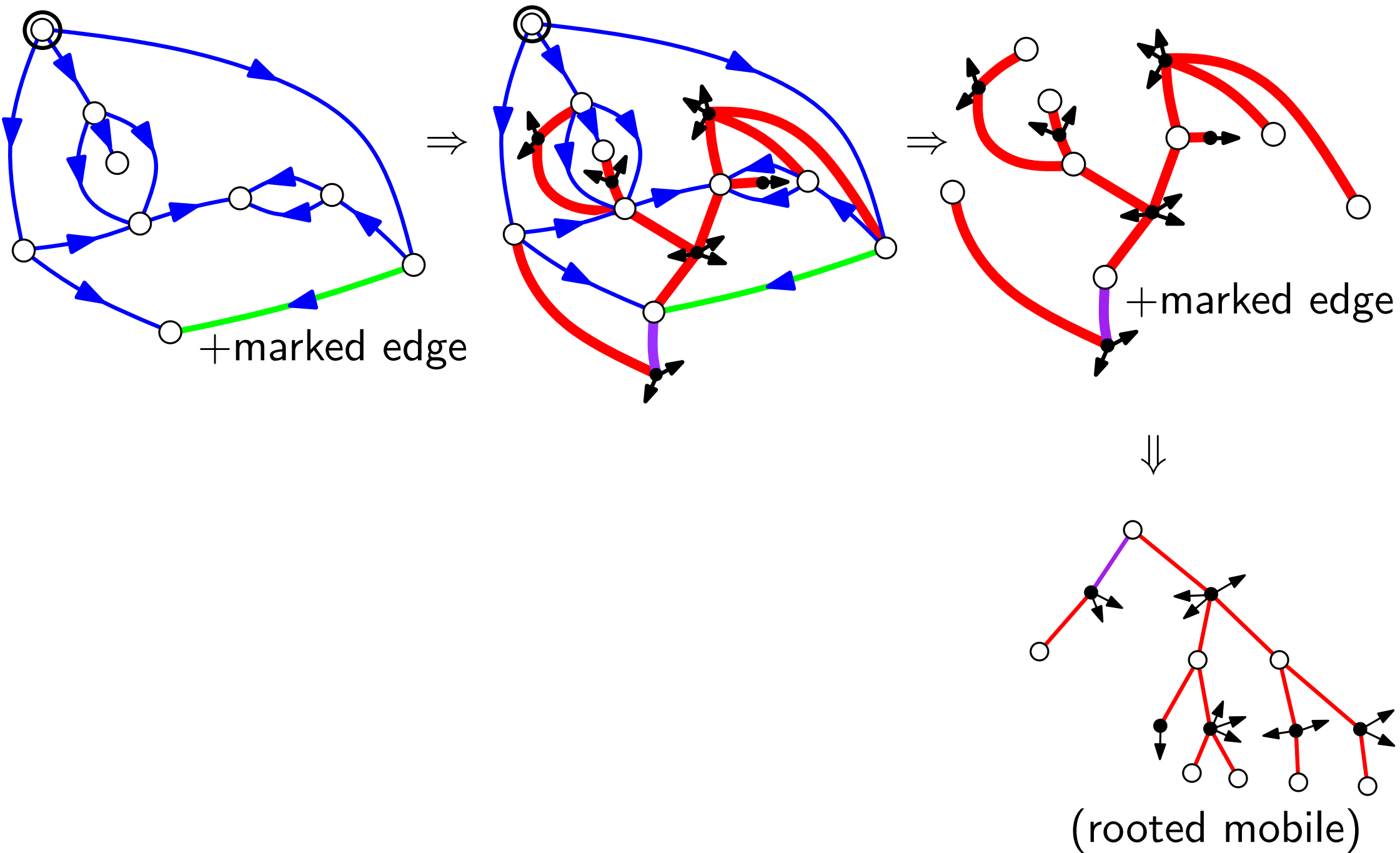
Condition:

Each black vertex has as many buds as neighbors

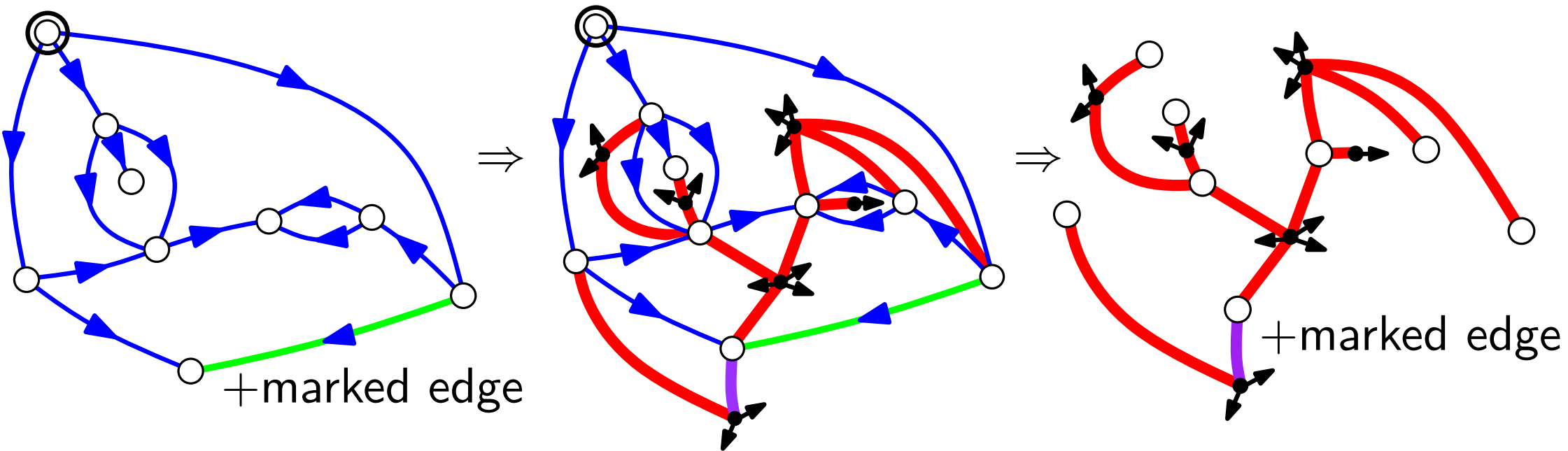
Theoreme: The mapping is a bijection between pointed bipartite maps and balanced blossoming mobiles

face of degree $2i$ \longleftrightarrow black vertex of degree $2i$

Proof of Tutte's slicings formula



Proof of Tutte's slicings formula



Let $B[n_1, n_2, \dots, n_k]$ be the number of rooted bipartite maps with n_i faces of degree $2i$ for $i \in [1..k]$

• **Bijection gives**

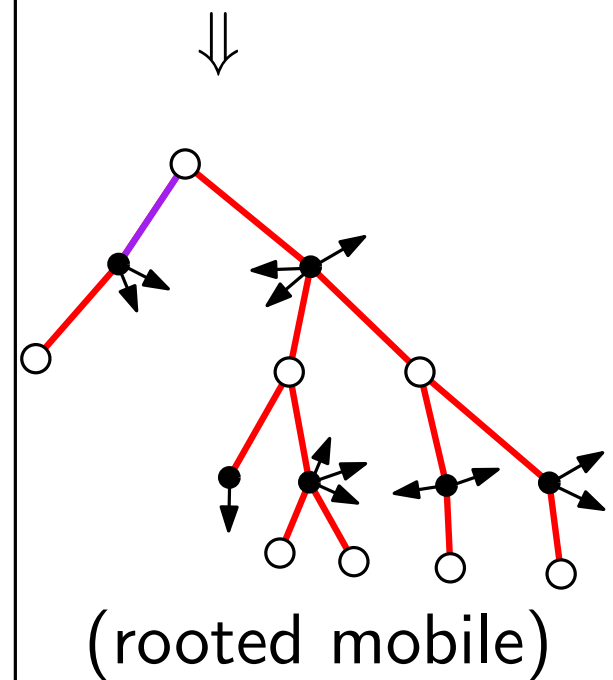
$$v \cdot B[n_1, \dots, n_k] = 2 \cdot \text{coeff } t_1^{n_1} \dots t_k^{n_k} \text{ in } R(t_1, t_2, \dots)$$

where $R \equiv R(t_1, t_2, \dots)$ is the GF of rooted mobiles

given by the equation
$$R = 1 + \sum_{i \geq 1} \binom{2i-1}{i-1} t_i R^i$$

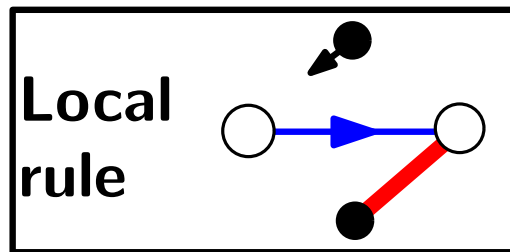
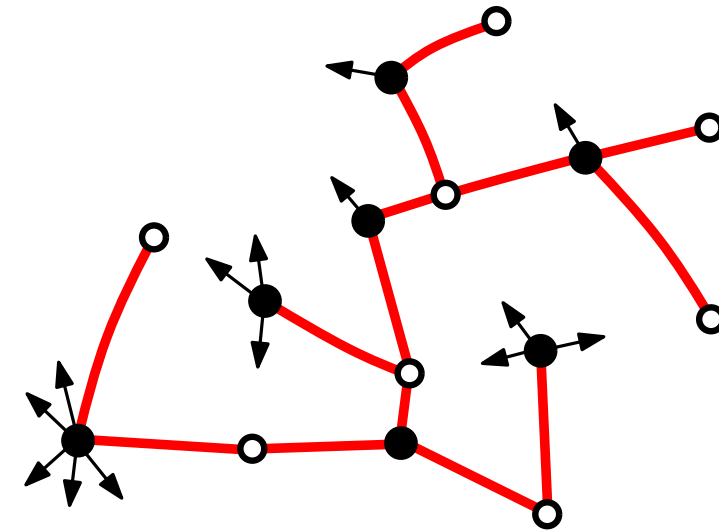
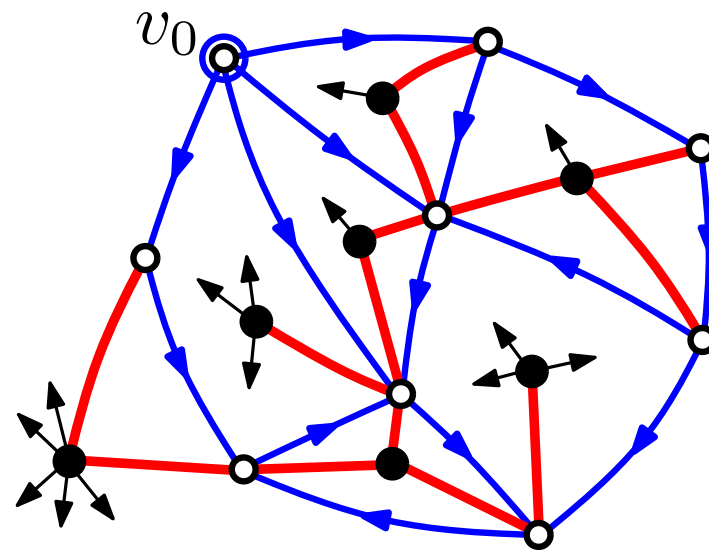
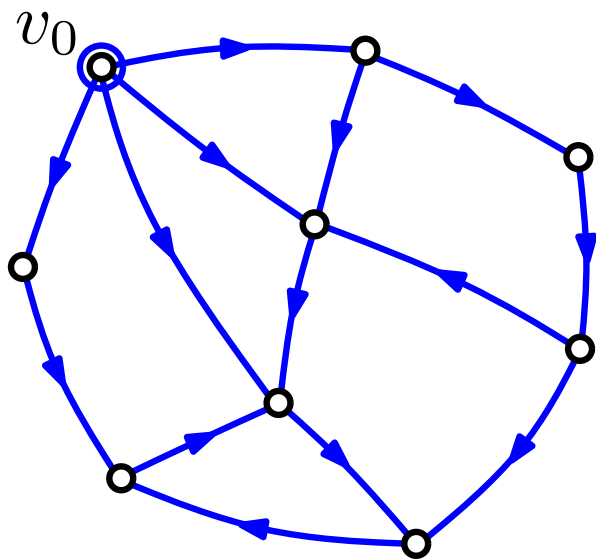
• **Lagrange inversion formula** gives:

$$[t_1^{n_1} \dots t_k^{n_k}] R = \frac{e!}{(v-1)!} \prod_{i=1}^k \frac{1}{n_i!} \binom{2i-1}{i-1}^{n_i}$$



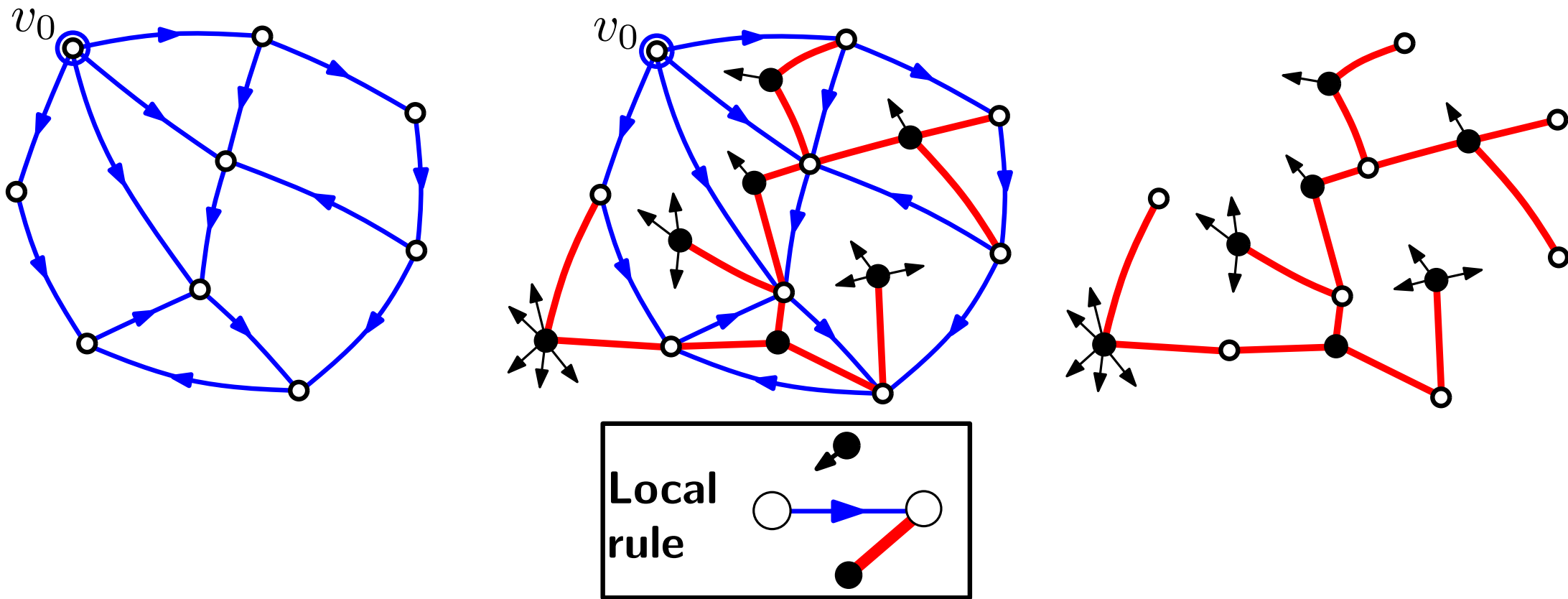
Extension for pointed orientations with no ccw cycle

- More generally, we **obtain a blossoming mobile** (of excess 0) if we start from a vertex-pointed orientation such that :
 - the marked vertex v_0 is a **“source”** (no incoming edge)
 - every vertex is **accessible** from v_0 by a directed path
 - **there is no ccw cycle** (with $v_0 \in$ outer face)



Extension for pointed orientations with no ccw cycle

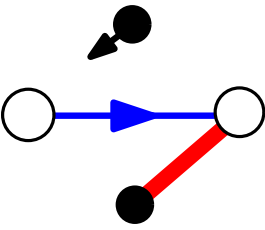
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Theorem : Let \mathcal{O}_0 be this family of orientations, then the correspondence is a bijection with mobiles of excess 0

Proof that it gives a tree

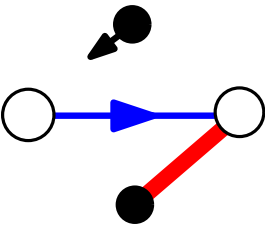
Start from an oriented map $M \in \mathcal{O}_0$ and apply the local rule



Let G be the graph of red edges and their incident vertices

Proof that it gives a tree

Start from an oriented map $M \in \mathcal{O}_0$ and apply the local rule

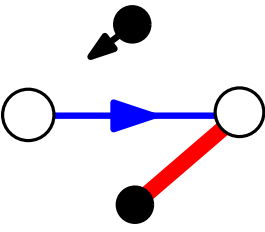


Let G be the graph of red edges and their incident vertices

G has $|V_M| - 1$, white vertices, $|F_M|$ black vertices, et $|E_M|$ edges

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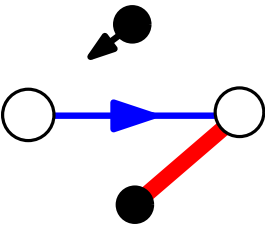
Euler relation: $|E_M| = |V_M| + |F_M| - 2$

$\Rightarrow G$ has one more vertices than edges

hence G is a tree iff G is acyclic

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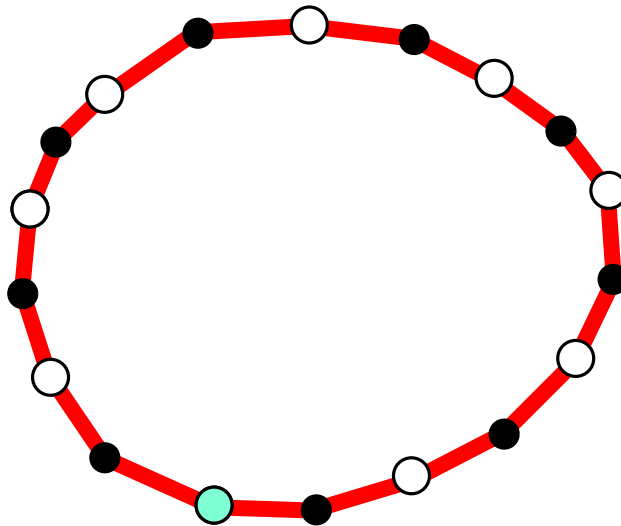
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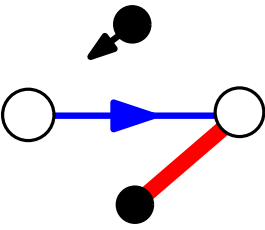
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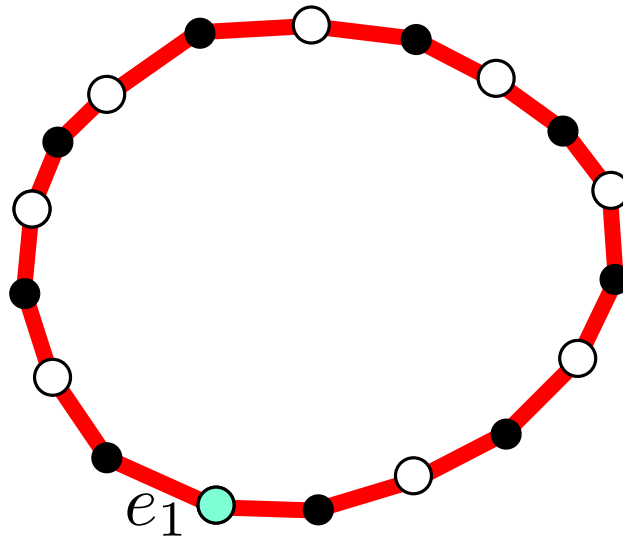
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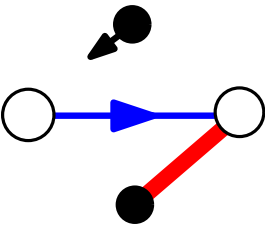
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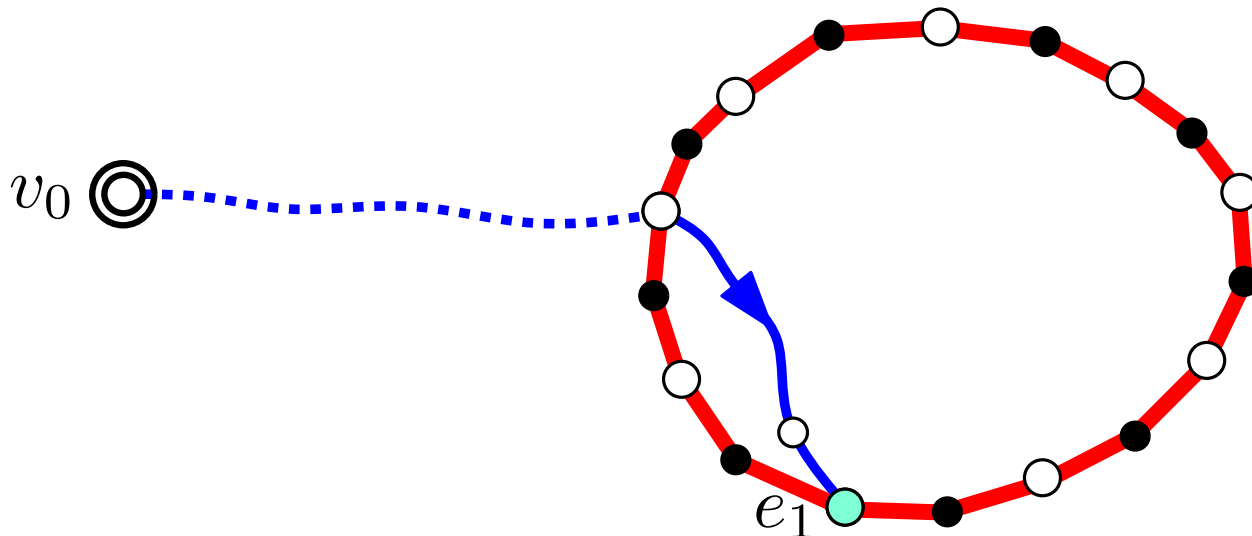
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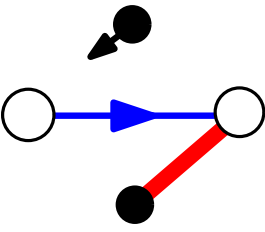
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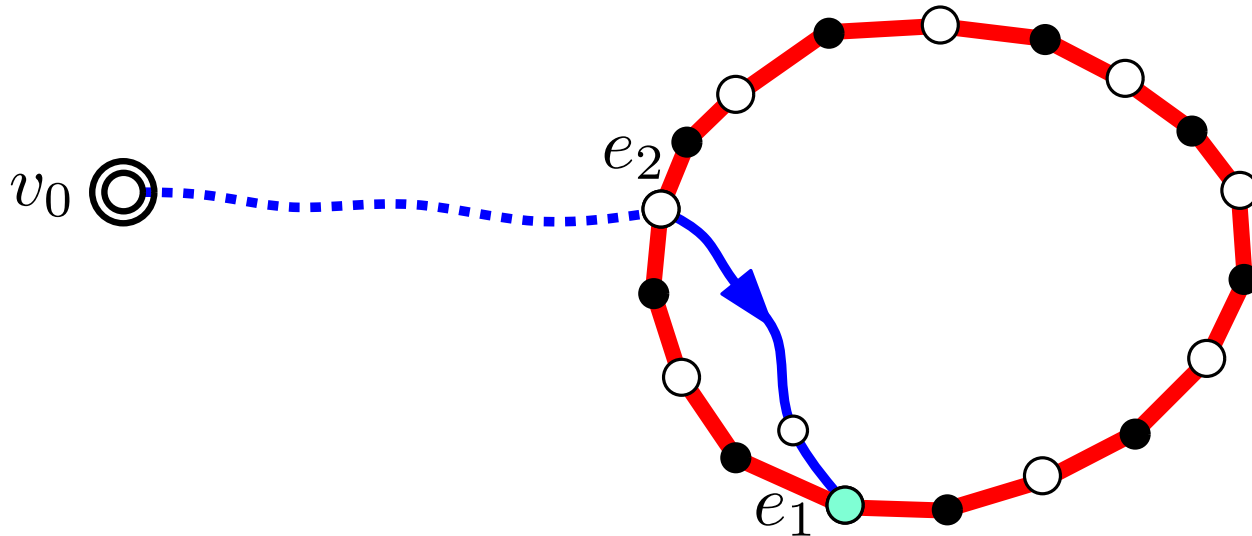
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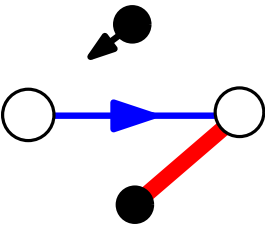
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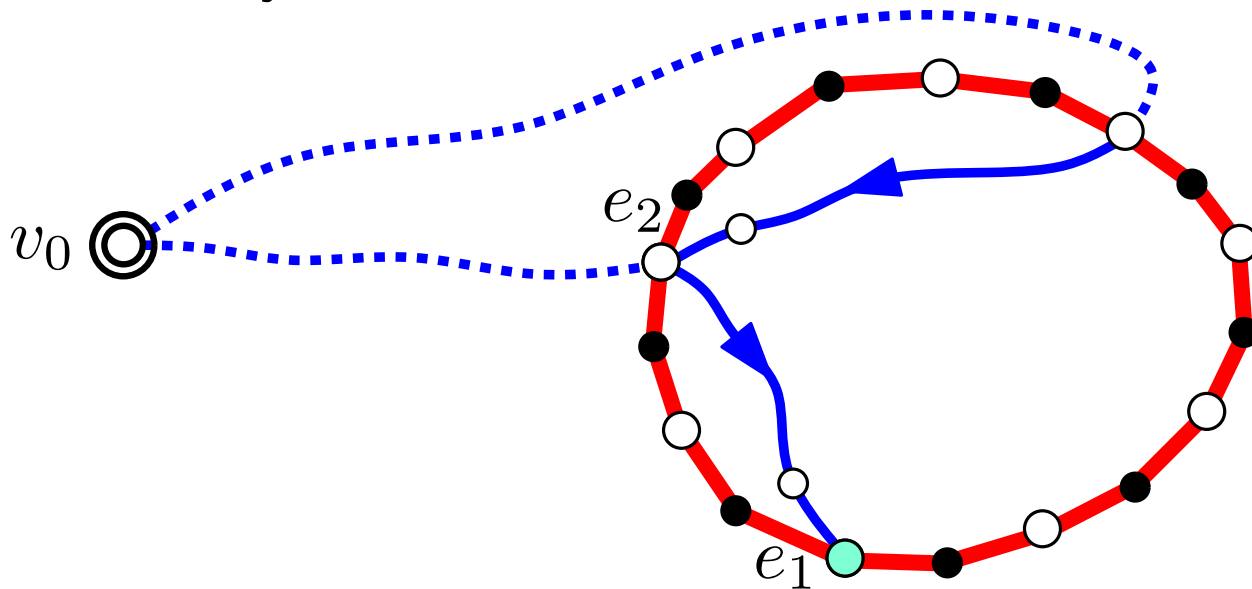
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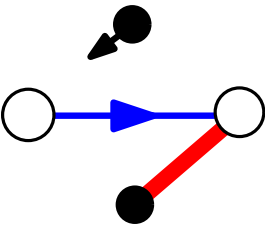
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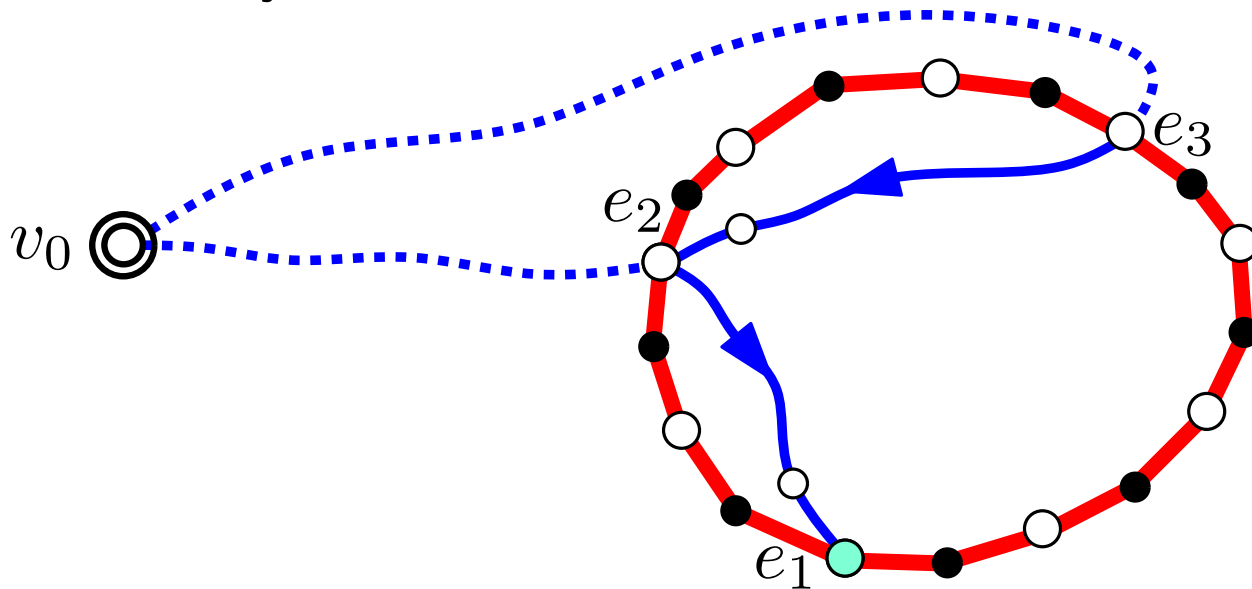
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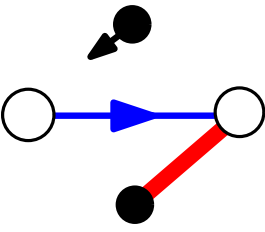
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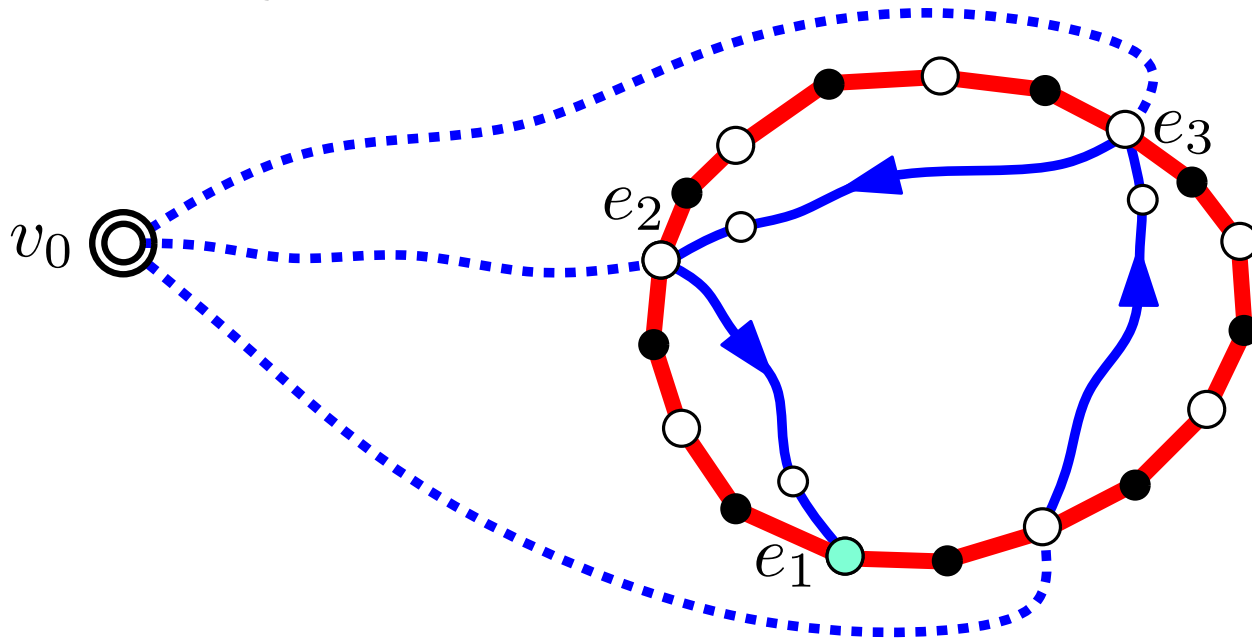
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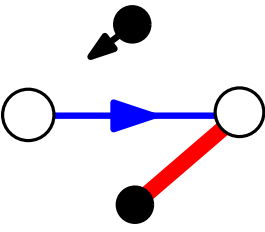
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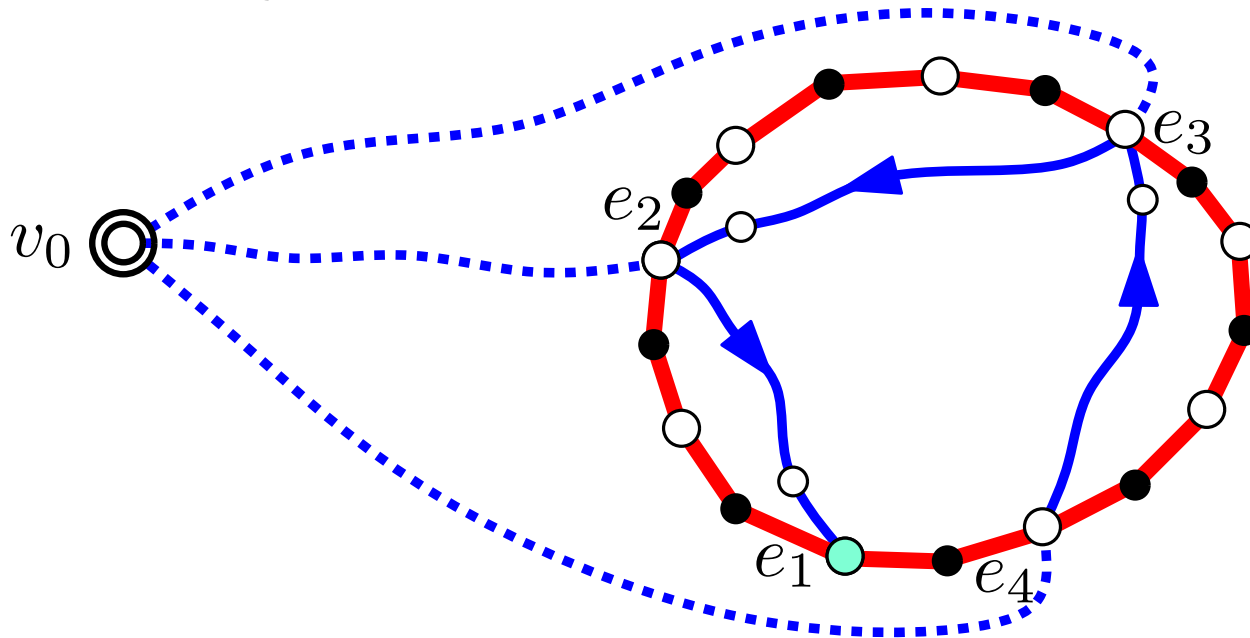
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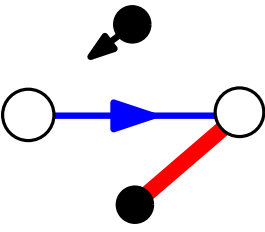
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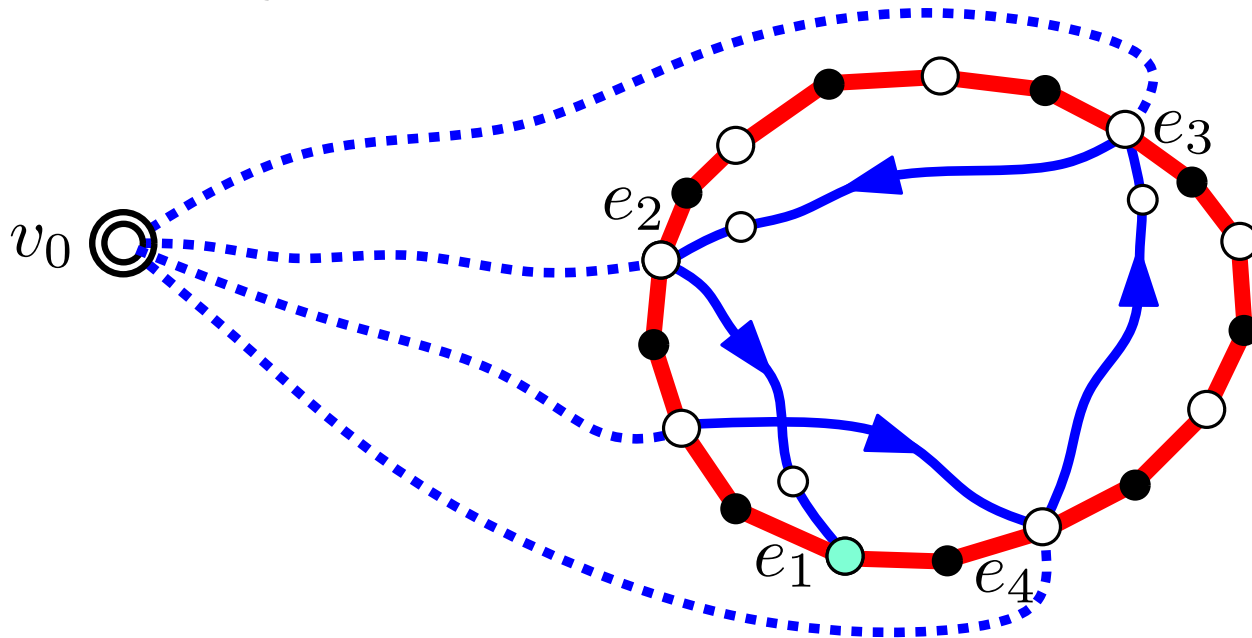
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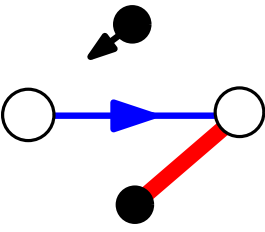
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Start from an oriented map $M \in \mathcal{O}_0$ and apply the local rule



Let G be the graph of red edges and their incident vertices

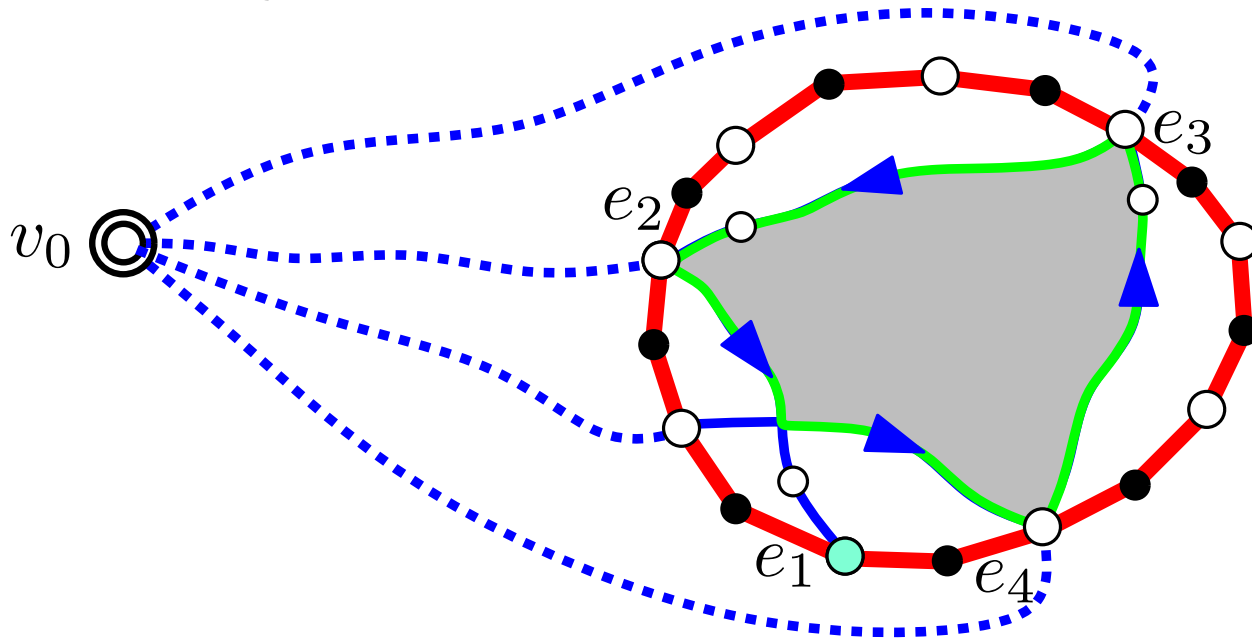
G has $|V_M| - 1$, white vertices, $|F_M|$ black vertices, et $|E_M|$ edges

Euler relation: $|E_M| = |V_M| + |F_M| - 2$

$\Rightarrow G$ has one more vertices than edges

hence G is a tree iff G is acyclic

Assume G has a cycle :



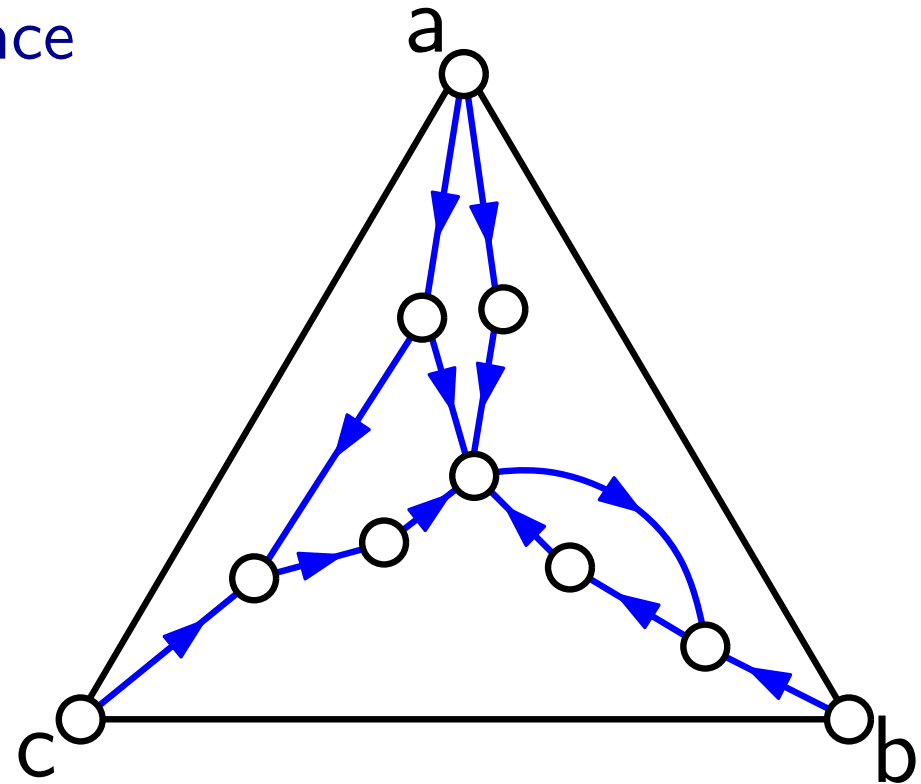
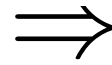
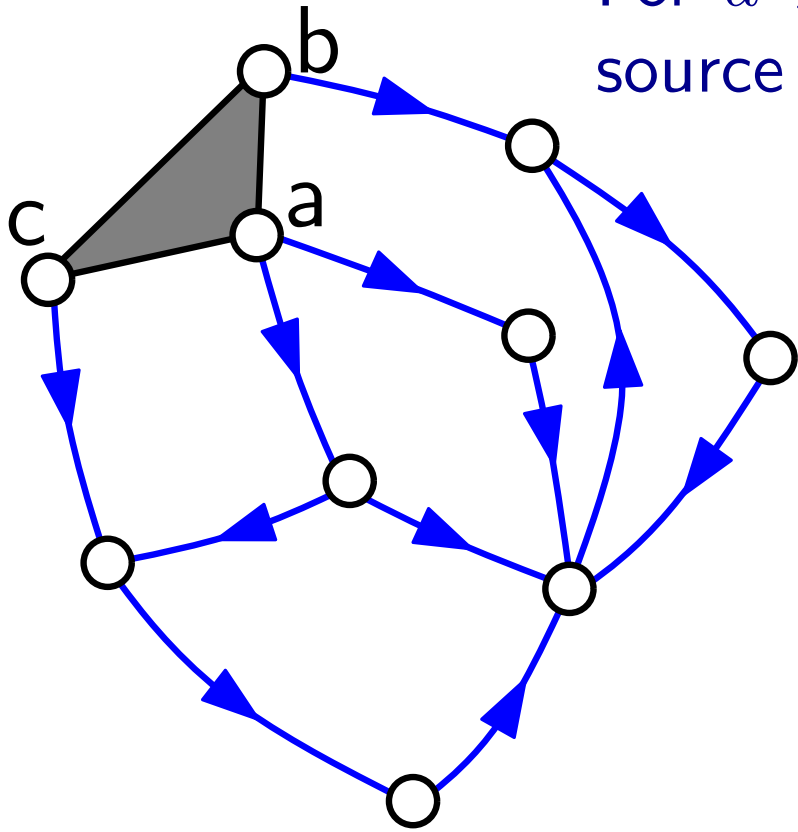
prisoner ccw cycle
 \Rightarrow contradiction

Extension for mobiles of negative excess

More generally the “source” can be a d -gon, for any $d \geq 0$

Example for $d = 3$

For $d > 0$, we take the d -gonal source as the outer face

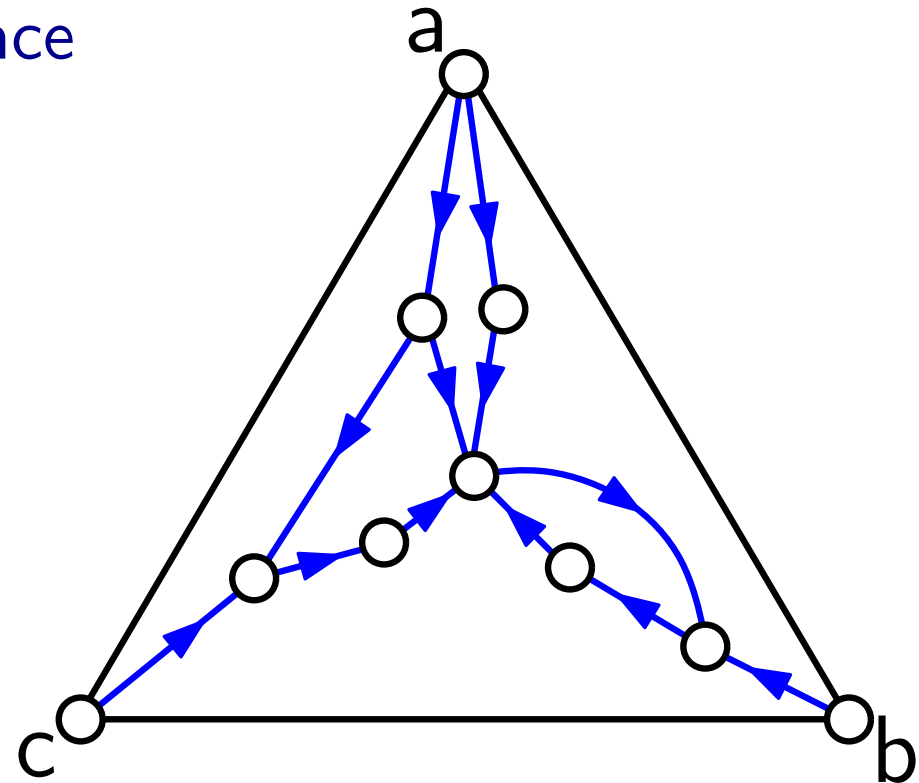
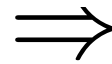
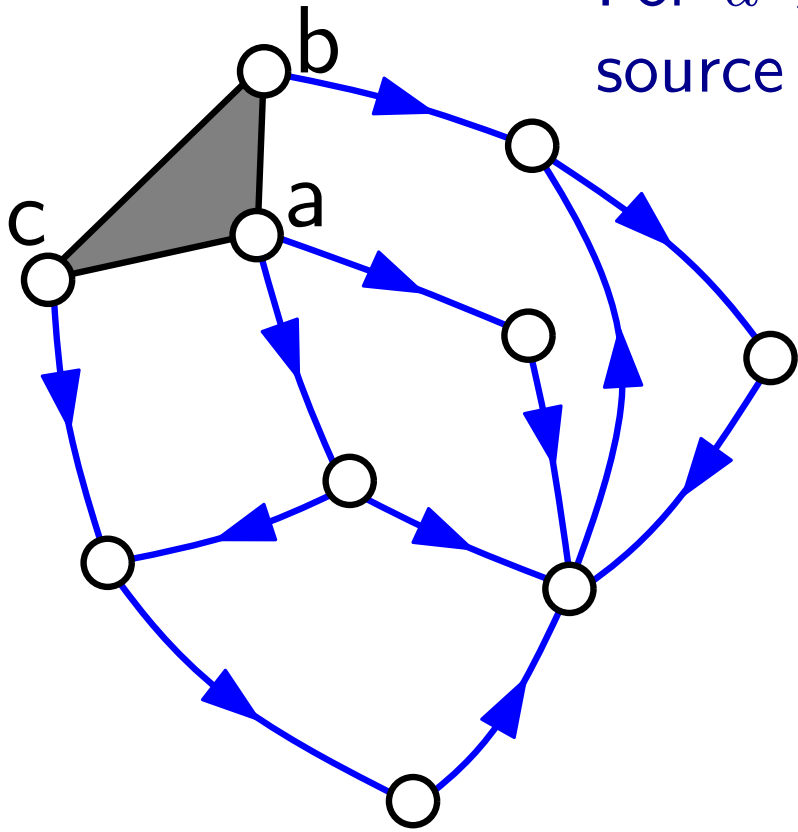


Extension for mobiles of negative excess

More generally the “source” can be a d -gon, for any $d \geq 0$

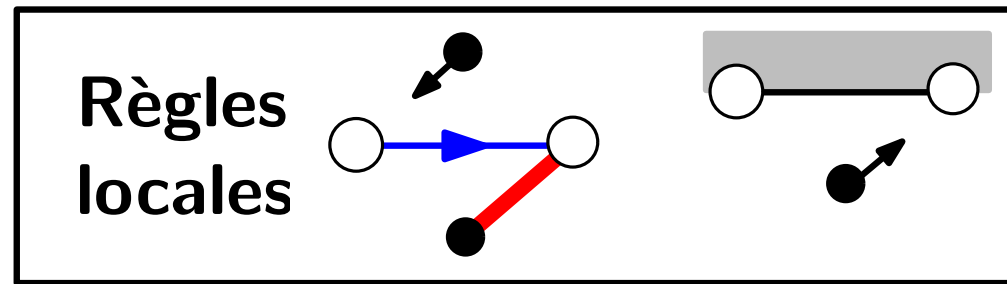
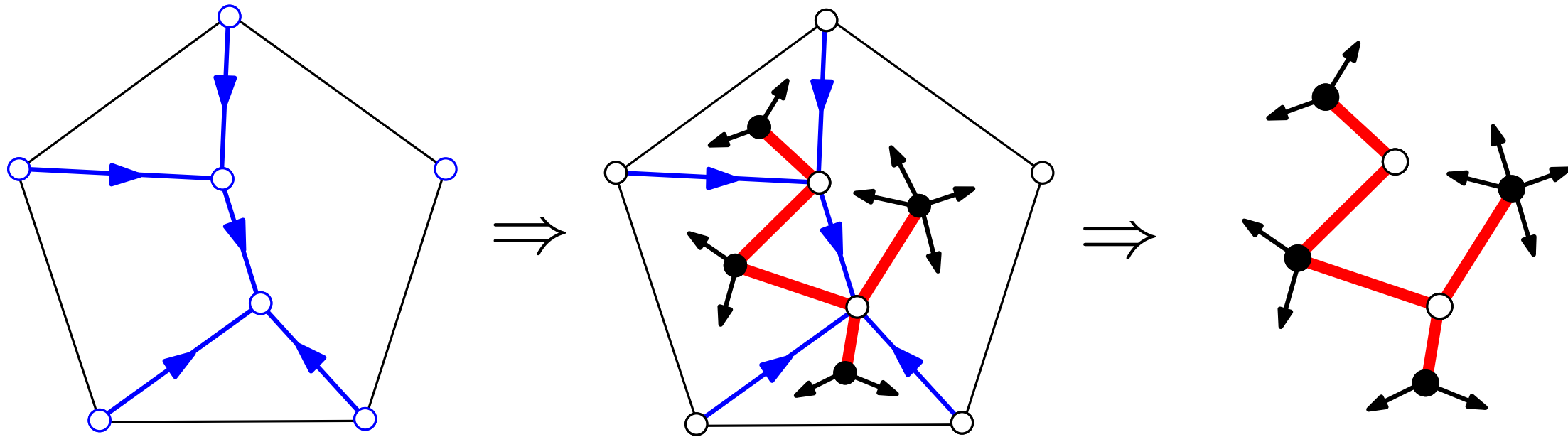
Example for $d = 3$

For $d > 0$, we take the d -gonal source as the outer face



- Let \mathcal{O}_{-d} be the family of these orientations, still with the conditions
- the d -gonal **source** has no ingoing edge
 - **accessibility** of every vertex from the source
 - **no ccw cycle**

Extension for mobiles of negative excess

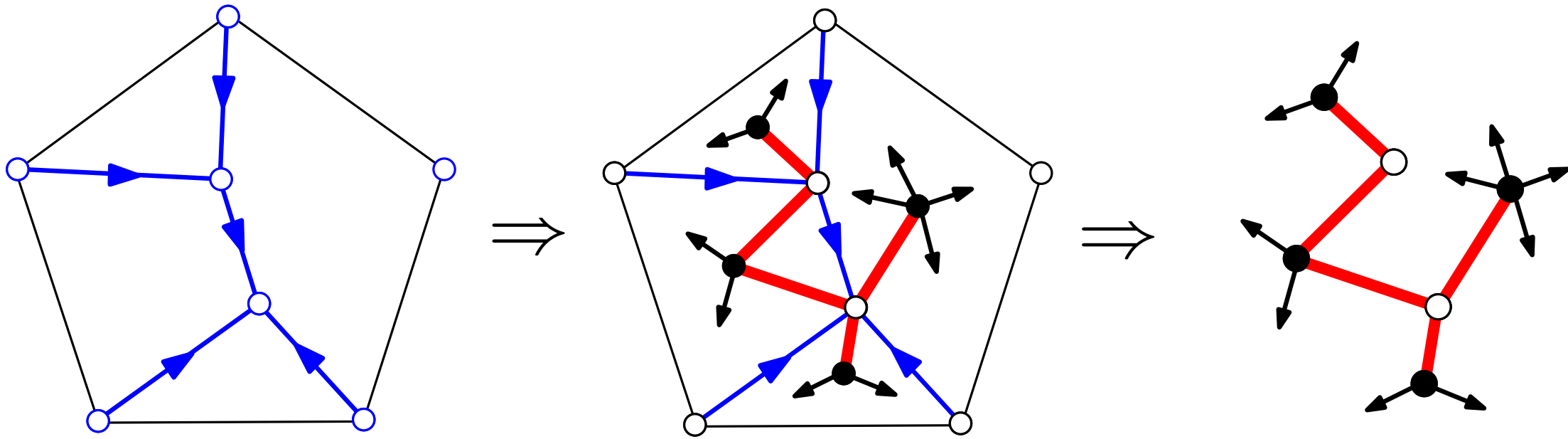


Theorem [Bernardi-F'10]: For $\delta \leq 0$, the correspondence Φ is a **bijection** between \mathcal{O}_δ and mobiles of excess δ .

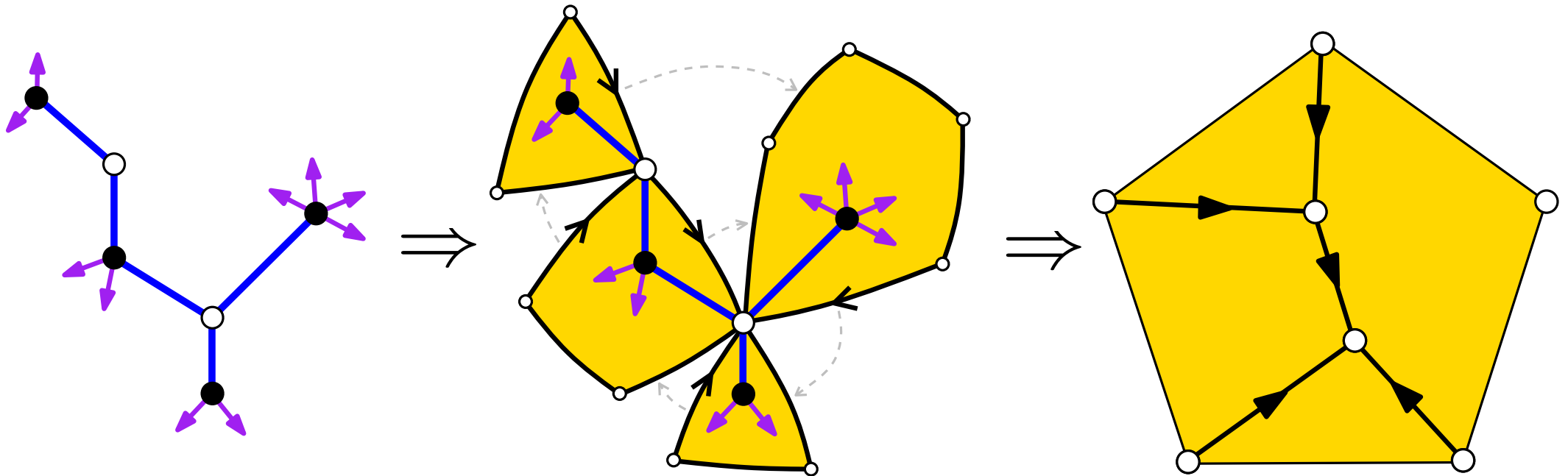
degrees of the inner faces \longleftrightarrow degrees of the black vertices
 indegrees of internal vertices \longleftrightarrow degrees of white vertices

cf [Bernardi'07], [Bernardi-Chapuy'10]

Extension for mobiles of negative excess



- Inverse mapping (tree \rightarrow cactus \rightarrow closure operations)



Specializing the correspondence

The correspondence Φ is a bijection between the family $\mathcal{O} = \bigcup_{d \geq 0} \mathcal{O}_{-d}$ of oriented maps and mobiles of nonpositive excess

Idea: Let \mathcal{F} be the family of planar maps we consider
(e.g. bipartite maps, simple triangulations, etc.)

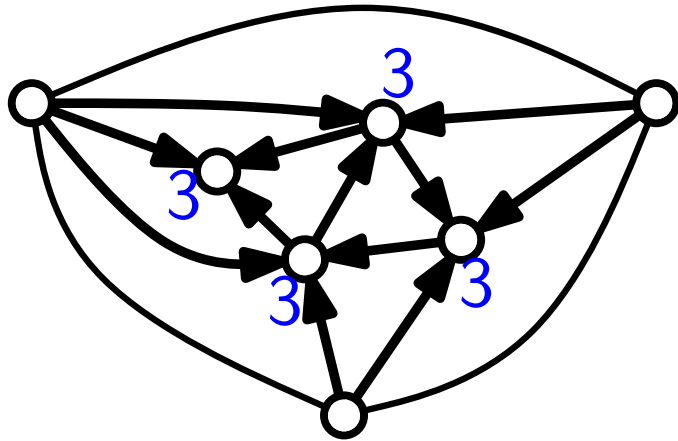
Prove that a map is in \mathcal{F} iff it admits a **canonical orientation** in \mathcal{O} specified by face-degrees and vertex-indegrees conditions

Specialize Φ to the corresponding subfamily $\mathcal{O}_{\mathcal{F}} \subseteq \mathcal{O}$

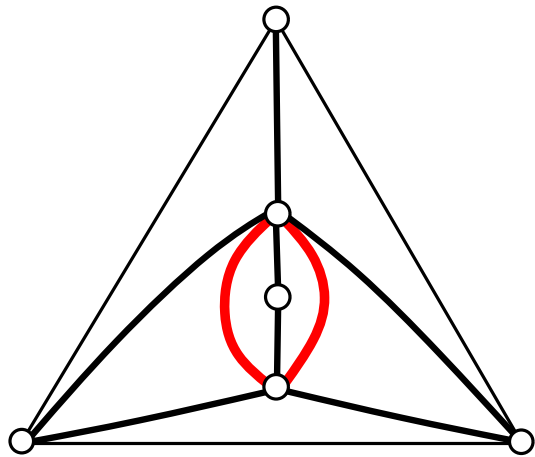
Gives a bijection between \mathcal{F} and a well characterized family of mobiles

Application to simple triangulations

For a triangulation T , a **3-orientation** of T is an orientation of the inner edges of T such that every inner vertex has **indegree 3**



Rk: If a triangulation T admits a 3-orientation, then T is simple



Assume there is a 2-cycle C

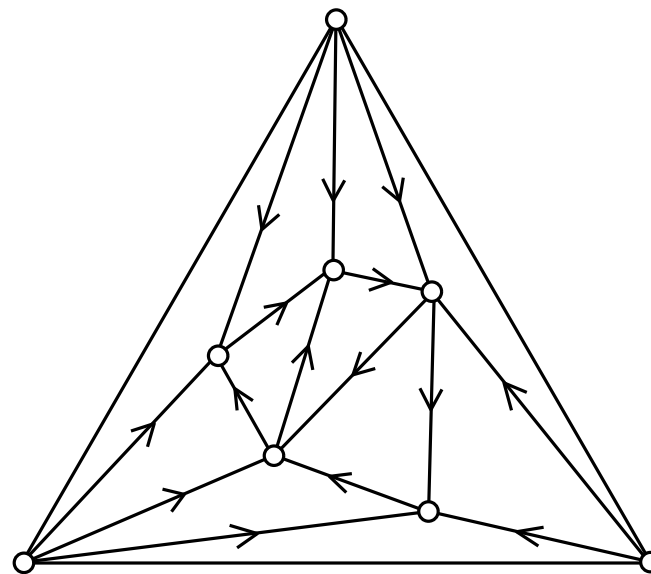
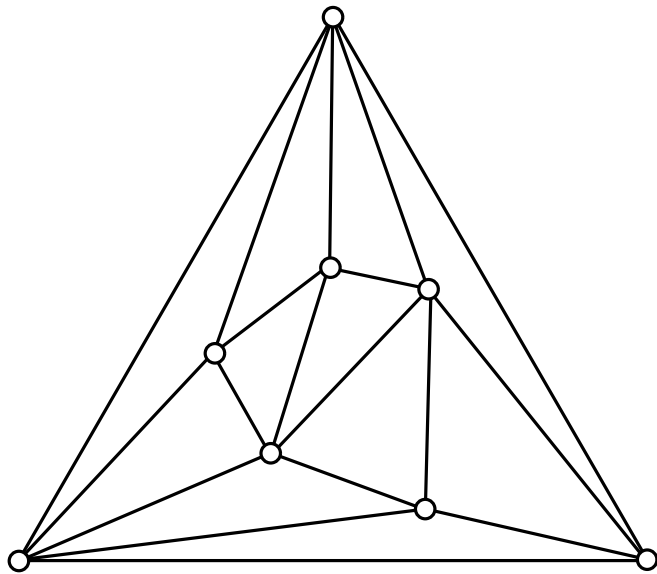
If there are k vertices inside C then there are $3k - 1$ edges inside C
 \Rightarrow total indegree is too large compared to the number of edges

Existence of a canonical 3-orientation

Theorem (Schnyder'89): Any simple triangulation admits a 3-orientation

Theorem: Let T be a simple triangulation. Then T has a unique 3-orientation with no ccw cycle, the **minimal 3-orientation** (set of 3-orientations is a lattice, flip = reverse cw to ccw)

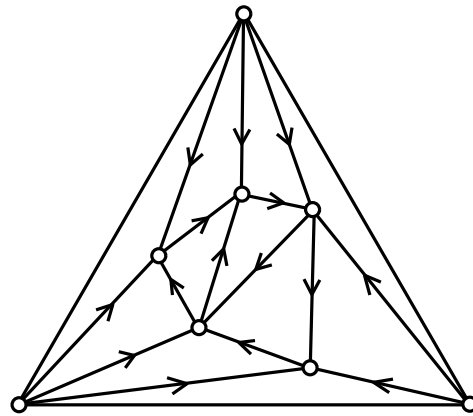
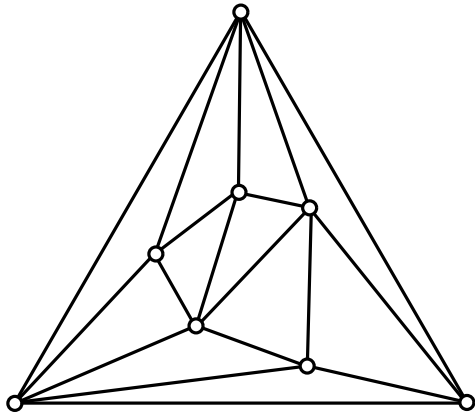
[Ossoana de Mendez'94], [Brehm'03], [[Felsner'03]]



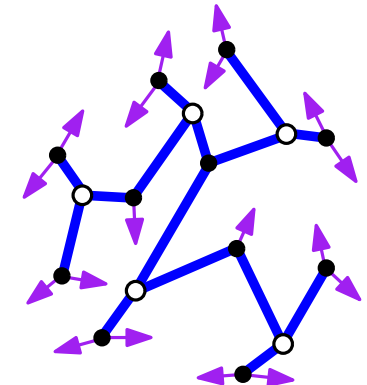
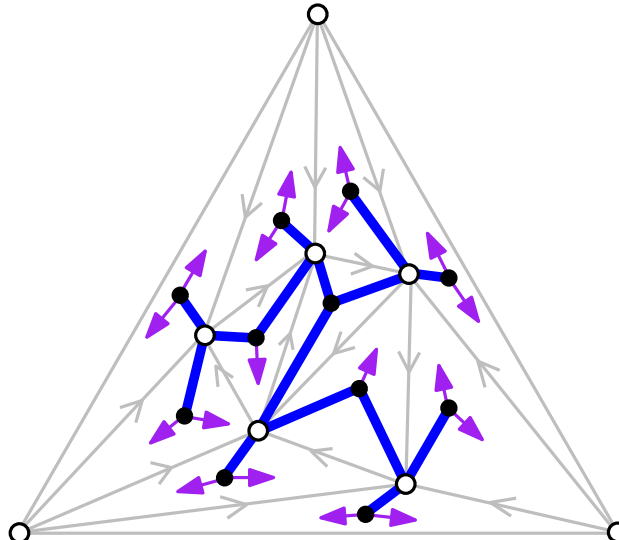
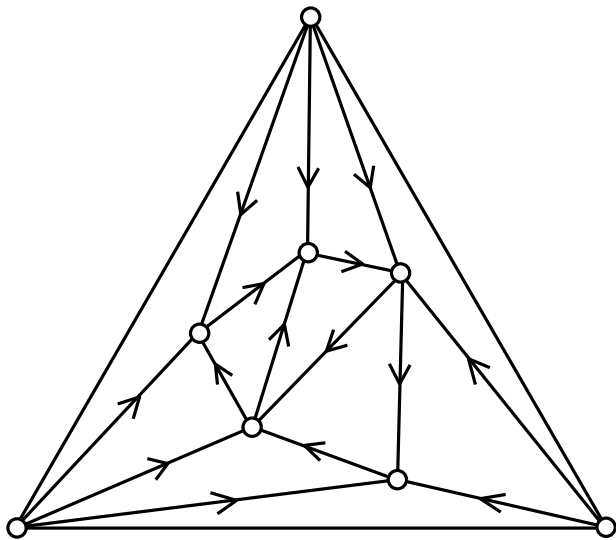
Bijection for simple triangulations

- From the lattice property (**taking the min**) we have family \mathcal{T} of simple triangulations \leftrightarrow subfamily $\mathcal{O}_{\mathcal{T}}$ of \mathcal{O} where:

- faces have degree 3
- inner vertices have indegree 3



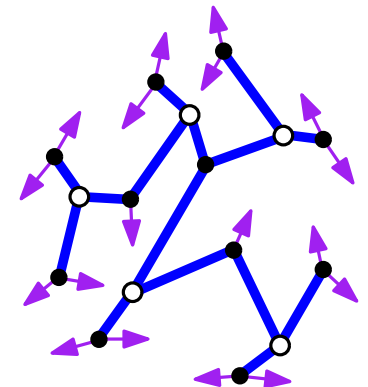
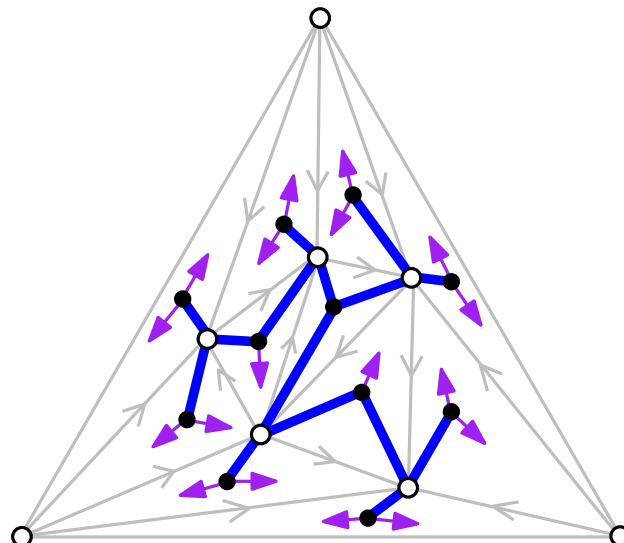
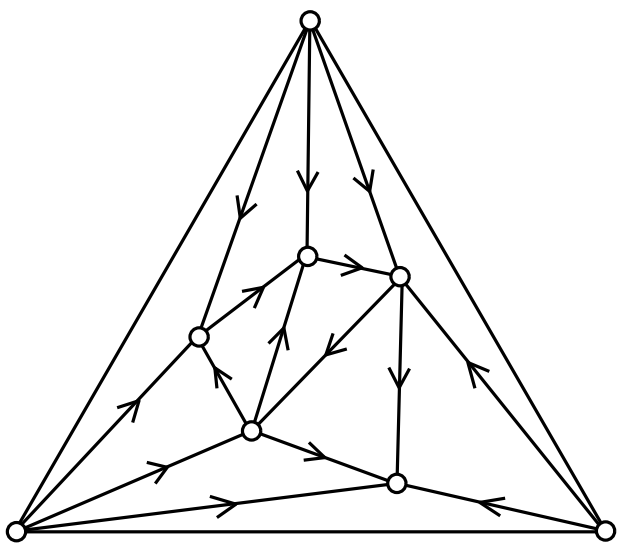
- From the **bijection Φ specialized to \mathcal{F}** , we have $\mathcal{F} \leftrightarrow$ **mobiles** where all vertices have **degree 3**



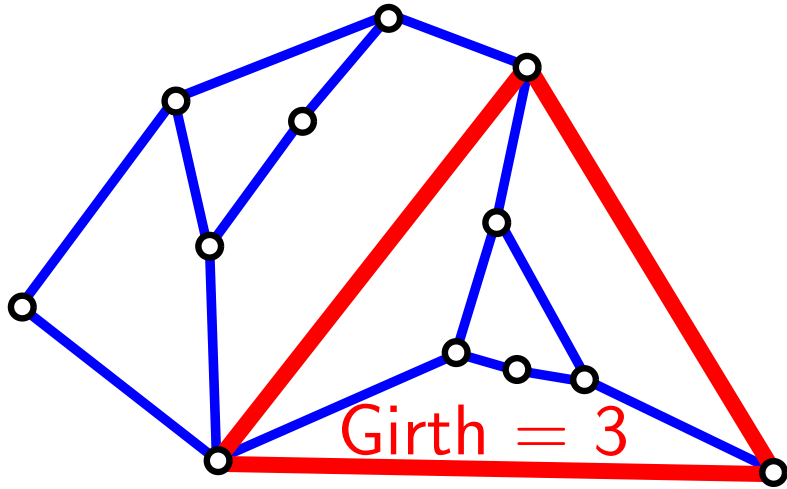
Counting simple triangulations

Counting: The generating function of mobiles with vertices of degree 3 rooted on a white corner is $T(x) = U(x)^3$, where $U(x) = 1 + xU(x)^4$.

Consequently, the number of (rooted) simple triangulations with $2n$ faces is $\frac{1}{n(2n-1)} \binom{4n-2}{n-1}$.



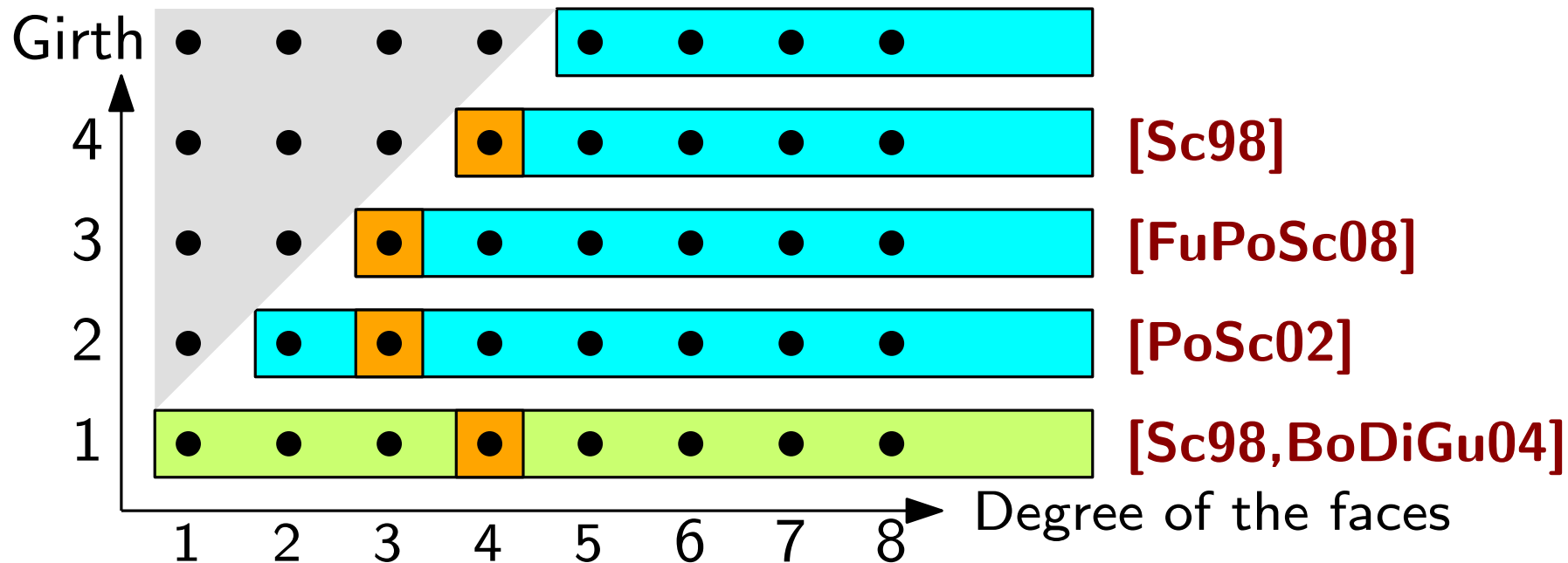
Extension to any girth and face-degrees



girth=length shortest cycle

Rk: $\text{girth} \leq \text{minimal face-degree}$

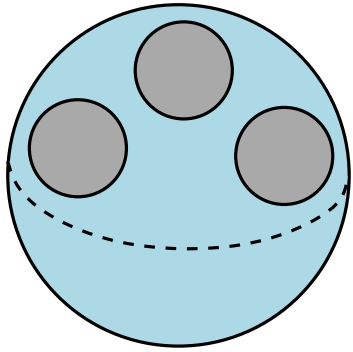
Our approach works in any girth d , with control on the face-degrees



Other approach using slice decompositions **[Bouttier, Guitter'15]**

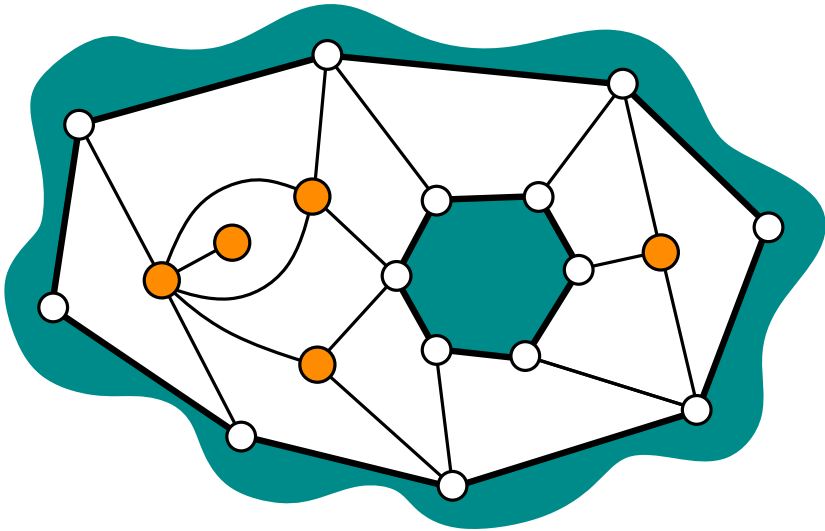
Maps with boundaries

- Sphere with k holes = sphere where k disks have been removed



sphere with 3 holes

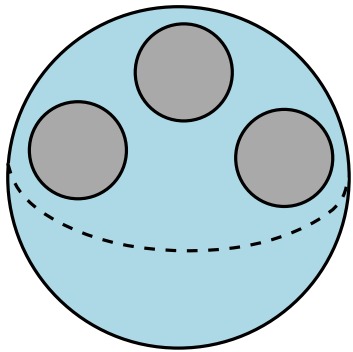
- Map with k boundaries = graph embedded on the sphere with k holes
the boundaries are occupied by cycles of edges



A quadrangulations with 2 boundaries
of lengths 8 and 6, and 5 internal vertices

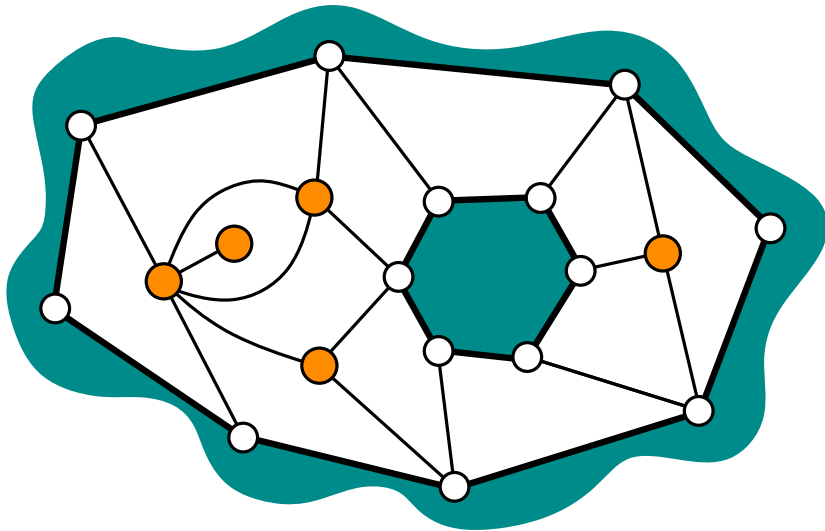
Maps with boundaries

- Sphere with k holes = sphere where k disks have been removed



sphere with 3 holes

- Map with k boundaries = graph embedded on the sphere with k holes
the boundaries are occupied by cycles of edges



A quadrangulations with 2 boundaries
of lengths 8 and 6, and 5 internal vertices

(also = planar map with k distinguished faces whose contours
are vertex-disjoint simple cycles)

Counting triangulations with boundaries

b boundaries of lengths k_1, \dots, k_b

n internal vertices

- Without loops and multiple edges, formula only for $b = 1$

No loops (girth=2)

$$t_n^{(k)} = \frac{2^{n+1}(2k-3)! (3n+2k-3)!}{(k-2)!^2 n!(2n+2k-2)!}$$

[Mullin'65] (recursive method)

bijective proof in [Poulalhon, Schaeffer'02]

No loops & multiple edges (girth=3)

$$s_n^{(k)} = \frac{2(2k-3)! (4n+2k-5)!}{(k-1)!(k-3)! n!(3n+2k-3)!}$$

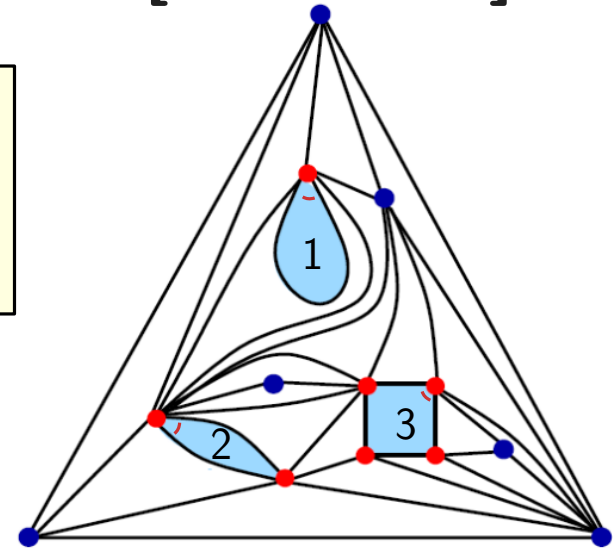
[Brown'64] (recursive method)

bijective proofs in [Poulalhon, Schaeffer'06]
[Bernardi, F'10]

- With loops and multiple edges, nice factorized formula [Krikun'07]

$$a_n^{(k_1, \dots, k_b)} = \frac{4^{n-1} (2k + 3n - 5)!!}{(n - b + 1)! (2k + n - 1)!!} \prod_{j=1}^b k_j \binom{2k_j}{k_j}$$

bijective proof in [Bernardi, F'15]



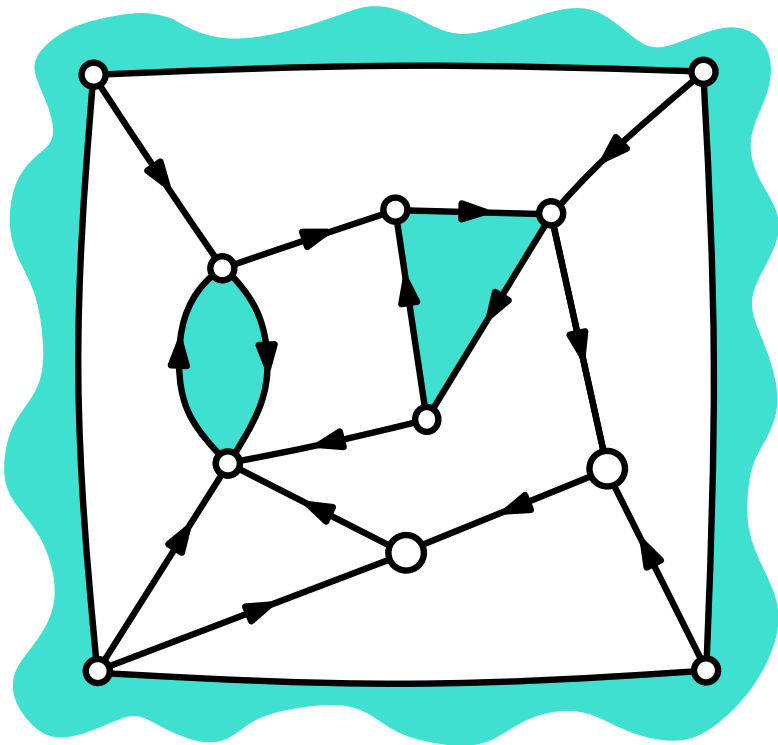
Orientations for maps with boundaries

For maps with boundaries we consider orientations such that **every inner boundary is a cw cycle** and the outer cycle is a boundary.

These are called **boundary-orientations**

To apply the mobile construction we still require the orientations to satisfy:

- the outer d -gon is a **source** (no ingoing edge)
- every vertex can be **reached** by a directed path starting from the source
- **there is no ccw cycle**

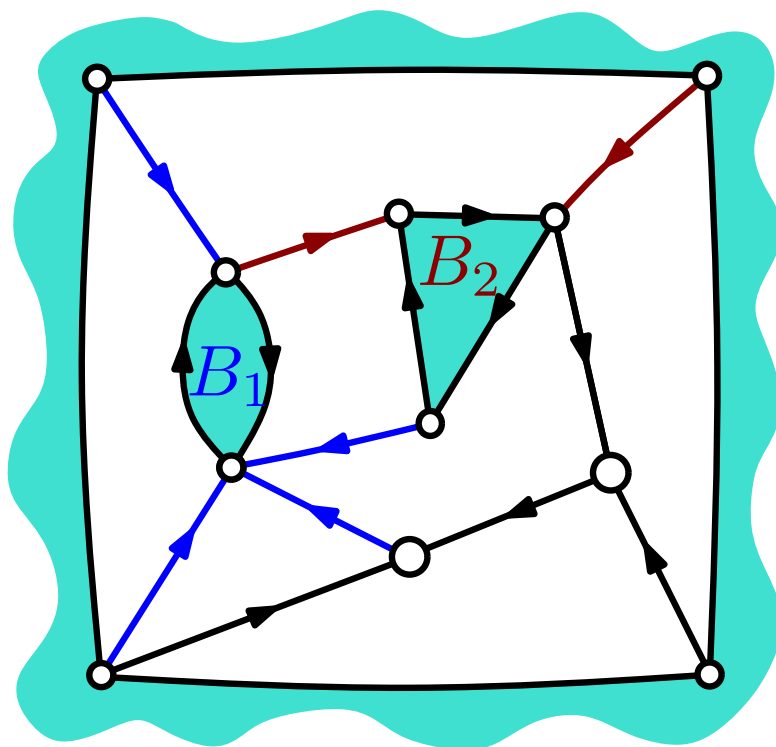


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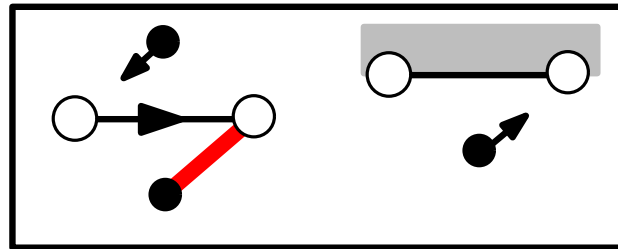
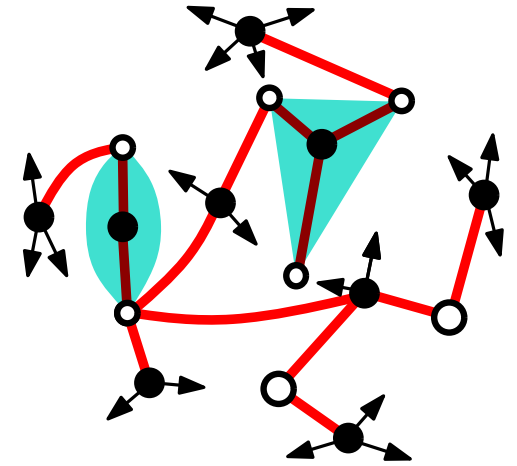
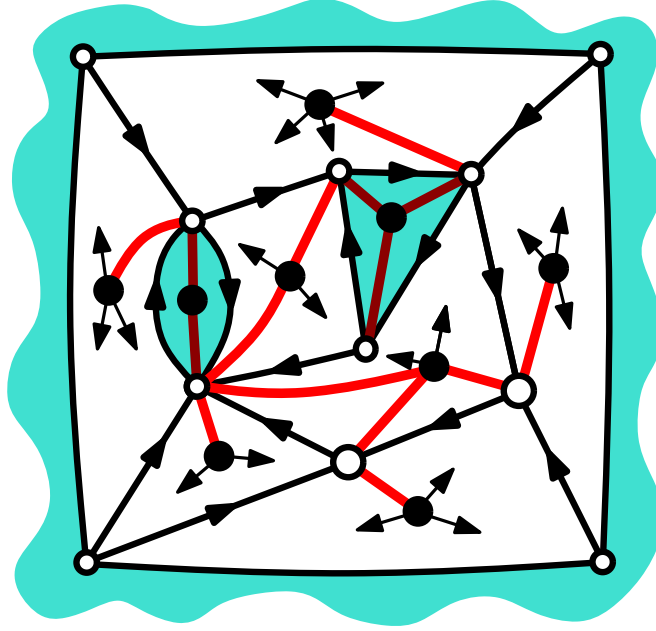
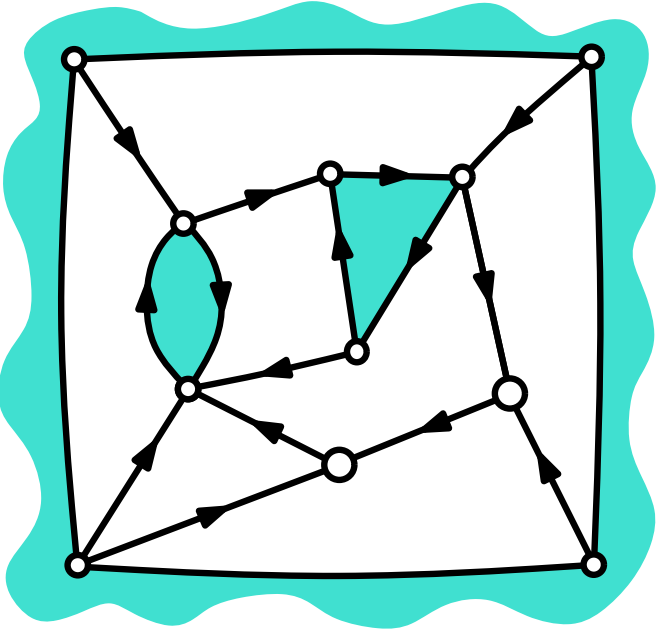
- the outer d -gon is a **source** (no ingoing edge)
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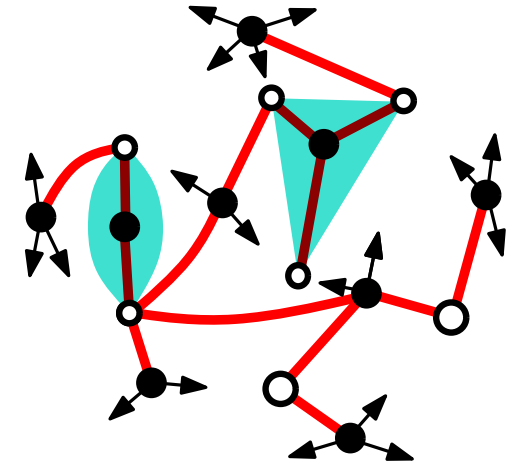
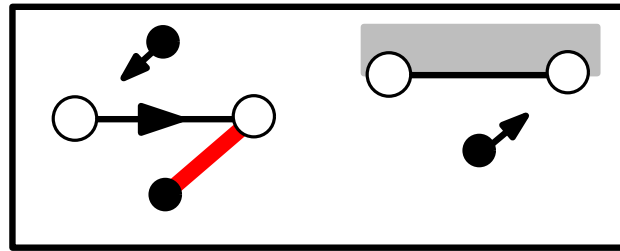
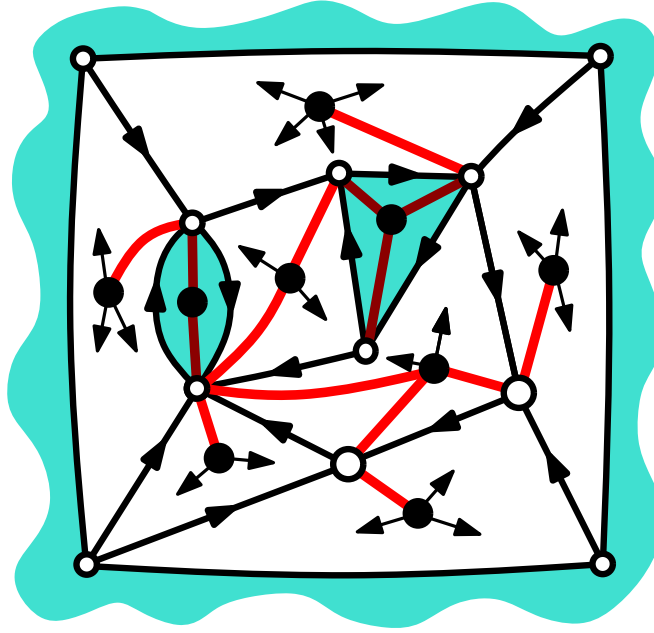
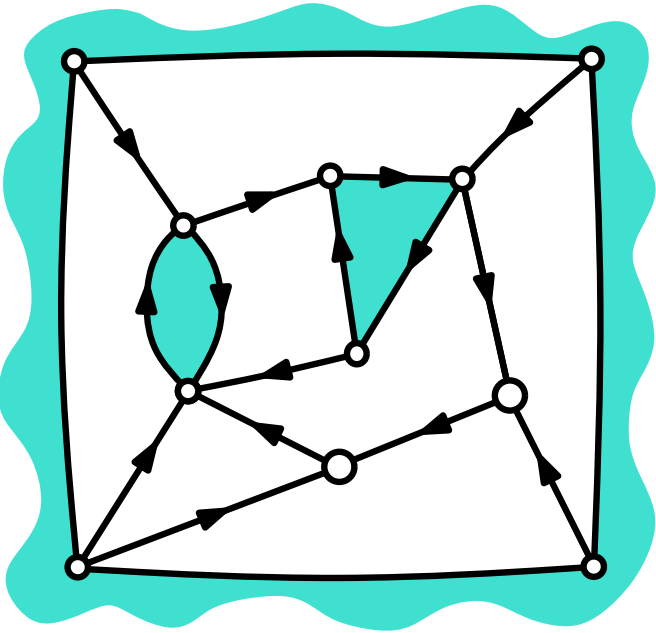
indegree of a boundary B :
total number of edges toward B

B_1 has indegree 4
 B_2 has indegree 2

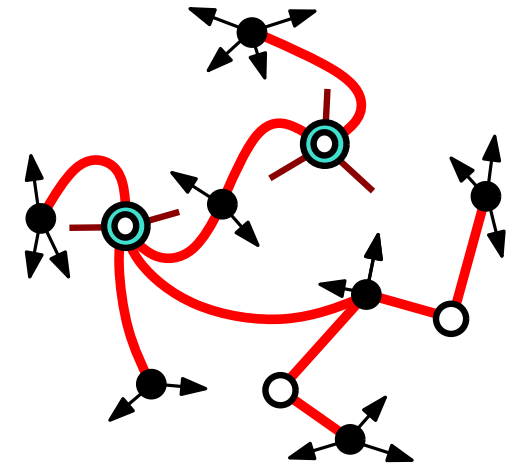
Extension of the bijection Φ to this setting



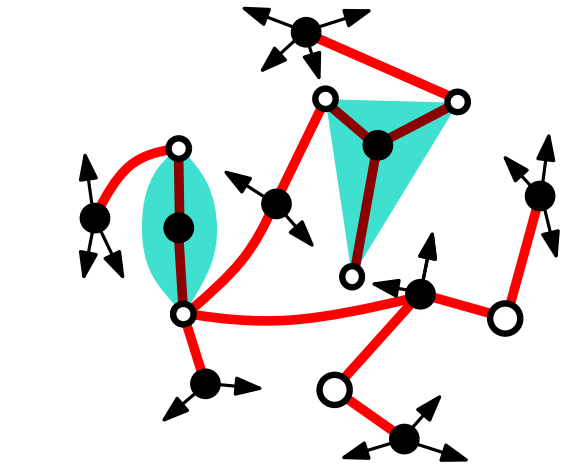
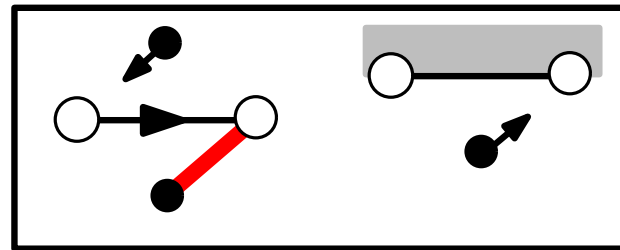
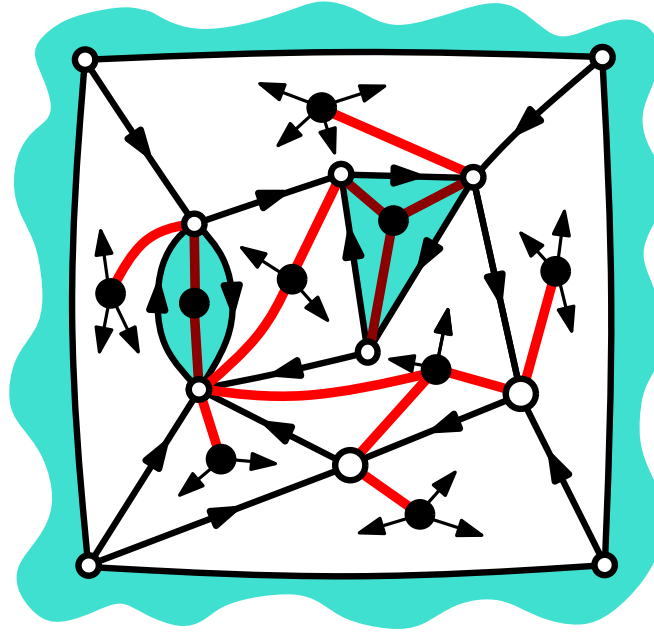
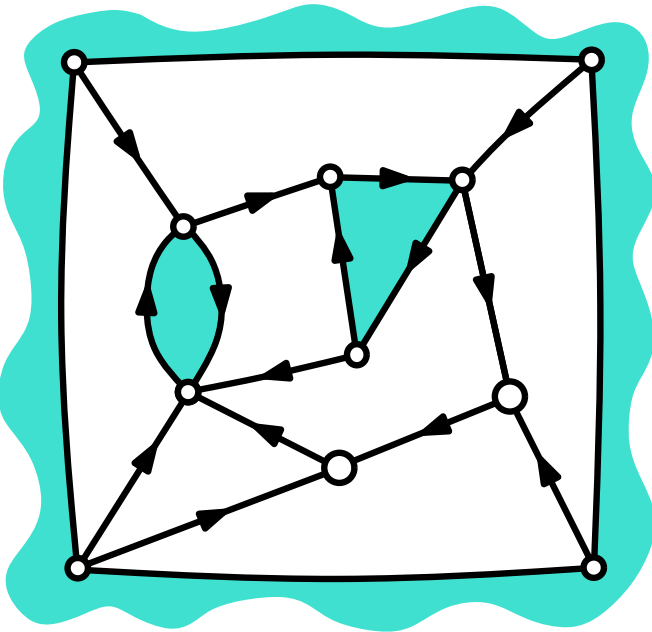
Extension of the bijection Φ to this setting



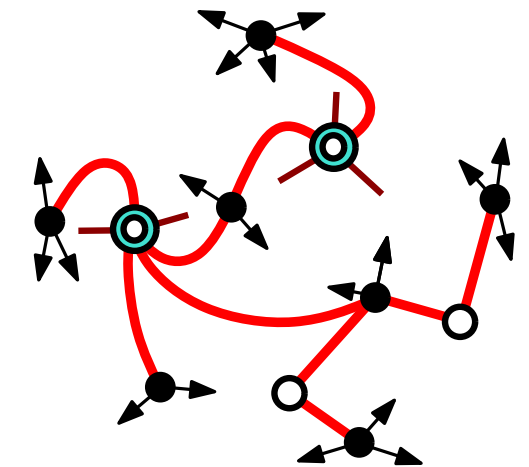
⇓ contraction of boundaries



Extension of the bijection Φ to this setting



⇓ contraction of boundaries

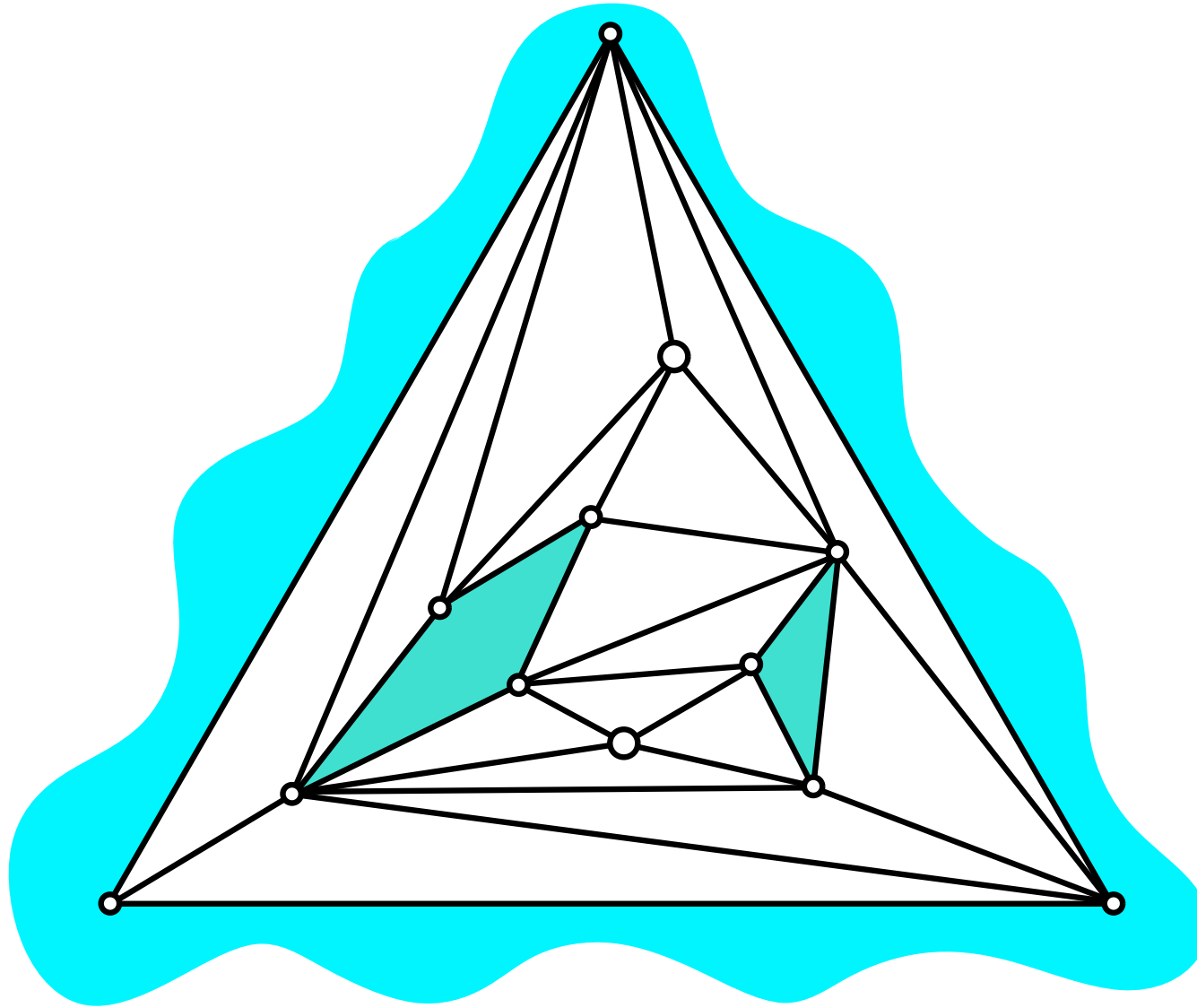


vertex \circ of indegree k
 internal boundary $\left\{ \begin{array}{l} \text{degré } r \\ b \text{ entrantes} \end{array} \right.$
 internal face of degree p



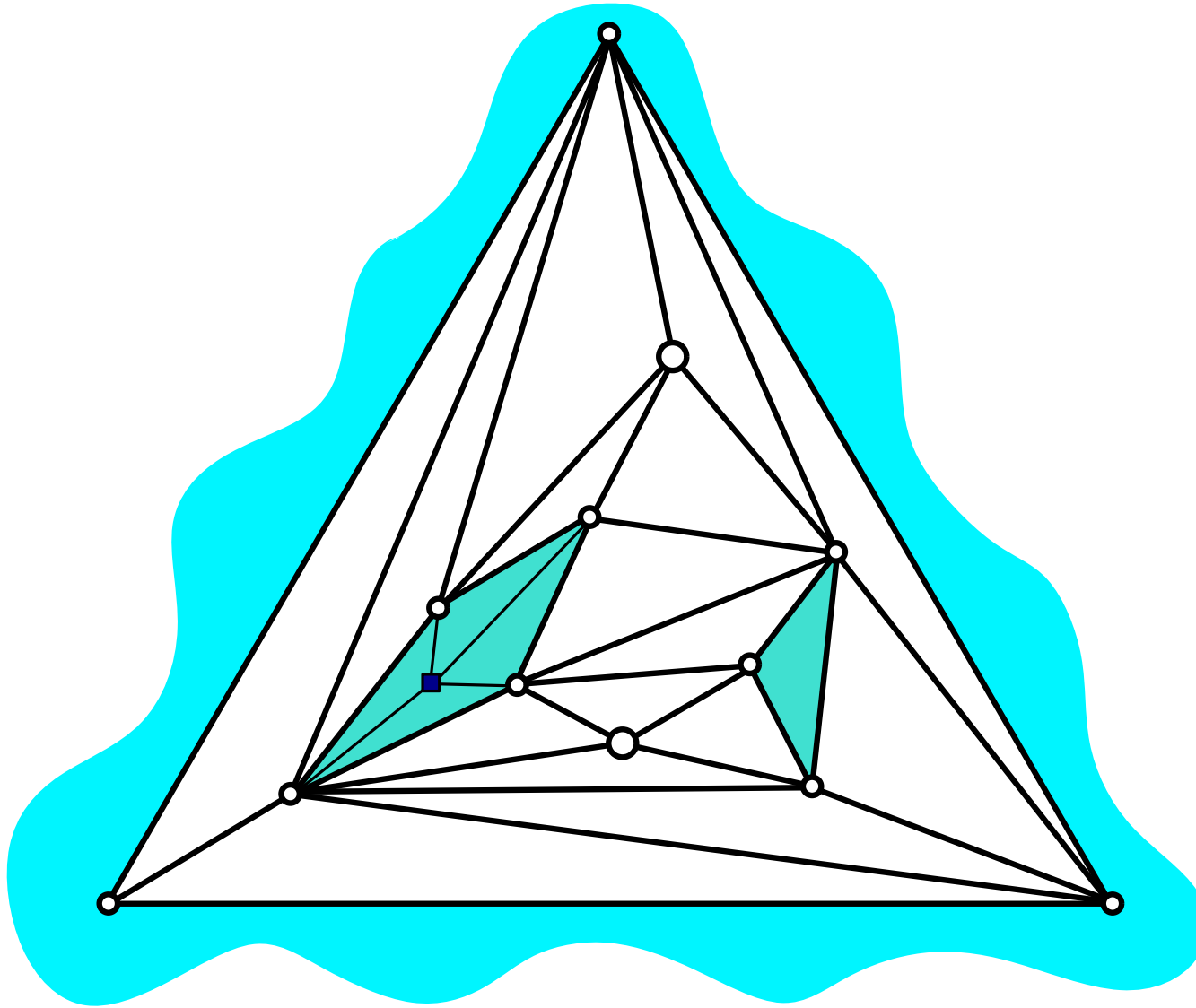
vertex \circ of degree k
 vertex \odot $\left\{ \begin{array}{l} r \text{ legs} \\ b \text{ neighbours} \end{array} \right.$
 vertex \bullet of degree p

Orientations for simple triangulations with boundaries



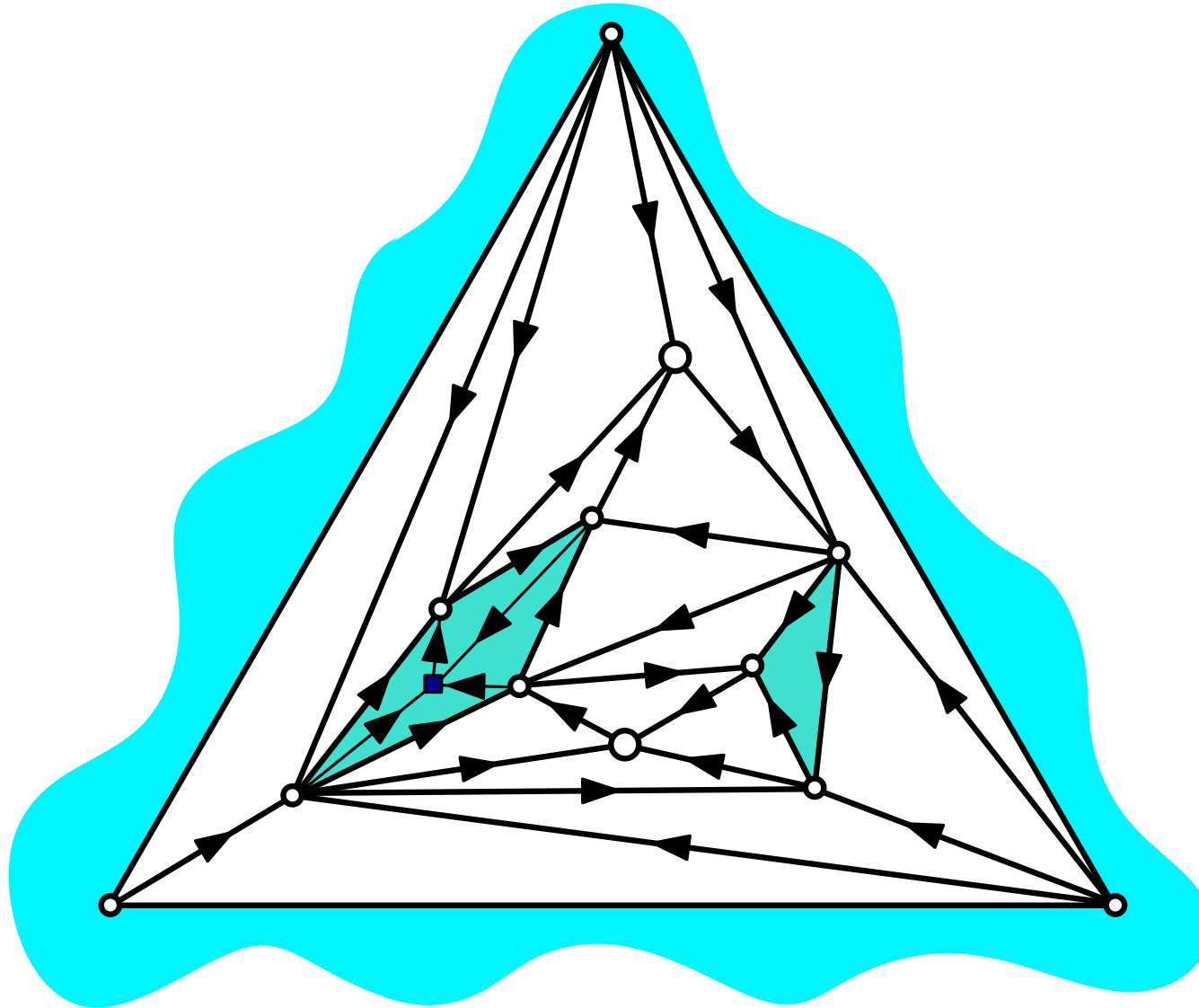
Orientations for simple triangulations with boundaries

Triangulate each inner boundary of length > 3



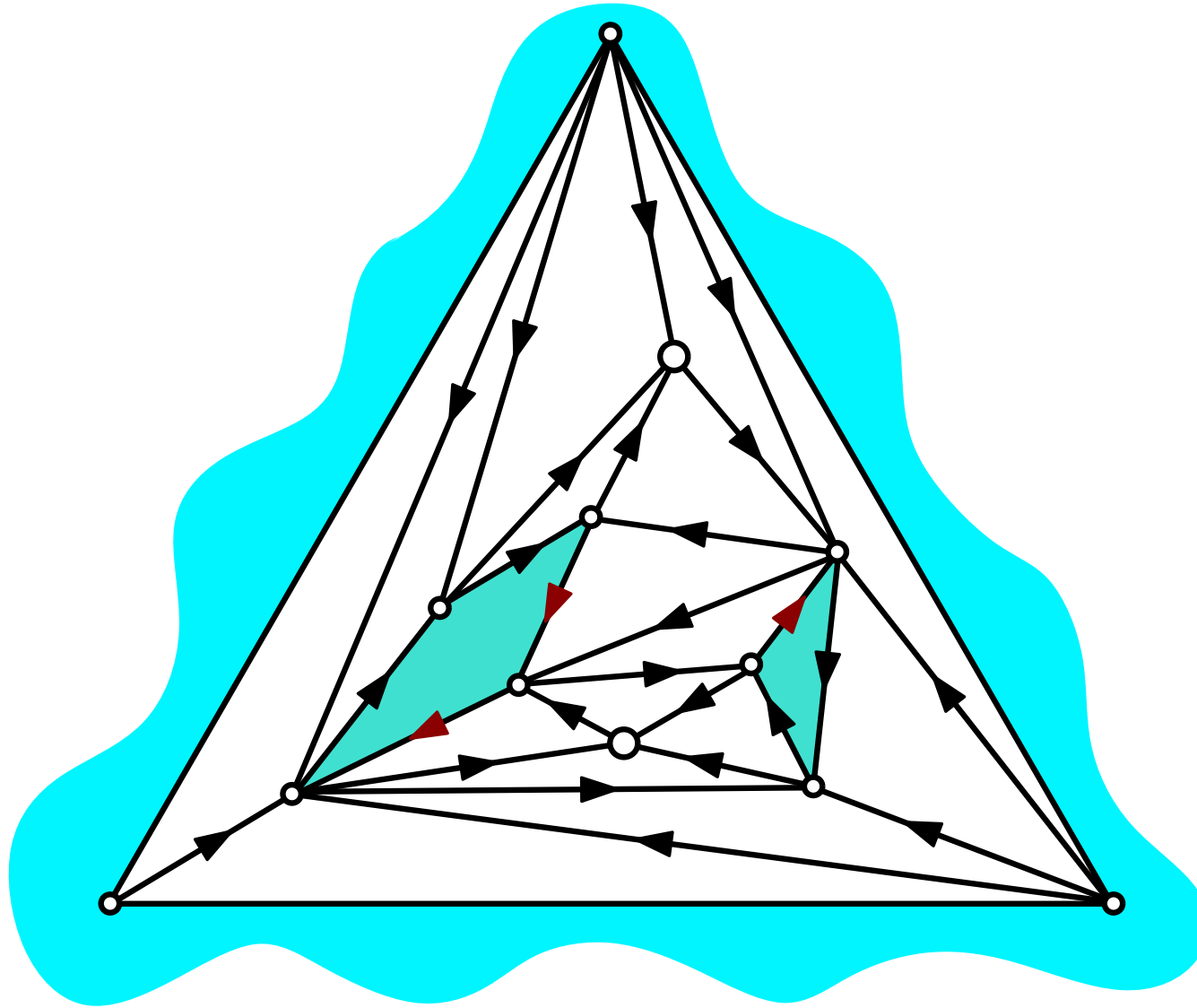
Orientations for simple triangulations with boundaries

Triangulate each inner boundary of length > 3
and compute the minimal 3-orientation



Orientations for simple triangulations with boundaries

delete the added edges inside boundaries
and reorient the inner boundaries as cw cycles

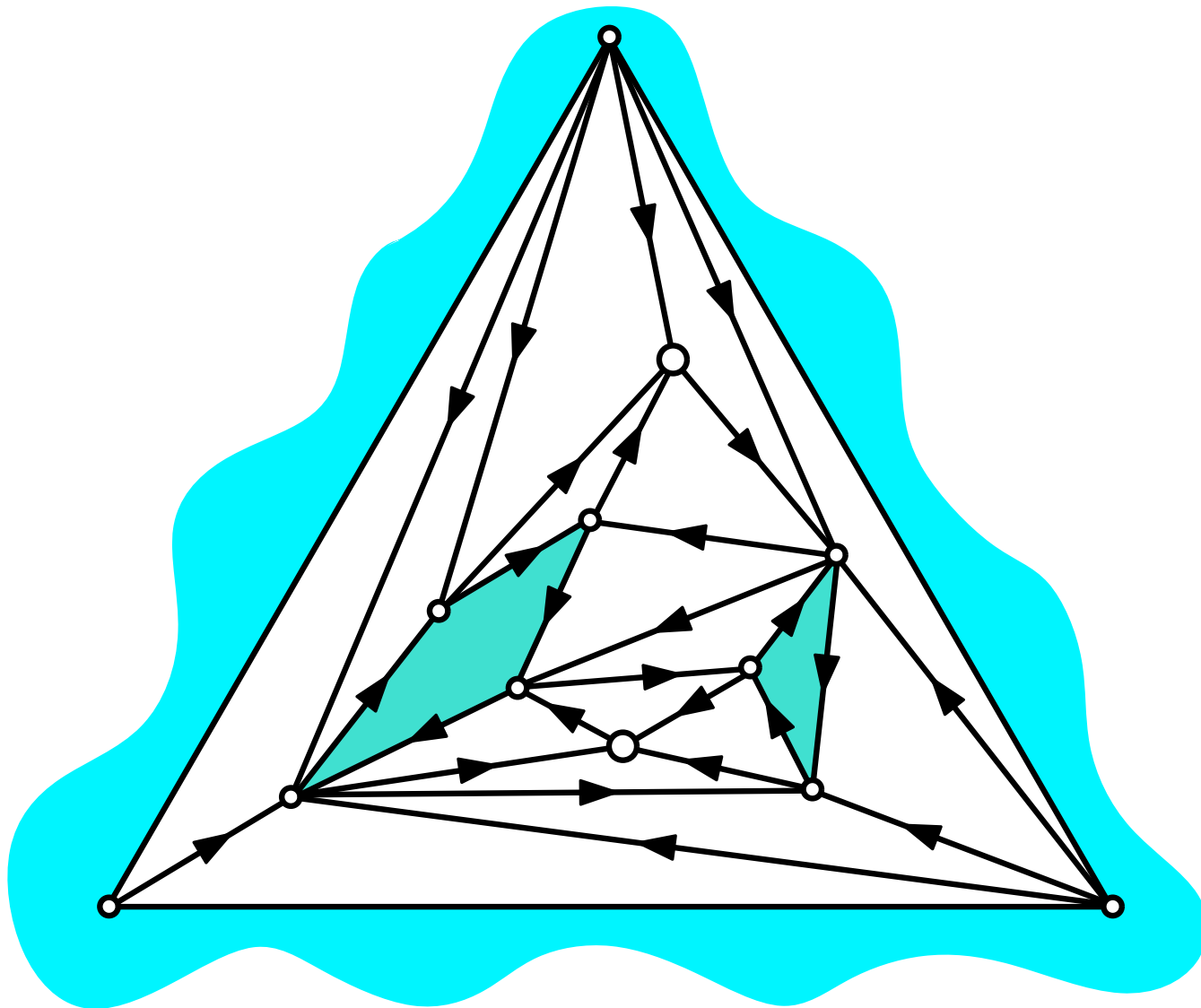


Orientations for simple triangulations with boundaries

Each inner boundary of length i has indegree $i + 3$

Each internal vertex has indegree 3

Such a boundary-orientation is called a **pseudo-3-orientation**



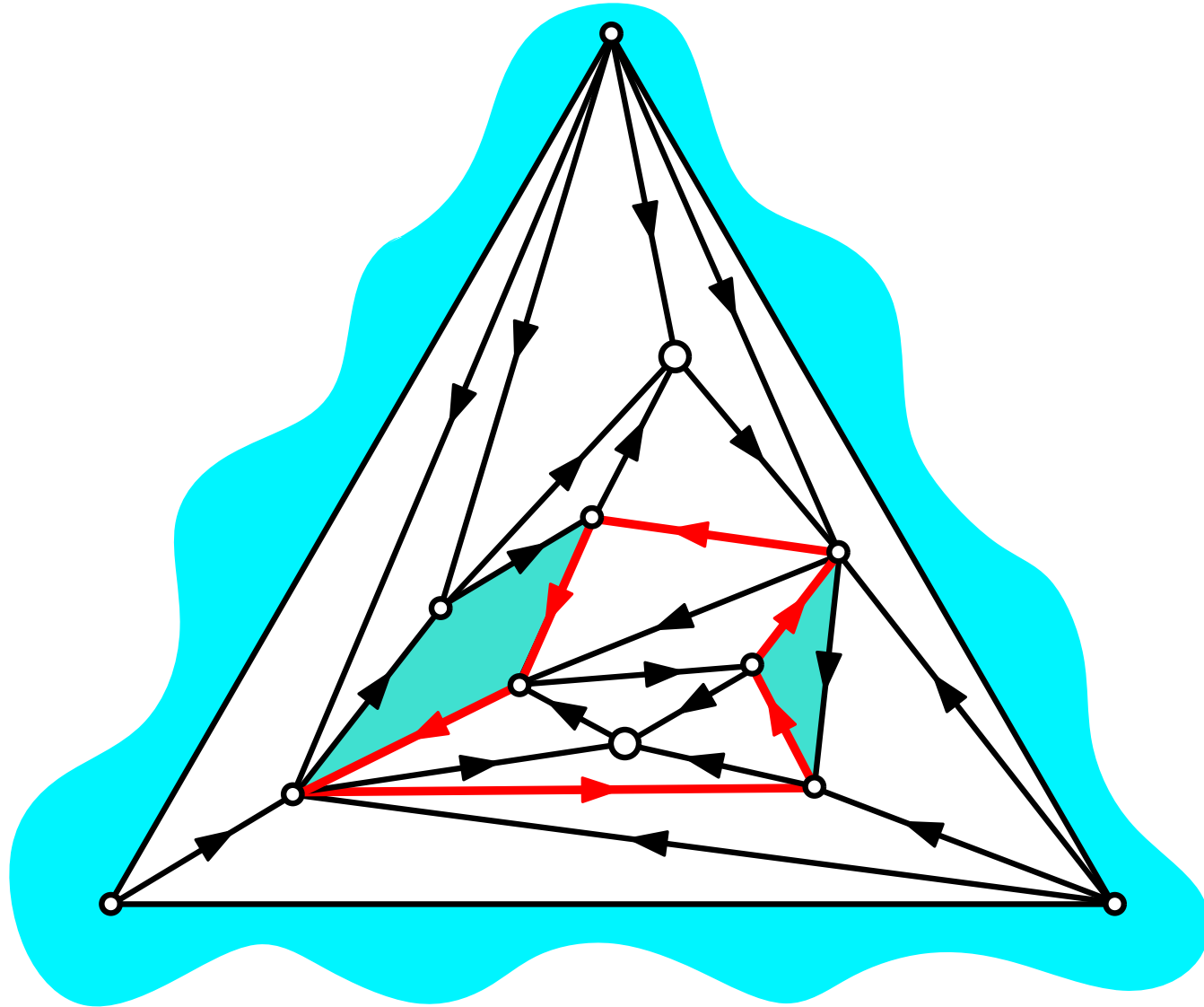
Orientations for simple triangulations with boundaries

Each inner boundary of length i has indegree $i + 3$

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Such a boundary-orientation is called a **pseudo-3-orientation**

Take the **minimal** such orientation (no ccw cycle)



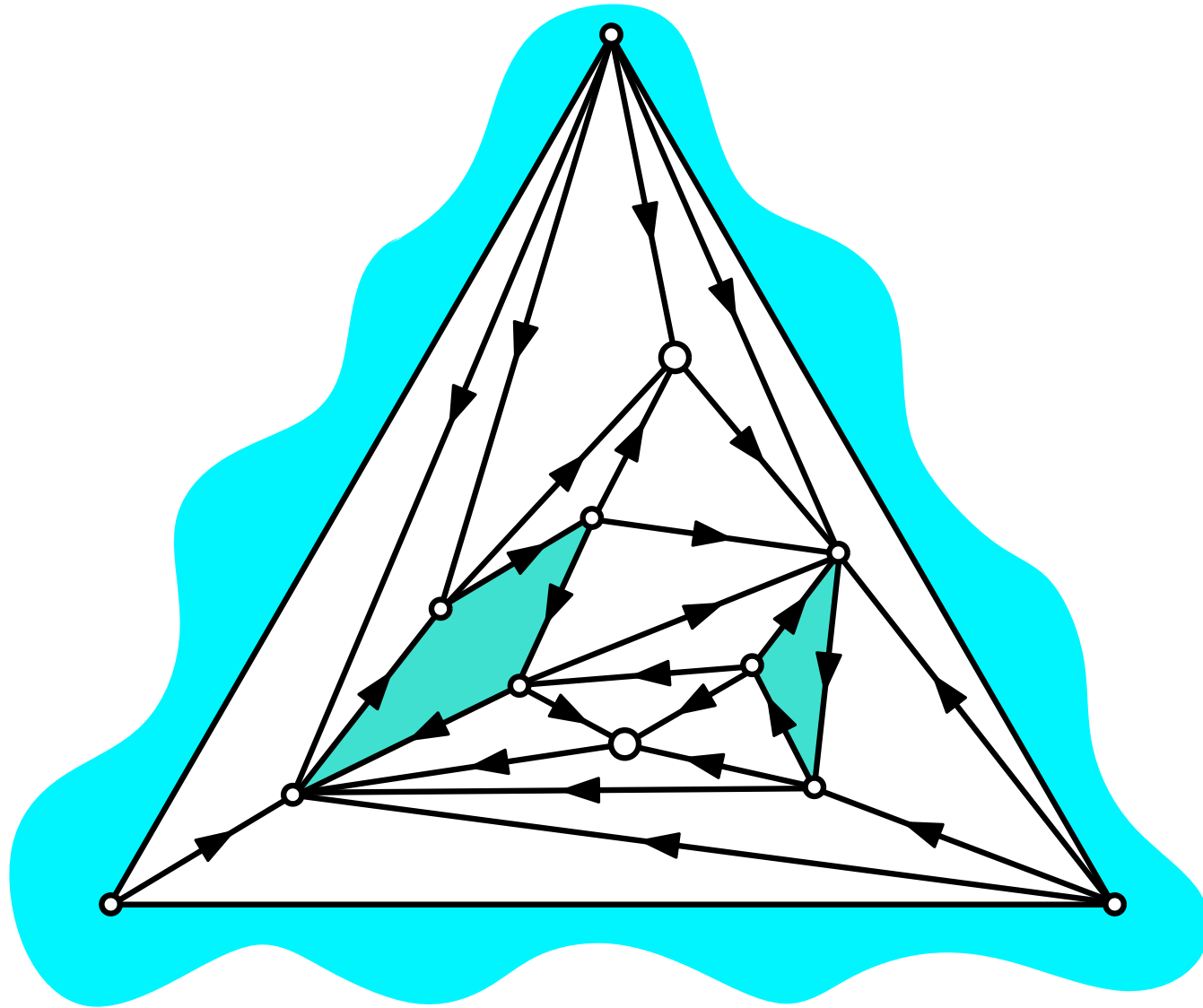
Orientations for simple triangulations with boundaries

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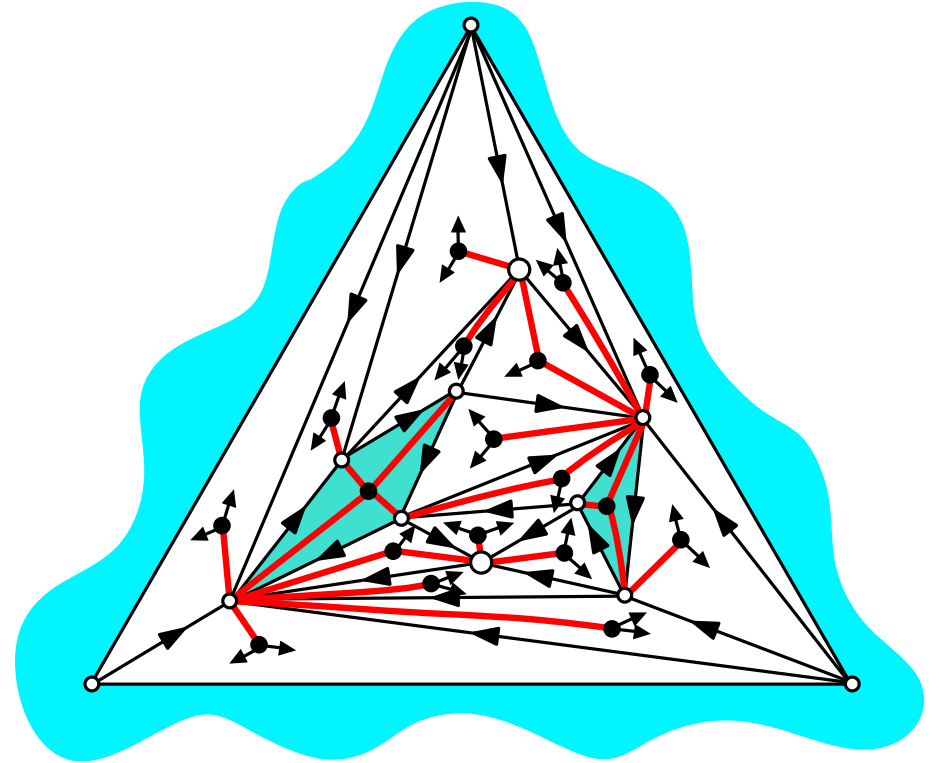
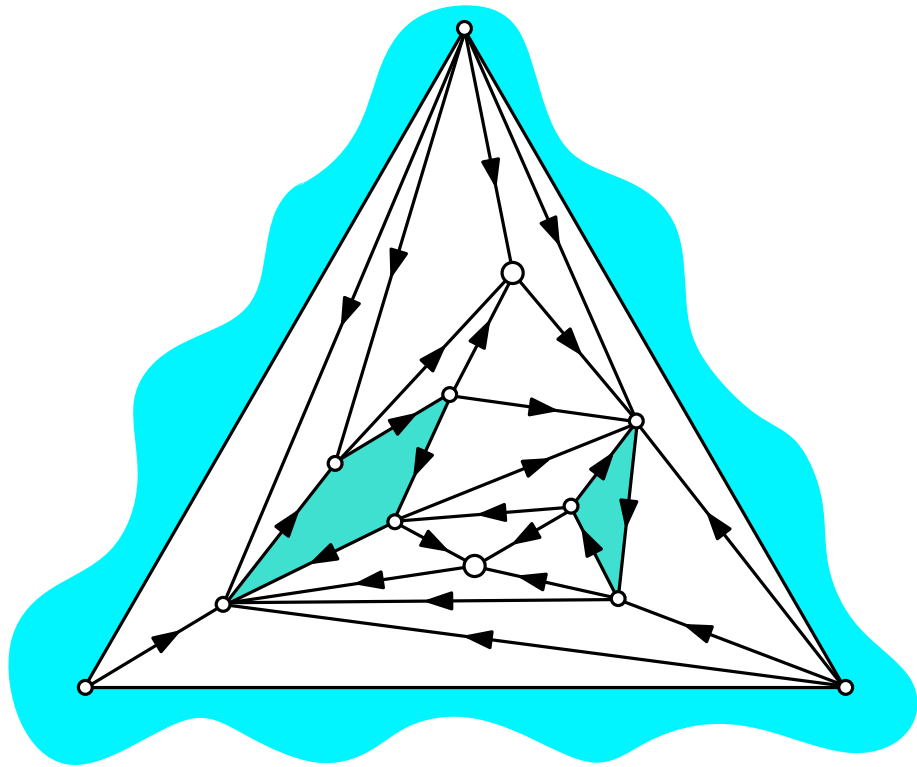
Such a boundary-orientation is called a **pseudo-3-orientation**

Take the **minimal** such orientation (no ccw cycle)



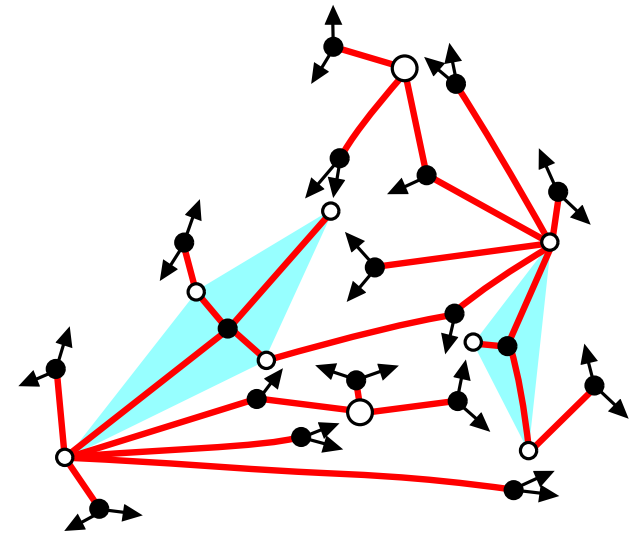
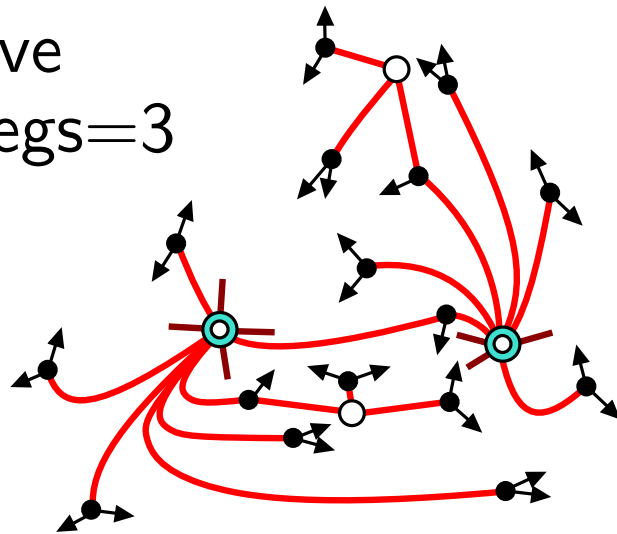
Mobiles for simple triangulations with boundaries

Apply the bijection Φ to the minimal pseudo-3-orientation



white vertices have
 $\#neighbours - \#legs = 3$

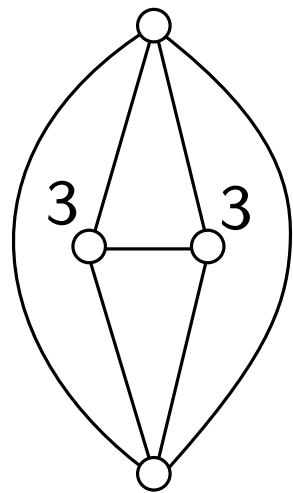
black vertices
have degree 3



Obstacles for the existence of pseudo-3-orientations

Not all 2-cycles are forbidden!

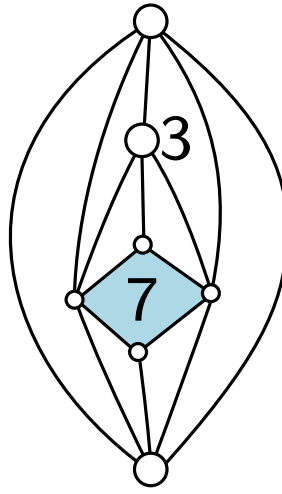
contractible 2-cycle



5 edges inside
total indegree 6 inside

Forbidden

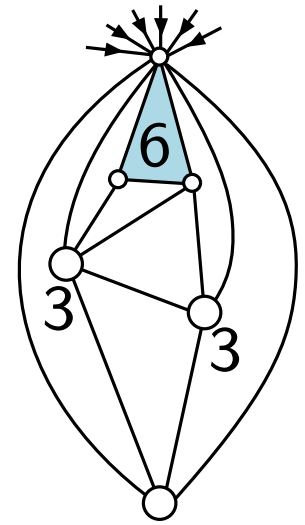
non-contractible 2-cycle
not touching any boundary
from the inside



9 non-boundary edges inside
total indegree 10 inside

Forbidden

non-contractible 2-cycle
touching a boundary
from the inside



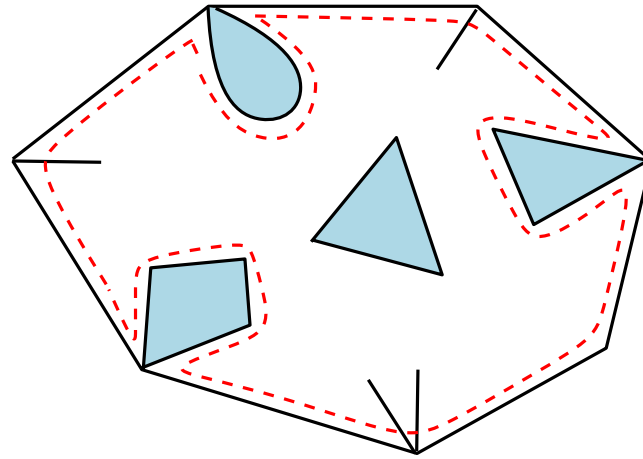
8 non-boundary edges inside
total indegree 6 inside

Not forbidden

Pseudo-girth parameter

For a map with boundaries that is planarly embedded

pseudo-girth = length of a shortest curve of the form



curve of length 15

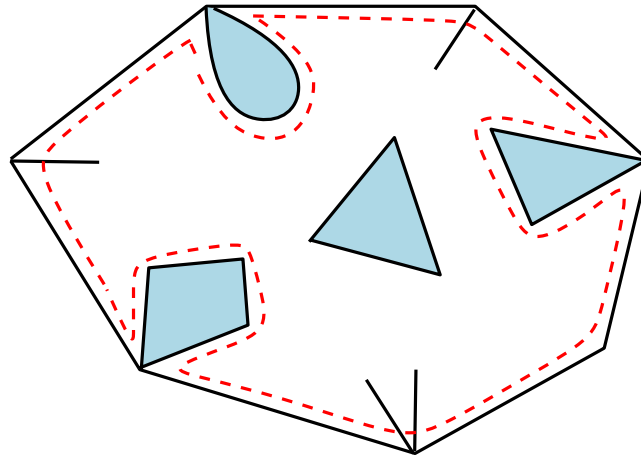
(curve that is the outer border of a region consisting of non-boundary faces)

Rk: $\text{girth} \leq \text{pseudo-girth} \leq \text{contractible girth}$

Pseudo-girth parameter

For a map with boundaries that is planarly embedded

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curve of length 15

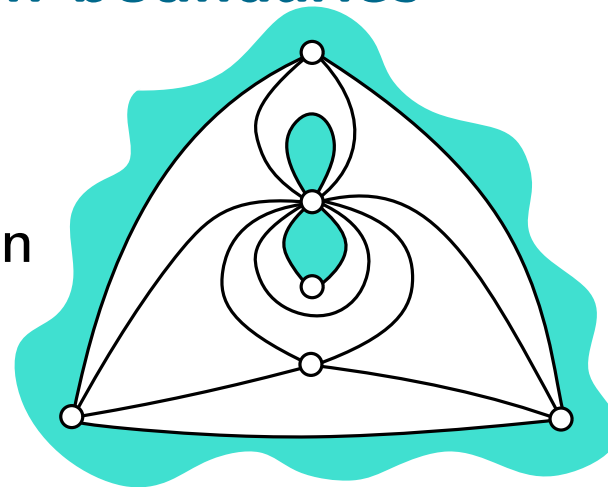
(curve that is the outer border of a region consisting of non-boundary faces)

Rk: $\text{girth} \leq \text{pseudo-girth} \leq \text{contractible girth}$

The map is called **pseudo-simple** if the pseudo-girth is ≥ 3

Results for pseudo-simple triangulations with boundaries

A triangulation with boundaries
(outer face being a triangular boundary-face)
is pseudo simple iff admits a pseudo 3-orientation



bijection with explicit mobiles

internal face (degree 3) \longleftrightarrow black vertex of degree 3
inner boundary of length i \longleftrightarrow white vertex with i legs
and $i + 3$ neighbours

Counting formula:

Let $N[n; a, k_1, \dots, k_r]$ be the number of pseudo-simple triangulations where:

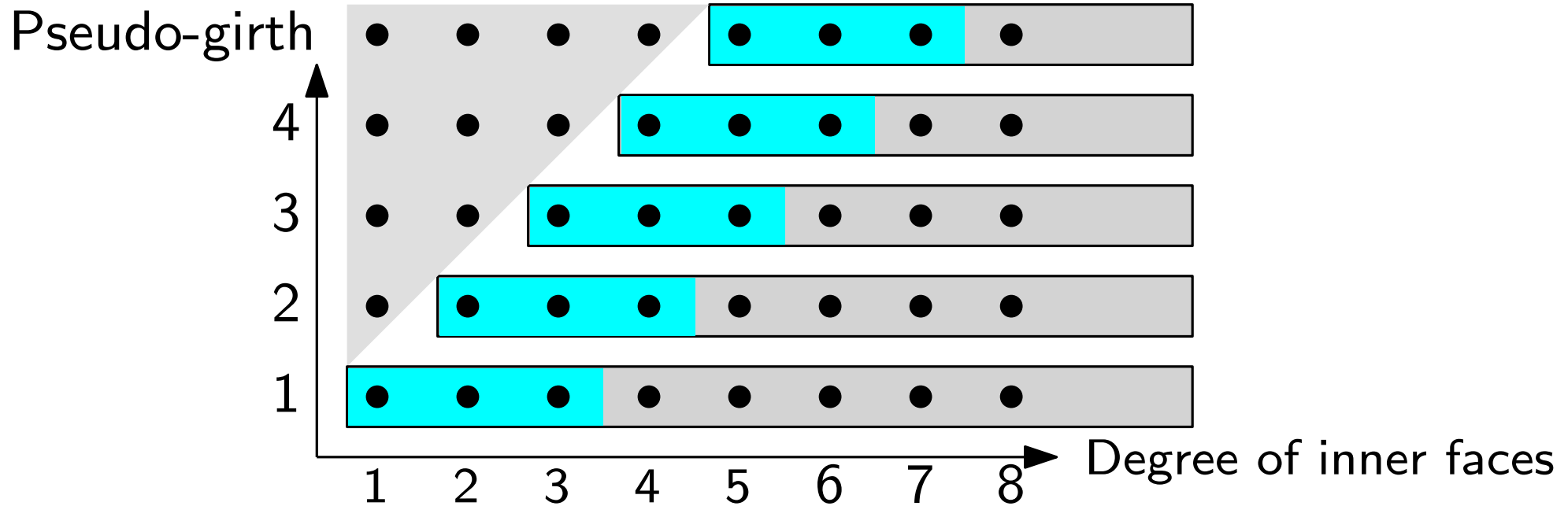
- the outer boundary has length a
- the inner boundaries B_1, \dots, B_r have lengths k_1, \dots, k_r
- there are n internal vertices
- in every boundary, a vertex is distinguished

$$N[n; a; k_1, \dots, k_r] = \frac{2(2a - 3)!}{(a - 3)!(a - 1)!} \frac{(4n + 4r + 2L - 5)!}{(n - 1)!(3m + 4r + 2L - 3)!} \prod_{i=1}^r i \binom{2i + 2}{i}$$

where $L = a + \sum_{i=1}^r k_i$ (total boundary length)

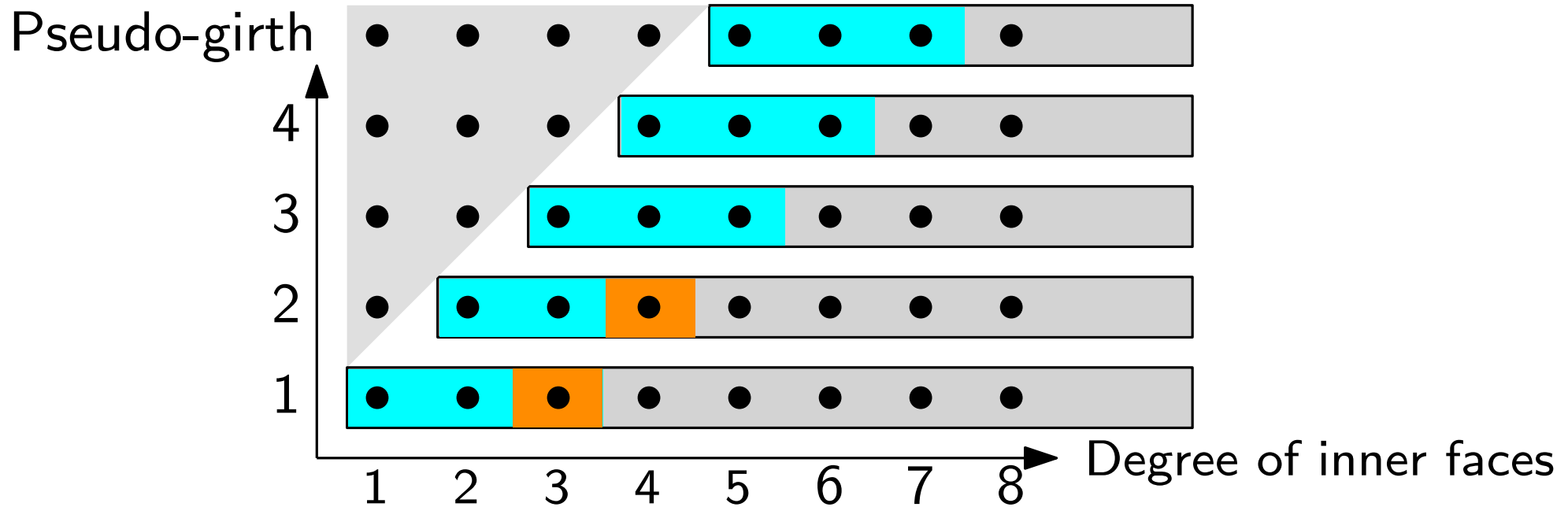
Results in any given pseudo-girth

We have a bijection in each pseudo-girth $d \geq 1$
for maps with boundaries, with inner face degrees in $\{d, d + 1, d + 2\}$



Results in any given pseudo-girth

[Bernardi, F'15] We have a bijection in each pseudo-girth $d \geq 1$
for maps with boundaries, with inner face degrees in $\{d, d+1, d+2\}$



Pseudo-girth-constraint is void for

$d = 1$ (recover Krikun's formula)

$d = 2$ bipartite case (new formula for quadrangulations with boundaries)

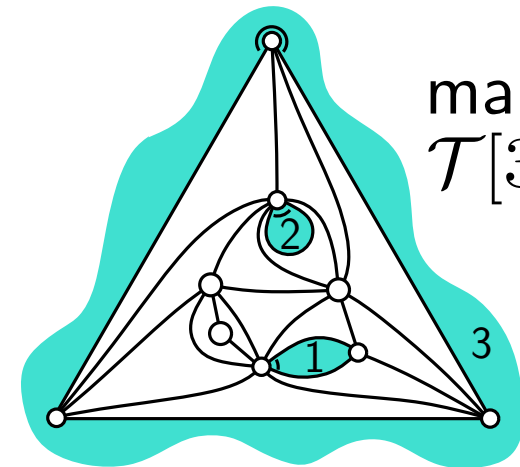
Factorized counting formulas

- Let $m \geq 0$ and ℓ_1, \dots, ℓ_r positive integers
- Let $\mathcal{T}[m; \ell_1, \dots, \ell_r]$ (resp. $\mathcal{Q}[m; \ell_1, \dots, \ell_r]$) be the set of **triangulations** (resp. **quadrangulations**) with r boundaries B_1, \dots, B_r s.t.
 - there are m internal vertices
 - every boundary B_i has length ℓ_i and a marked corner

Triangulations : Krikun's formula (2007)

$$|\mathcal{T}[m; a_1, \dots, a_r]| = \frac{4^k (e-2)!!}{m! (2b+k)!!} \prod_{i=1}^r a_i \binom{2a_i}{a_i}$$

with $b = \sum_i a_i$, $k = r + m - 2$, and $e = 2b + 3k$

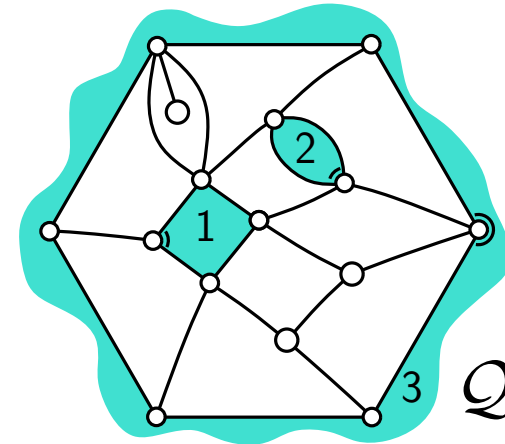


map \in
 $\mathcal{T}[3; 2, 1, 3]$

Quadrangulations : [Bernardi, F'15]

$$|\mathcal{Q}[m; 2a_1, \dots, 2a_r]| = \frac{3^k (e-1)!}{m! (3b+k)!} \prod_{i=1}^r 2a_i \binom{3a_i}{a_i}$$

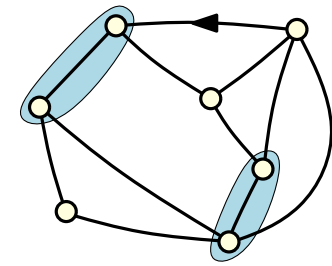
with $b = \sum_i a_i$, $k = r + m - 2$, and $e = 2b + 3k$



map \in
 $\mathcal{Q}[3; 4, 2, 6]$

Solution of the dimer model on quadrangulations

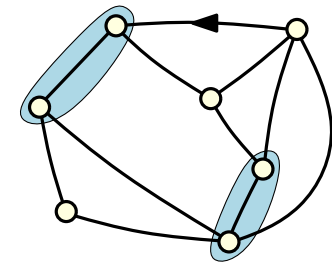
Map with dimers = pair (M, X) where M is a map and X is a subset of edges giving a partial-matching



map
with
2 dimers

Solution of the dimer model on quadrangulations

Map with dimers = pair (M, X) where M is a map and X is a subset of edges giving a partial-matching



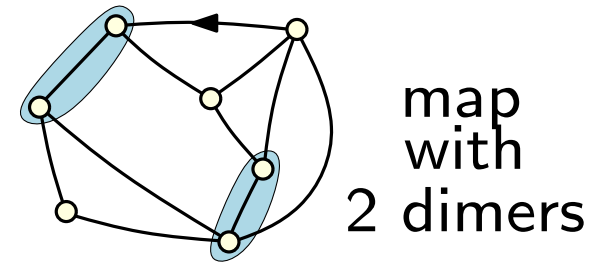
map
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- Dimer model on (rooted) quadrangulations

Generating function :
$$F(t, w) = \sum_{\text{configurations}} t^{\#\text{faces}} w^{\#\text{dimers}}$$

Solution of the dimer model on quadrangulations

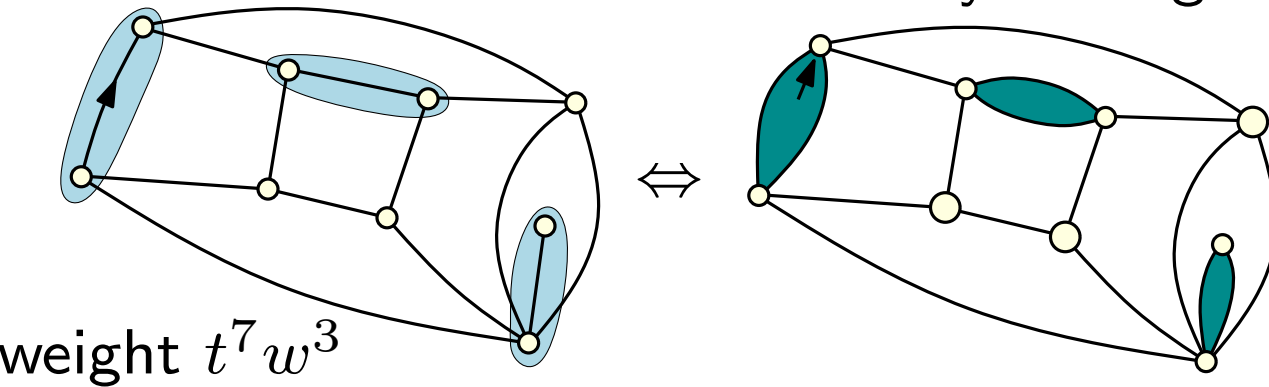
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- Dimer model on (rooted) quadrangulations

Generating function : $F(t, w) = \sum_{\text{configurations}} t^{\#\text{faces}} w^{\#\text{dimers}}$

dimer \leftrightarrow boundary of length 2



bijection

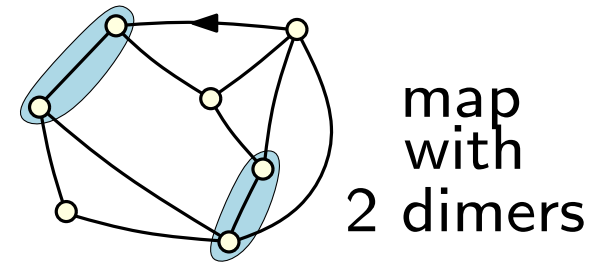


$$F(t, w) = R - 1 - t R^3 - 6wt^2 R^6$$

$$\text{où } R = 1 + 3tR^2 + 9wt^2 R^5$$

Solution of the dimer model on quadrangulations

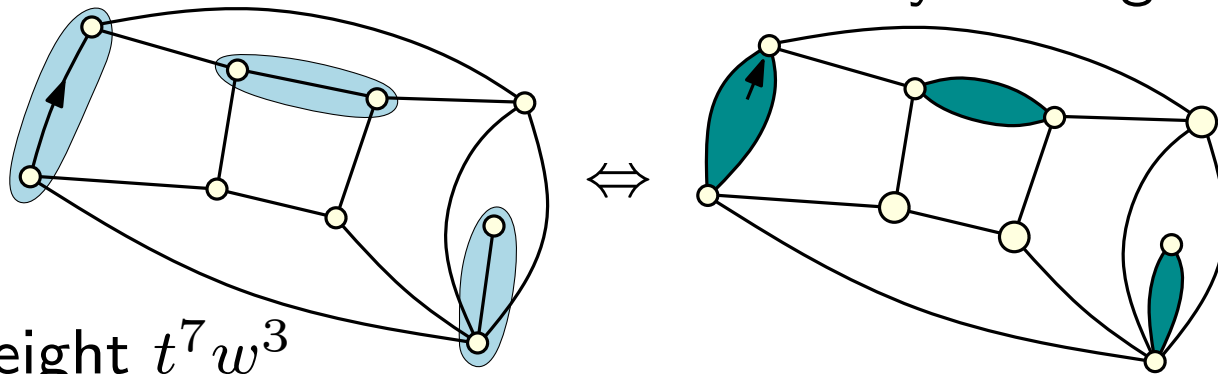
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bijection
 \Downarrow

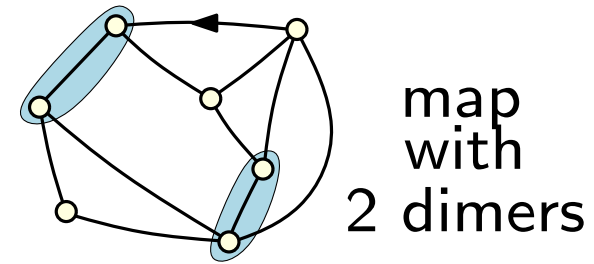
$$F(t, w) = R - 1 - t R^3 - 6wt^2 R^6$$

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Asymptotics : for $w \in \mathbb{R}$ fixed, $[t^n]F \sim c_w \gamma_w^n n^{-5/2}$
 except at **critical weight** $w_0 = -3/125$ where $[t^n]F \sim c_0 \gamma_0^n n^{-7/3}$

Solution of the dimer model on quadrangulations

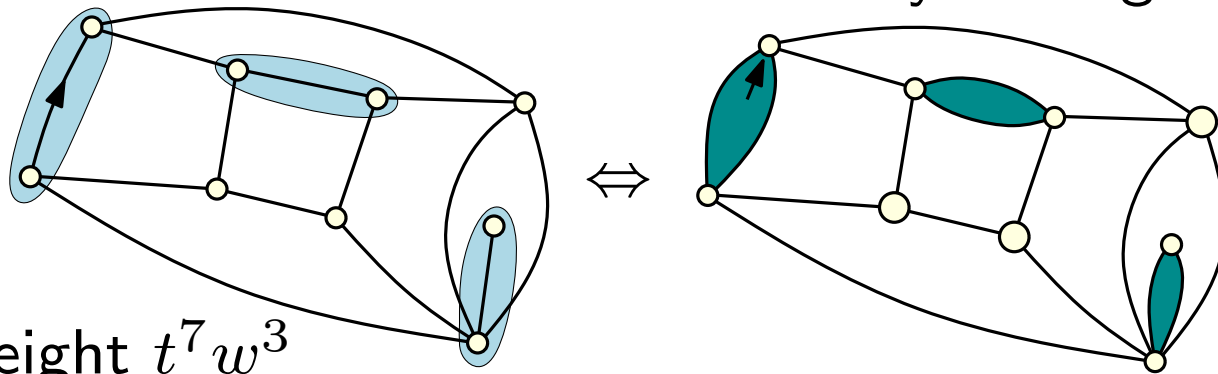
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weight $t^7 w^3$

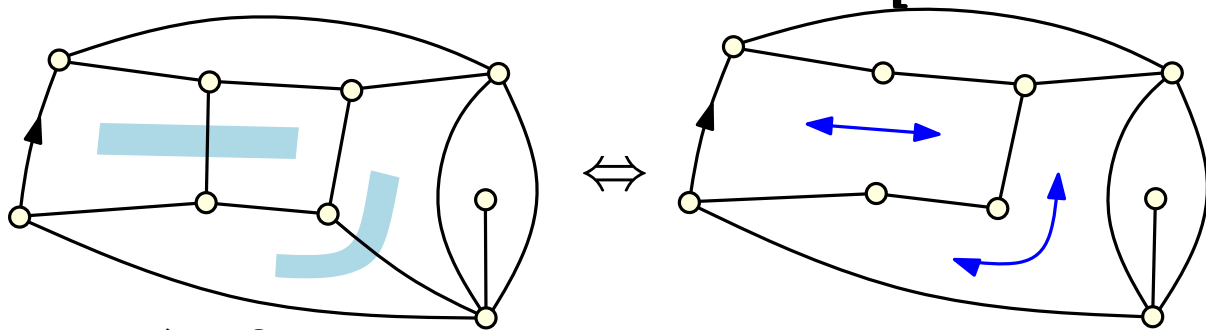
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- Solution of the dual model in [Bouttier, Di Francesco, Guitter'03]



poids $t^7 w^2$

weight t per square face
 weight $3t^2 w$ per hexagonal face

$$F(t, w) = R - 1 - t R^3 - 15wt^2 R^4$$

$$\text{où } R = 1 + 3tR^2 + 30wt^2 R^3$$

critical weight $w_0 = -1/10$
 where **typical distance** $\approx n^{1/6}$