

# Combinatorics of characters and continuation of Li.

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Collaboration at various stages of the work  
and in the framework of the Project

*Evolution Equations in Combinatorics and Physics* :

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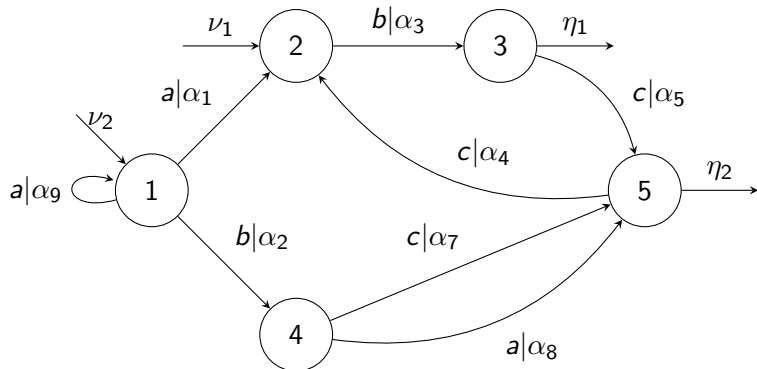
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# Multiplicity Automaton (Eilenberg, Schützenberger)



1 S. Eilenberg, *Automata, Languages, and Machines (Vol. A)* Acad. Press, New York, 1974

2 M.P. Schützenberger, *On the definition of a family of automata*, *Inf. and Contr.*, 4 (1961), 245-270.

# Multiplicity automaton (linear representation) & behaviour

## Linear representation

$$\nu = (\nu_2 \quad \nu_1 \quad 0 \quad 0 \quad 0), \quad \eta = (0 \quad 0 \quad \eta_1 \quad 0 \quad \eta_2)^T$$

$$\mu(a) = \begin{pmatrix} \alpha_9 & \alpha_1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & \alpha_8 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix} \quad \mu(b) = \begin{pmatrix} 0 & 0 & 0 & \alpha_2 & 0 \\ 0 & 0 & \alpha_3 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

$$\mu(c) = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & \alpha_5 \\ 0 & 0 & 0 & 0 & \alpha_7 \\ 0 & \alpha_4 & 0 & 0 & 0 \end{pmatrix}$$

## Behaviour

$$\mathcal{A}(w) = \nu \mu(w) \eta = \sum_{\substack{i,j \\ \text{states}}} \nu(i) \underbrace{\left( \sum \text{weight}(p) \right)}_{\substack{\text{weight of all paths } \textcircled{i} \rightarrow \textcircled{j} \\ \text{with label } w}} \eta(j)$$

# Operations and definitions on series

**Addition, Scaling:** as for functions because  $R\langle\langle X \rangle\rangle = R^{X^*}$

**Concatenation:**  $f.g(w) = \sum_{w=uv} f(u)g(v)$

**Polynomials:** Series s.t.  $\text{supp}(f) = \{w\}_{f(w) \neq 0}$  is finite.

The set of polynomials will be denoted  $R\langle X \rangle$ .

**Pairing:**  $\langle S | P \rangle = \sum_{w \in X^*} S(w)P(w)$  ( $S$  series,  $P$  polynomial)

**Summation:**  $\sum_{i \in I} S_i$  summable iff for all  $w \in X^*$ ,  $i \mapsto \langle S_i | w \rangle$  is finitely supported. This corresponds to the product topology (with  $R$  discrete). In particular, we have

$$\sum_{i \in I} S_i := \sum_{w \in X^*} \left( \sum_{i \in I} \langle S_i | w \rangle \right) w$$

**Star:** For all series  $S$  s.t.  $\langle S | 1_{X^*} \rangle = 0$ , the family  $(S^n)_{n \geq 0}$  is summable and we set  $S^* := \sum_{n \geq 0} S^n = 1 + S + S^2 + \dots (= (1 - S)^{-1})$ .

**Shifts:**  $\langle u^{-1}S | w \rangle = \langle S | uw \rangle$ ,  $\langle Su^{-1} | w \rangle = \langle S | wu \rangle$

# Rational series (Sweedler & Schützenberger)

## Theorem A

Let  $S \in k\langle\langle X \rangle\rangle$  TFAE

- i) The family  $(Su^{-1})_{u \in X^*}$  is of finite rank.
- ii) The family  $(u^{-1}S)_{u \in X^*}$  is of finite rank.
- iii) The family  $(u^{-1}Sv^{-1})_{u,v \in X^*}$  is of finite rank.
- iv) It exists  $n \in \mathbb{N}$ ,  $\lambda \in k^{1 \times n}$ ,  $\mu : X^* \rightarrow k^{n \times n}$  (a multiplicative morphism) and  $\gamma \in k^{n \times 1}$  such that, for all  $w \in X^*$

$$(S, w) = \lambda \mu(w) \gamma \quad (1)$$

- v) The series  $S$  is in the closure of  $k\langle X \rangle$  for  $(+, conc, *)$  within  $k\langle\langle X \rangle\rangle$ .

## Definition

A series which fulfill one of the conditions of Theorem A will be called *rational*. The set of these series will be denoted by  $k^{rat}\langle\langle X \rangle\rangle$ .

# Sweedler's duals

## Remarks

- 1 (i  $\leftrightarrow$  iii) needs  $k$  to be a field.
- 2 (iv) needs  $X$  to be finite, (iv  $\leftrightarrow$  v) is known as the theorem of Kleene-Schützenberger (M.P. Schützenberger, *On the definition of a family of automata, Inf. and Contr.*, 4 (1961), 245-270.)
- 3 For the sake of Combinatorial Physics (where the alphabets can be infinite), **(iv)** has been extended to infinite alphabets and replaced by **iv')** The series  $S$  is in the rational closure of  $k^X$  (linear series) within  $k\langle\langle X \rangle\rangle$ .
- 4 This theorem is linked to the following: Representative functions on  $X^*$  (see Eichii Abe, Chari & Pressley), Sweedler's duals &c.
- 5 In the vein of (v) expressions like  $ab^*$  or identities like  $(ab^*)^*a^* = (a + b)^*$  (Lazard's elimination) will be called rational.

# From theory to practice: Schützenberger's calculus

## From series to automata

Starting from a series  $S$ , one has a way to construct an automaton (finite-stated iff the series is rational) providing that we know how to compute on shifts and one-letter-shifts will be sufficient due to the formula  $u^{-1}v^{-1}S = (vu)^{-1}S$ .

## Calculus on rational expressions

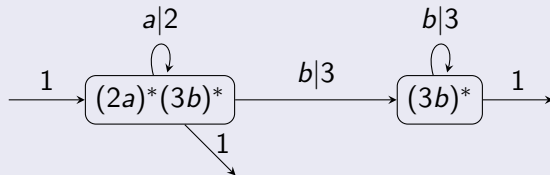
In the following,  $x$  is a letter,  $E, F$  are rational expressions (i.e. expressions built from letters by scalings, concatenations and stars)

- 1  $x^{-1}$  is (left and right) linear
- 2  $x^{-1}(E.F) = x^{-1}(E).F + \langle E \mid 1_{x^*} \rangle x^{-1}(F)$
- 3  $x^{-1}(E^*) = x^{-1}(E).E^*$

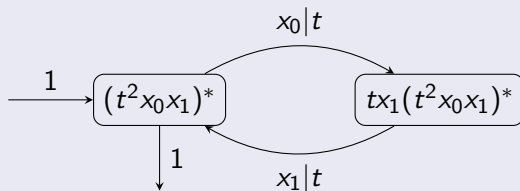


# Examples

With  $(2a)^*(3b)^*$  ;  $X = \{a, b\}$



With  $(t^2x_0x_1)^*$  ;  $X = \{x_0, x_1\}$



## From theory to practice: construction starting from $S$ .

- **States**  $\boxed{u^{-1}S}$  (constructed step by step)
- **Edges** We shift every state by letters (length) level by level (knowing that  $x^{-1}(u^{-1}S) = (ux)^{-1}S$ ). Two cases:  
**Returning state:** The state is a linear combination of the already created ones i.e.  $x^{-1}(u^{-1}S) = \sum_{v \in F} \alpha(ux, v)v^{-1}S$  (with  $F$  finite), then we set the edges

$$\boxed{u^{-1}S} \xrightarrow{x|\alpha_v} \boxed{v^{-1}S}$$

**The created state is new:** Then

$$\boxed{u^{-1}S} \xrightarrow{x|1} \boxed{x^{-1}(u^{-1}S)}$$

- **Input**  $\boxed{S}$  with the weight 1
- **Outputs** All states  $\boxed{T}$  with weight  $\langle T | 1_{X^*} \rangle$

## Link with conc-bialgebras (CAP 17)

We call here conc-bialgebras, structures such that  $\mathcal{B} = (k\langle X \rangle, \text{conc}, 1_{X^*}, \Delta, \epsilon)$  is a bialgebra and  $\Delta(X) \subset (kX \oplus k1_{X^*})^{\otimes 2}$ . For this, as  $k\langle X \rangle$  is a free algebra, it suffices to define  $\Delta$  and check the axioms on letters. Below, some examples

**Shuffle:**  $X$  is arbitrary  $\Delta(x) = x \otimes 1 + 1 \otimes x$  and

$$\Delta(w) = \sum_{I+J=[1 \dots |w|]} w[I] \otimes w[J]$$

**Stuffle:**  $Y = \{y_i\}_{i \geq 1}$ ,  $\Delta(y_k) = y_k \otimes 1 + 1 \otimes y_k + \sum_{i+j=k} y_i \otimes y_j$   
 **$q$ -infiltration:**  $X$  is arbitrary,  $\Delta(x) = x \otimes 1 + 1 \otimes x + qx \otimes x$  and

$$\Delta(w) = \sum_{I \cup J = [1 \dots |w|]} q^{|\mathcal{I} \cap \mathcal{J}|} w[I] \otimes w[J]$$

## Link with conc-bialgebras/2

In case  $\epsilon(P) = \langle P \mid 1_{X^*} \rangle^a$ , the restricted (graded) dual is  $\mathcal{B}^V = (k\langle X \rangle, *, 1_{X^*}, \Delta_{\text{conc}}, \epsilon)$  and we can write, for  $x \in X$

$$\Delta(x) = x \otimes 1_{X^*} + 1_{X^*} \otimes x + \Delta_+(x) \quad (2)$$

then, the dual law  $*$  ( $=^t \Delta$ ) can be defined by recursion

$$\begin{aligned} w * 1_{X^*} &= 1_{X^*} * w = w \\ au * bv &= a(u * bv) + b(au * v) + \varphi(a, b)(u * v) \end{aligned} \quad (3)$$

where  $\varphi =^t \Delta_+ : k.X \otimes k.X \rightarrow k.X$  is an associative law.

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<sup>a</sup>which covers all usual combinatorial cases, save Hadamard

# Some dual laws

Name	Formula (recursion)	$\varphi$	Type
Shuffle [21]	$au \sqcup bv = a(u \sqcup bv) + b(au \sqcup v)$	$\varphi \equiv 0$	I
Stuffle [19]	$x_i u \sqcup x_j v = x_i(u \sqcup x_j v) + x_j(x_i u \sqcup v) + x_{i+j}(u \sqcup v)$	$\varphi(x_i, x_j) = x_{i+j}$	I
Min-stuffle [7]	$x_i u \sqcup x_j v = x_i(u \sqcup x_j v) + x_j(x_i u \sqcup v) - x_{i+j}(u \sqcup v)$	$\varphi(x_i, x_j) = -x_{i+j}$	III
Muffle [14]	$x_i u \sqcup x_j v = x_i(u \sqcup x_j v) + x_j(x_i u \sqcup v) + x_{i \times j}(u \sqcup v)$	$\varphi(x_i, x_j) = x_{i \times j}$	I
$q$ -shuffle [3]	$x_i u \sqcup_q x_j v = x_i(u \sqcup_q x_j v) + x_j(x_i u \sqcup_q v) + qx_{i+j}(u \sqcup_q v)$	$\varphi(x_i, x_j) = qx_{i+j}$	III
$q$ -shuffle <sub>2</sub>	$x_i u \sqcup_q x_j v = x_i(u \sqcup_q x_j v) + x_j(x_i u \sqcup_q v) + q^{i \cdot j} x_{i+j}(u \sqcup_q v)$	$\varphi(x_i, x_j) = q^{i \cdot j} x_{i+j}$	II
LDIAG(1, $q_s$ ) [10] (non-crossed, non-shifted)	$au \sqcup bv = a(u \sqcup bv) + b(au \sqcup v) + q_s^{ a  b } a \cdot b(u \sqcup v)$	$\varphi(a, b) = q_s^{ a  b }(a \cdot b)$	II
$q$ -Infiltration [12]	$au \uparrow bv = a(u \uparrow bv) + b(au \uparrow v) + q\delta_{a,b} a(u \uparrow v)$	$\varphi(a, b) = q\delta_{a,b}$	III
AC-stuffle	$au \sqcup_\varphi bv = a(u \sqcup_\varphi bv) + b(au \sqcup_\varphi v) + \varphi(a, b)(u \sqcup_\varphi v)$	$\varphi(a, b) = \varphi(b, a)$ $\varphi(\varphi(a, b), c) = \varphi(a, \varphi(b, c))$	IV
Semigroup-stuffle	$x_t u \sqcup_{\perp} x_s v = x_t(u \sqcup_{\perp} x_s v) + x_s(x_t u \sqcup_{\perp} v) + x_{t \perp s}(u \sqcup_{\perp} v)$	$\varphi(x_t, x_s) = x_{t \perp s}$	I
$\varphi$ -shuffle	$au \sqcup_\varphi bv = a(u \sqcup_\varphi bv) + b(au \sqcup_\varphi v) + \varphi(a, b)(u \sqcup_\varphi v)$	$\varphi(a, b)$ law of AAU	V

Of course, the  $q$ -shuffle is equal to the (classical) shuffle when  $q = 0$ . As for the  $q$ -infiltration when  $a = 1$  one recovers the infiltration product defined in [6].

# A useful property

## Proposition B

Let  $\mathcal{B} = (k\langle X \rangle, \text{conc}, 1_{X^*}, \Delta, \epsilon)$  be a conc-bialgebra, then

- 1 The space  $k^{\text{rat}}\langle X \rangle$  is closed by the convolution product  $\diamond$  (here  ${}^t\Delta$ ) given by

$$\langle S \diamond T \mid w \rangle = \langle S \otimes T \mid \Delta(w) \rangle \quad (4)$$

- 2 If  $k$  is a  $\mathbb{Q}$ -algebra and  $\Delta_+ : k.X \rightarrow k.X \otimes k.X$  cocommutative,  $\mathcal{B}$  is an enveloping algebra iff  $\Delta_+$  is moderate<sup>a</sup>.
- 3 If, moreover  $k$  is without zero divisors, the characters  $(x^*)_{x \in X}$  are algebraically independant over  $(k\langle X \rangle, \diamond, 1_{X^*})$  within  $(k\langle\langle X \rangle\rangle, \diamond, 1_{X^*})$ .

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<sup>a</sup>See CAP 2017

# A useful property/2

## Independence of characters with respect to polynomials

▲ I came across the following property :

5  $\blacktriangledown$  Let  $\mathfrak{g}$  be a Lie algebra over a ring  $k$  without zero divisors,  $\mathcal{U} = \mathcal{U}(\mathfrak{g})$  be its enveloping algebra. As such,  $\mathcal{U}$  is a Hopf algebra and  $\epsilon$ , its counit, is the only character of  $\mathcal{U} \rightarrow k$  which vanishes on  $\mathfrak{g}$ .

☆ Set  $\mathcal{U}_+ = \ker(\epsilon)$ . We build the following filtrations ( $N \geq 1$ )

1 
$$\mathcal{U}_N = \mathcal{U}_+^N = \underbrace{\mathcal{U}_+ \dots \mathcal{U}_+}_{N \text{ times}} \quad (1)$$

and

$$\mathcal{U}_N^* = \mathcal{U}_{N+1}^\perp = \{f \in \mathcal{U}^* \mid (\forall u \in \mathcal{U}_{N+1})(f(u) = 0)\} \quad (2)$$

the first one is decreasing and the second one increasing. One shows easily that (with  $\diamond$  as the convolution product)

$$\mathcal{U}_p^* \diamond \mathcal{U}_q^* \subset \mathcal{U}_{p+q}^*$$

so that  $\mathcal{U}_\infty^* = \bigcup_{n \geq 1} \mathcal{U}_n^*$  is a convolution subalgebra of  $\mathcal{U}^*$ .

Now, we can state the

**Theorem** : The set of characters of  $(\mathcal{U}, \cdot, \mathbb{1}_{\mathcal{U}})$  is linearly free w.r.t.  $\mathcal{U}_\infty^*$ .

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## A useful property/3

### Remark

Property (3) is no longer true if  $\Delta$  is not moderate. For example with the Hadamard coproduct and  $x \neq y$ , one has  $y \odot (x)^* = 0$ .

### Examples

**Shuffle:**  $(\alpha x)^* \sqcup (\beta y)^* = (\alpha x + \beta y)^*$

**Stuffle:**  $(\alpha y_i)^* \sqcup (\beta y_j)^* = (\alpha y_i + \beta y_j + \alpha \beta y_{i+j})^*$

**$q$ -infiltration:**  $(\alpha x)^* \uparrow_q (\beta y)^* = (\alpha x + \beta y + \alpha \beta \delta_{x,y} x)^*$

**Hadamard:**  $(\alpha a)^* \odot (\beta b)^* = 1_{X^*}$  if  $a \neq b$  and  $(\alpha a)^* \odot (\beta a)^* = (\alpha \beta a)^*$



## Starting the ladder

$$\begin{array}{ccc}
 (\mathbb{C}\langle X \rangle, \sqcup, \mathbf{1}_{X^*}) & \xrightarrow{\text{Li}\bullet} & \mathbb{C}\{\text{Li}_w\}_{w \in X^*} \\
 \downarrow & & \downarrow \\
 (\mathbb{C}\langle X \rangle, \sqcup, \mathbf{1}_{X^*})[x_0^*, (-x_0)^*, x_1^*] & \xrightarrow{\text{Li}\bullet^{(1)}} & \mathcal{C}_{\mathbb{Z}}\{\text{Li}_w\}_{w \in X^*}
 \end{array}$$

## Domain of Li (definition)

In order to extend Li to series, we define  $Dom(Li; \Omega)$  (or  $Dom(Li)$  if the context is clear) as the set of series  $S = \sum_{n \geq 0} S_n$  (decomposition by homogeneous components) such that  $\sum_{n \geq 0} Li_{S_n}(z)$  converges for the compact convergence in  $\Omega$ . One sets

$$Li_S(z) := \sum_{n \geq 0} Li_{S_n}(z) \quad (5)$$

## Examples

$$Li_{x_0^*}(z) = z, \quad Li_{x_1^*}(z) = (1 - z)^{-1}; \quad Li_{\alpha x_0^* + \beta x_1^*}(z) = z^\alpha (1 - z)^{-\beta}$$

# Properties of the extended Li

## Proposition

With this definition, we have

- 1  $Dom(Li)$  is a shuffle subalgebra of  $\mathbb{C}\langle\langle X \rangle\rangle$  and then so is  $Dom^{rat}(Li) := Dom(Li) \cap \mathbb{C}^{rat}\langle\langle X \rangle\rangle$
- 2 For  $S, T \in Dom(Li)$ , we have

$$Li_{S \sqcup T} = Li_S \cdot Li_T$$

## Examples and counterexamples

For  $|t| < 1$ , one has  $(tx_0)^*x_1 \in Dom(Li, D)$  ( $D$  is the open unit slit disc), whereas  $x_0^*x_1 \notin Dom(Li, D)$ .

Indeed, we have to examine the convergence of  $\sum_{n \geq 0} Li_{x_0^n x_1}(z)$ , but, for  $z \in ]0, 1[$ , one has  $0 < z < Li_{x_0^n x_1}(z) \in \mathbb{R}$  and therefore, for these values  $\sum_{n \geq 0} Li_{x_0^n x_1}(z) = +\infty$ .

## Coefficients in the Ladder

$$\begin{array}{ccc}
 (\mathbb{C}\langle X \rangle, \sqcup, 1_{X^*}) & \xleftarrow{\text{Li}_\bullet} & \mathbb{C}\{\text{Li}_w\}_{w \in X^*} \\
 \downarrow & & \downarrow \\
 (\mathbb{C}\langle X \rangle, \sqcup, 1_{X^*})[x_0^*, (-x_0)^*, x_1^*] & \xrightarrow{\text{Li}_\bullet^{(1)}} & \mathcal{C}_{\mathbb{Z}}\{\text{Li}_w\}_{w \in X^*} \\
 \downarrow & & \downarrow \\
 \mathbb{C}\langle X \rangle \sqcup \mathbb{C}^{\text{rat}}\langle\langle x_0 \rangle\rangle \sqcup \mathbb{C}^{\text{rat}}\langle\langle x_1 \rangle\rangle & \xrightarrow{\text{Li}_\bullet^{(2)}} & \mathcal{C}_{\mathbb{C}}\{\text{Li}_w\}_{w \in X^*}
 \end{array}$$

Were, for every additive subgroup  $(H, +) \subset (\mathbb{C}, +)$ ,  $\mathcal{C}_H$  has been set to the following subring of  $\mathbb{C}$

$$\mathcal{C}_H := \mathbb{C}\{z^\alpha(1-z)^{-\beta}\}_{\alpha, \beta \in H} . \tag{6}$$

## Examples

$$\text{Li}_{x_0^*}(z) = z, \quad \text{Li}_{x_1^*}(z) = (1-z)^{-1}; \quad \text{Li}_{\alpha x_0^* + \beta x_1^*}(z) = z^\alpha(1-z)^{-\beta}$$

# The arrow $\text{Li}_{\bullet}^{(1)}$

## Proposition

- i. The family  $\{x_0^*, x_1^*\}$  is algebraically independent over  $(\mathbb{C}\langle X \rangle, \sqcup, 1_{X^*})$  within  $(\mathbb{C}\langle\langle X \rangle\rangle^{\text{rat}}, \sqcup, 1_{X^*})$ .
- ii.  $(\mathbb{C}\langle X \rangle, \sqcup, 1_{X^*})[x_0^*, x_1^*, (-x_0)^*]$  is a free module over  $\mathbb{C}\langle X \rangle$ , the family  $\{(x_0^*)^{\sqcup k} \sqcup (x_1^*)^{\sqcup l}\}_{(k,l) \in \mathbb{Z} \times \mathbb{N}}$  is a  $\mathbb{C}\langle X \rangle$ -basis of it.
- iii. As a consequence,  $\{w \sqcup (x_0^*)^{\sqcup k} \sqcup (x_1^*)^{\sqcup l}\}_{\substack{w \in X^* \\ (k,l) \in \mathbb{Z} \times \mathbb{N}}}$  is a  $\mathbb{C}$ -basis of it.
- iv.  $\text{Li}_{\bullet}^{(1)}$  is the unique morphism from  $(\mathbb{C}\langle X \rangle, \sqcup, 1_{X^*})[x_0^*, (-x_0)^*, x_1^*]$  to  $\mathcal{H}(\Omega)$  such that

$$x_0^* \rightarrow z, \quad (-x_0)^* \rightarrow z^{-1} \quad \text{and} \quad x_1^* \rightarrow (1 - z)^{-1}$$

- v.  $\text{Im}(\text{Li}_{\bullet}^{(1)}) = \mathcal{C}_{\mathbb{Z}}\{\text{Li}_w\}_{w \in X^*}$ .
- vi.  $\ker(\text{Li}_{\bullet}^{(1)})$  is the (shuffle) ideal generated by  $x_0^* \sqcup x_1^* - x_1^* + 1_{X^*}$ .

## Sketch of the proof for vi.

Let  $\mathcal{J}$  be the ideal generated by  $x_0^* \sqcup x_1^* - x_1^* + 1_{X^*}$ . It is easily checked, from the following formulas<sup>a</sup>, for  $k \geq 1$ ,

$$\begin{aligned}w \sqcup x_0^* \sqcup (x_1^*)^{\sqcup k} &\equiv w \sqcup (x_1^*)^{\sqcup k} - w \sqcup (x_1^*)^{\sqcup k-1} [\mathcal{J}], \\w \sqcup (-x_0)^* \sqcup (x_1^*)^{\sqcup k} &\equiv w \sqcup (-x_0)^* \sqcup (x_1^*)^{\sqcup k-1} + w \sqcup (x_1^*)^{\sqcup k} [\mathcal{J}],\end{aligned}$$

that one can rewrite  $[\text{mod } \mathcal{J}]$  any monomial  $w \sqcup (x_0^*)^{\sqcup l} \sqcup (x_1^*)^{\sqcup k}$  as a linear combination of such monomials with  $kl = 0$ . Observing that the image, through  $\text{Li}_{\bullet}^{(1)}$ , of the following family is free in  $\mathcal{H}(\Omega)$

$$\left\{ w \sqcup (x_1^*)^{\sqcup l} \sqcup (x_0^*)^{\sqcup k} \right\}_{(w,l,k) \in (X^* \times \mathbb{N} \times \{0\}) \sqcup (X^* \times \{0\} \times \mathbb{Z})} \quad (7)$$

we get the result.

---

<sup>a</sup>In the Figure below,  $(w, l, k)$  codes the element  $w \sqcup (x_0^*)^{\sqcup l} \sqcup (x_1^*)^{\sqcup k}$ .

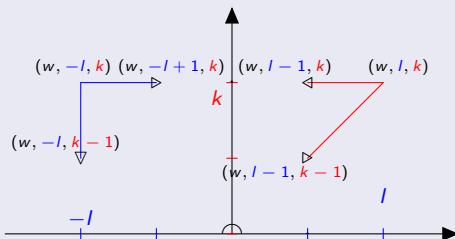


Figure: Rewriting mod  $\mathcal{J}$  of  $\{w \sqcup (x_0^*) \sqcup l \sqcup (x_1^*) \sqcup k\}_{k \in \mathbb{N}, l \in \mathbb{Z}, w \in X^*}$ .

## End of the ladder: pushing coefficients to $\mathcal{C}_{\mathbb{C}}$

$$\begin{array}{ccc}
 (\mathbb{C}\langle X \rangle, \sqcup, 1_{X^*}) & \xleftarrow{\text{Li}\bullet} & \mathbb{C}\{\text{Li}_w\}_{w \in X^*} \\
 \downarrow & & \downarrow \\
 (\mathbb{C}\langle X \rangle, \sqcup, 1_{X^*})[x_0^*, (-x_0)^*, x_1^*] & \xrightarrow{\text{Li}\bullet^{(1)}} & \mathbb{C}_{\mathbb{Z}}\{\text{Li}_w\}_{w \in X^*} \\
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 \end{array}$$

### Exchangeable (rational) series

The power series  $S$  belongs to  $\mathbb{C}_{\text{exc}}\langle X \rangle$ , iff

$$(\forall u, v \in X^*)((\forall x \in X)(|u|_x = |v|_x) \Rightarrow \langle S|u \rangle = \langle S|v \rangle). \quad (8)$$

We will note  $\mathbb{C}_{\text{exc}}^{\text{rat}}\langle X \rangle$ , the set of exchangeable rational series i.e.

$$\mathbb{C}_{\text{exc}}^{\text{rat}}\langle X \rangle := \mathbb{C}_{\text{exc}}\langle X \rangle \cap \mathbb{C}^{\text{rat}}\langle X \rangle \quad (9)$$

## Lemma (D., HNM, Ngô, 2016)

$$\textcircled{1} \mathbb{C}_{\text{exc}}^{\text{rat}} \langle\langle X \rangle\rangle := \mathbb{C}^{\text{rat}} \langle\langle X \rangle\rangle \cap \mathbb{C}_{\text{exc}} \langle\langle X \rangle\rangle = \mathbb{C}^{\text{rat}} \langle\langle x_0 \rangle\rangle \sqcup \mathbb{C}^{\text{rat}} \langle\langle x_1 \rangle\rangle.$$

$\textcircled{2}$  For any  $x \in X$ , from a theorem by Kronecker, one has  $\mathbb{C}^{\text{rat}} \langle\langle x \rangle\rangle = \text{span}_{\mathbb{C}} \{(ax)^* \sqcup \mathbb{C} \langle x \rangle \mid a \in \mathbb{C}\}$  and

$$\{(ax)^* \sqcup x^n\}_{(a,n) \in \mathbb{C} \times \mathbb{N}} \quad (10)$$

is a basis of it. When restricted to  $(\mathbb{C}^* \times \mathbb{N}) \cup \{(0,0)\}$  this family spans  $\mathbb{C}_{\text{const}}^{\text{rat}} \langle\langle x \rangle\rangle$  (fractions being constant at infinity)

$$\textcircled{3} \mathbb{C} \langle X \rangle \sqcup \mathbb{C}_{\text{exc}}^{\text{rat}} \langle\langle X \rangle\rangle \simeq \mathbb{C} \langle X \rangle \otimes_{\mathbb{C}} \mathbb{C}_{\text{const}}^{\text{rat}} \langle\langle x_0 \rangle\rangle \otimes_{\mathbb{C}} \mathbb{C}_{\text{const}}^{\text{rat}} \langle\langle x_1 \rangle\rangle$$

$$\textcircled{4} \text{Im}(\text{Li}_{\bullet}^{(2)}) = \mathcal{C}_{\mathbb{C}} \{\text{Li}_w\}_{w \in X^*}.$$

$\textcircled{5}$   $\ker(\text{Li}_{\bullet}^{(2)})$  is the (shuffle) ideal generated by  $x_0^* \sqcup x_1^* - x_1^* + 1_{X^*}$  (prospective).



# Concluding remarks/1

- 1 We have coded classical (and extended) polylogarithms with words obtaining a Noncommutative generating series which is a shuffle character
- 2 This character can be extended by continuity to certain series forming a shuffle subalgebra of Noncommutative formal power series.
- 3 We have found some remarkable subalgebras of  $Dom^{rat}(Li)$ , given their bases and described the kernel of the so extended  $Li_{\bullet}$ .
- 4 Definition of  $Dom(Li)$  and  $Dom^{rat}(Li)$  have to be refined and their exploration pushed further.
- 5 Combinatorics of discrete Dyson integrals for various sets of differential forms has to be implemented

## Concluding remarks/2

- 6 Drinfeld-Kohno Lie algebras i.e. algebras presented by

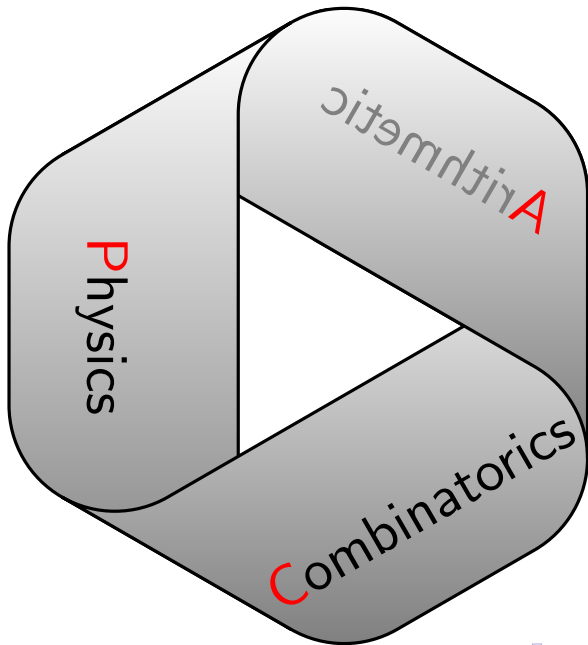
$$DK(A; k) = \langle A \times A ; \mathbf{R}_A \rangle_{k\text{-Lie algebras}} \quad (11)$$

with  $\mathbf{R}_A$ , the relator

$$\mathbf{R}_A = \begin{cases} (a, a) & = 0 \text{ for } a \in A \\ (a, b) & = (b, a) \text{ for } a, b \in A \\ [(a, c), (a, b) + (b, c)] & = 0 \text{ for } |\{a, b, c\}| = 3, \\ [(a, b), (c, d)] & = 0 \text{ for } |\{a, b, c, d\}| = 4 \end{cases} \quad (12)$$

can be decomposed in several ways as a direct sum of Free Lie algebras giving rise to product of MRS factorisations

$$\chi = \prod_{I \in \mathcal{Lyn}(X)} e^{\chi(S_I) P_I} \quad (13)$$



THANK YOU FOR YOUR ATTENTION !